

A mathematician's vocabulary

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Chapter 1

1.1 Introduction

1.1.1 Structures on a set

1.1.1.1

A very general class of mathematical structures is obtained by equipping a set X with one or more subsets $\Gamma \subseteq F(X)$ where $F(X)$ is a set *naturally* associated with set X . ‘Naturally’ here means that any map $f: X \rightarrow Y$ induces a map

$$f_*: F(X) \rightarrow F(Y) \quad (1.1)$$

or a map

$$f^*: F(Y) \rightarrow F(X) \quad (1.2)$$

1.1.1.2

In the first case we expect that

$$(f \circ g)_* = f_* \circ g_*, \quad (1.3)$$

and we speak of *covariant* dependence on X , in the second case we require that

$$(f \circ g)^* = g^* \circ f^*, \quad (1.4)$$

and we speak of *contravariant* dependence on X .

1.1.1.3

In modern Mathematics, such associations are called *covariant* and *contravariant functors* from the category of sets to the category of sets.

1.1.2 A few examples of such functors

1.1.2.1 Cartesian powers

Given a set I , consider the correspondence that associates with a set X its I -th Cartesian power

$$X \rightsquigarrow X^I := \{(x_i)_{i \in I} \mid x_i \in X\}. \quad (1.5)$$

The Cartesian power is a covariant functor, a map $f: X \rightarrow Y$ induces the map

$$f_*: X^I \rightarrow Y^I, \quad f_*((x_i)_{i \in I}) := (f(x_i))_{i \in I}. \quad (1.6)$$

1.1.2.2 Exponents

Given a set A , consider the correspondence that associates with a set X the set of maps from X to A

$$X \rightsquigarrow A^X := \{\phi: X \rightarrow A\}. \quad (1.7)$$

This functor is contravariant:

$$f^*: A^Y \rightarrow A^X, \quad f^*(\phi) := \phi \circ f. \quad (1.8)$$

1.1.2.3 The power set as a covariant functor

This is the functor that associates with a set X the set $\mathcal{P}(X)$ of all of its subsets and, with a map $f: X \rightarrow Y$, the map $f_*: \mathcal{P}(X) \rightarrow \mathcal{P}(Y)$ that sends a subset $A \subseteq X$ to its *image* under f ,

$$f(A) := \{y \in Y \mid y = f(x) \text{ for some } x \in A\}.$$

1.1.2.4 The power set as a contravariant functor

This functor associates with a set X , the same set $\mathcal{P}(X)$, and with $f: X \rightarrow Y$, the map $f^*: \mathcal{P}(Y) \rightarrow \mathcal{P}(X)$ that sends a subset $B \subseteq Y$ to its *preimage* under f ,

$$f^{-1}(B) := \{x \in X \mid f(x) \in B\}.$$

1.1.2.5

For any $A \subseteq X$ and $B \subseteq Y$, one has

$$f(A) \subseteq B \quad \text{if and only if} \quad A \subseteq f^{-1}(B). \quad (1.9)$$

This means that the pair of maps (f_*, f^*) forms a *Galois connection* between partially ordered sets $(\mathcal{P}(X), \subseteq)$ and $(\mathcal{P}(Y), \subseteq)$ (cf. *Notes on Partially Ordered Sets*).

1.1.2.6

For any set X , there exists a *natural bijection*¹

$$\chi^X: \mathcal{P}(X) \rightarrow 2^X, \quad A \mapsto \chi_A^X, \quad (1.10)$$

where

$$\chi_A^X(x) := \begin{cases} 1 & \text{if } x \in A \\ 0 & \text{otherwise} \end{cases} \quad (1.11)$$

is the *characteristic* function of a subset $A \subseteq X$. In the interest of simplifying notation when possible, the superscript X is dropped when X is clear from the context.

1.1.2.7

‘Naturality’ of (1.10) means that, given a map $f: X \rightarrow Y$, the following diagram commutes,

$$\begin{array}{ccc} \mathcal{P}(X) & \xleftarrow{f^*} & \mathcal{P}(Y) \\ \chi^X \downarrow & & \downarrow \chi^Y \\ 2^X & \xleftarrow{f^*} & 2^Y \end{array}, \quad (1.12)$$

i.e., the composition of arrows either way produces the same result

$$\chi^X \circ f^* = f^* \circ \chi^Y.$$

In categorical language, we could say that χ is a *natural transformation* of the contravariant power-set functor $\mathcal{P}(\)$ into the exponent functor $2^{(\)}$ (in this case an *isomorphism* of functors, since all the maps χ^X are isomorphisms in the category of sets, i.e., they are invertible maps).

¹In the language of sets, $0 = \emptyset$ and $n = \{0, \dots, n-1\}$.

1.1.2.8

Besides the category of sets there are other categories of interest in Mathematics, and there exist several interesting functors between them. Categorical language allows one to see various ‘natural’ constructions in a clear light, and it facilitates noticing connections between seemingly distant concepts and subjects. For this reason, it became very popular in modern Mathematics to the point of being indispensable, and a ‘must-learn’ for a beginner. We shall use it too.

1.1.2.9

You are encouraged to familiarize yourself with the language of categories and functors as soon as possible and, after mastering the basics of categorical grammar, to learn also at least the concepts of an equivalence of categories and of a pair of adjoint functors, and study numerous fundamentally important examples these two concepts. To facilitate this, I include the most basic definitions below.

Like with any language, acquiring proficiency requires constant use, so you, after learning the basic concepts, should be constantly observing these concepts at work in various branches of Mathematics.

1.2 First terms in the vocabulary

1.2.1 Families

1.2.1.1 Families of sets

The term a *family of sets* is used in two meanings: as a subset $\mathcal{X} \subseteq \mathcal{P}(U)$ of the power set of some set U or, as a map

$$I \rightarrow \mathcal{P}(U), \quad i \mapsto X_i,$$

which assigns a set X_i to every element $i \in I$ of certain set I . In the latter case we speak of a family of subsets of U *indexed* by set I . The indexing set can be arbitrary and it may come equipped with additional structure like ordering.

1.2.1.2 Notation

It is customary to denote indexed families by $(X_i)_{i \in I}$.

1.2.1.3

A family of subsets of U viewed as a subset of $\mathcal{P}(U)$ is conceptually simpler, as its definition does not rely on the notion of a map yet it can be viewed as a special case of an indexed family, namely as a family indexed by itself:

$$\mathcal{X} \rightarrow \mathcal{P}(U), \quad X \mapsto X.$$

1.2.1.4 Families of elements of a set

A family of elements of a set X will be always used in the sense of a family indexed by some set I . By definition it is a map

$$I \rightarrow X, \quad i \mapsto x_i.$$

Conceptually, there is no difference between a *family of elements* of X and a *map* $I \rightarrow X$. The difference is exclusively in notation and in the points of emphasis.

In the language of *families of elements* the focus is on X and its elements. The nature of the indexing set is secondary and generally not very important.

In the language of *maps*, the source and the target of a map are on equal footing, and the map itself is usually sufficiently important to merit its own symbol in notation.

1.2.1.5 Natural numbers

We shall frequently identify natural numbers with the sets:

$$0 := \emptyset, \quad 1 := \{0\}, \quad 2 := \{0, 1\}, \quad \dots, \quad n := \{0, \dots, n-1\}, \quad \dots \quad (1.13)$$

Exercise 1 Show that $2 \neq 3$.

1.2.1.6 Sequences

Families indexed by subsets of the set of natural numbers or, more generally, by ordered countable sets, are called *sequences*.

1.2.1.7 n -tuples

Families indexed by $I = \{1, \dots, n\}$ are called *ordered n -tuples* of elements of X , and notation

$$(x_1, \dots, x_n) \quad \text{instead of} \quad (x_i)_{i \in \{1, \dots, n\}}$$

is generally used. Ordered 2-, 3-, 4-tuples are respectively called *ordered pairs, triples, quadruples*.

1.2.2 Rings of sets and algebras of subsets

1.2.2.1 Rings of sets

A nonempty family of sets \mathcal{R} , i.e., a nonempty set whose elements are sets, is said to be a *ring of sets* if the union, $R \cup R'$, and the difference, $R \setminus R'$, belongs to \mathcal{R} for any $R, R' \in \mathcal{R}$.

Exercise 2 Show that in every ring of sets \mathcal{R} one has

$$R \cap R' \in \mathcal{R} \quad \text{for any} \quad R, R' \in \mathcal{R}.$$

1.2.2.2 Algebras of subsets

A nonempty family $\mathcal{A} \subseteq \mathcal{P}(X)$ of subsets of a set X is said to be an *algebra of subsets of X* if the intersection, $A \cap A'$, and the complement, $A^c := X \setminus A$, belongs to \mathcal{A} for any $A, A' \in \mathcal{A}$.

Exercise 3 Show that $\mathcal{A} \subseteq \mathcal{P}(X)$ is an algebra of subsets of a set X if and only if \mathcal{A} is a ring of sets which contains X .

1.2.2.3

The family of all *finite* subsets $\mathcal{P}_{\text{fin}}(X)$ of a set X is a ring of sets which is an algebra of subsets of X if and only if X is finite.

1.2.3 Operations involving families of sets

1.2.3.1 Union

The union of a family $\mathcal{X} \subseteq \mathcal{P}(U)$ is the set

$$\{u \in U \mid u \in X \text{ for some } X \in \mathcal{X}\}. \quad (1.14)$$

This set is denoted

$$\bigcup \mathcal{X} \quad \text{or} \quad \bigcup_{X \in \mathcal{X}} X. \quad (1.15)$$

1.2.3.2

The union of an indexed family $(X_i)_{i \in I}$ is defined similarly

$$\bigcup_{i \in I} X_i := \{u \in U \mid u \in X_i \text{ for some } i \in I\}. \quad (1.16)$$

1.2.3.3 Intersection

The intersection a family $\mathcal{X} \subseteq \mathcal{P}(U)$ is the set

$$\{u \in U \mid u \in X \text{ for every } X \in \mathcal{X}\}. \quad (1.17)$$

This set is denoted

$$\bigcap \mathcal{X} \quad \text{or} \quad \bigcap_{X \in \mathcal{X}} X. \quad (1.18)$$

1.2.3.4

The intersection of an indexed family $(X_i)_{i \in I}$ is defined similarly

$$\bigcap_{i \in I} X_i := \{u \in U \mid u \in X_i \text{ for every } i \in I\}. \quad (1.19)$$

Exercise 4 Show that the intersection

$$\bigcap_{i \in I} \mathcal{R}_i$$

of any family of rings of sets $(\mathcal{R}_i)_{i \in I}$ is a ring of sets. Likewise, show that the intersection

$$\bigcap_{i \in I} \mathcal{A}_i$$

of any family of algebras of subsets $(\mathcal{A}_i)_{i \in I}$ of a given set X is an algebra of subsets of X .

1.2.3.5 Cartesian product

The Cartesian product of an indexed family $(X_i)_{i \in I}$ is the set of all families $\xi = (x_i)_{i \in I}$ of elements of $\bigcup_{i \in I} X_i$ such that $x_i \in X_i$:

$$\prod_{i \in I} X_i := \{(x_i)_{i \in I} \mid x_i \in X_i\}. \quad (1.20)$$

Equivalently,

$$\prod_{i \in I} X_i := \left\{ \xi: I \rightarrow \bigcup_{i \in I} X_i \mid \xi(i) \in X_i \right\}. \quad (1.21)$$

1.2.3.6 Notation

The Cartesian product of a finite family (X_1, \dots, X_n) is usually denoted

$$X_1 \times \cdots \times X_n. \quad (1.22)$$

1.2.3.7 Comment

It is important to observe that one can replace $\bigcup_{i \in I} X_i$ in the definition of the Cartesian product by *any* set that contains all X_i . The corresponding ‘products’ will be essentially identical sets. This is due to the observation that there exists a canonical identification between maps $A \rightarrow B$ whose image is contained in a subset $B' \subseteq B$, and maps $A \rightarrow B'$.

1.2.3.8 Canonical projections

The Cartesian product comes equipped with the family of surjective maps,

$$\pi_i: \prod_{j \in I} X_j \rightarrow X_i \quad \xi \mapsto x_i \quad (i \in I), \quad (1.23)$$

which send a map $\xi: I \rightarrow \bigcup_{i \in I} X_i$ to its value at each i . When $I = \{1, \dots, n\}$, then π_i is the i -th coordinate map

$$\pi_i: (x_1, \dots, x_n) \mapsto x_i \quad (i = 1, \dots, n).$$

1.2.3.9 A universal property of the Cartesian product

Given any set Y and a family $(f_i)_{i \in I}$ of maps $f_i: Y \rightarrow X_i$, there exists a unique map $\tilde{f}: Y \rightarrow \prod_{i \in I} X_i$ such that

$$f_i = \pi_i \circ \tilde{f} \quad (i \in I). \quad (1.24)$$

Exercise 5 Verify that the map

$$\tilde{f}: y \mapsto (f_i(y))_{i \in I} \quad (y \in Y) \quad (1.25)$$

satisfies (1.24), and that any map $g: Y \rightarrow \prod_{i \in I} X_i$ which satisfies (1.24) coincides with \tilde{f} .

1.2.3.10 Disjoint unions of sets

The disjoint union of an indexed family $(X_i)_{i \in I}$ should be thought of as the union of all sets X_i except that we keep as many distinct ‘copies’ of an element $x \in \bigcup_{i \in I} X_i$ as there are sets X_i which contain x . We achieve this by ‘tagging’ every element in $\bigcup_{i \in I} X_i$ by the index of the set it belongs to:

$$\coprod_{i \in I} X_i := \{(i, x) \in I \times \bigcup_{i \in I} X_i \mid x \in X_i\}. \quad (1.26)$$

1.2.3.11 Notation

The disjoint union of a finite family (X_1, \dots, X_n) is usually denoted

$$X_1 \sqcup \dots \sqcup X_n. \quad (1.27)$$

Exercise 6 Denote by p the composition of the inclusion map and the canonical projection

$$\coprod_{i \in I} X_i \hookrightarrow I \times \bigcup_{i \in I} X_i \rightarrow \bigcup_{i \in I} X_i. \quad (1.28)$$

Show that p is surjective. Show that the fiber $p^{-1}(x)$ at $x \in \bigcup_{i \in I} X_i$ is

$$p^{-1}(x) = \{(i, x) \mid x \in X_i\}.$$

In particular, $p^{-1}(x)$ is in on-to-one correspondence with the set

$$\{i \in I \mid x \in X_i\}.$$

1.2.3.12

It follows that the disjoint union of a family of sets $(X_i)_{i \in I}$ is canonically identified with their union if and only if sets X_i are disjoint for distinct $i \in I$:

$$X_i \cap X_j = \emptyset \quad (i \neq j).$$

1.2.3.13 Canonical inclusions

The disjoint union comes equipped with the family of injective maps,

$$\iota_i: X_i \rightarrow \coprod_{j \in I} X_j \quad x \mapsto (i, x) \quad (i \in I). \quad (1.29)$$

1.2.3.14 A universal property of the disjoint union

Given any set Y and a family $(f_i)_{i \in I}$ of maps $f_i: X_i \rightarrow Y$, there exists a unique map $\tilde{f}: \coprod_{i \in I} X_i \rightarrow Y$ such that

$$f_i = \tilde{f} \circ \iota_i \quad (i \in I). \quad (1.30)$$

Exercise 7 Verify that the map

$$\tilde{f}: (i, x) \mapsto f_i(x) \quad (i \in I; x \in X_i) \quad (1.31)$$

satisfies (1.30), and that any map $g: \coprod_{i \in I} X_i \rightarrow Y$ which satisfies (1.30) coincides with \tilde{f} .

1.2.3.15

Map p defined in (1.28) is precisely such universal map \tilde{f} for the family of inclusion maps

$$f_i: X_i \hookrightarrow \bigcup_{j \in I} X_j \quad (i \in I).$$

1.2.3.16

Note that the properties of the Cartesian product and of the disjoint union of a family of sets are *dual* to each other. We shall explain this concept of *duality* later.

1.3 Associativity properties of operations on families of sets

1.3.1 Associativity of union

1.3.1.1 The indexed families case

Suppose we have two families of sets

$$(X_i)_{i \in I} \quad \text{and} \quad (X_k)_{k \in K}.$$

The iterated union

$$\bigcup_{i \in I} X_i \cup \bigcup_{k \in K} X_k$$

and the union

$$\bigcup_{l \in I \sqcup K} X_l$$

are *equal* as sets. In the case of a pair of finite families (A_1, \dots, A_m) and (B_1, \dots, B_n) , this equality acquires the form

$$(A_1 \cup \dots \cup A_m) \cup (B_1 \cup \dots \cup B_n) = A_1 \cup \dots \cup A_m \cup B_1 \cup \dots \cup B_n.$$

1.3.1.2 The total family

In general, given any family of families of sets

$$\left((X_{i_j})_{i_j \in I_j} \right)_{j \in J}, \quad (1.32)$$

the universal property of the disjoint union allows us to form the *total* family

$$(X_l)_{l \in L} \quad \text{where} \quad L = \coprod_{j \in J} I_j. \quad (1.33)$$

Indeed, regarding all sets to be subsets of a common set U , family of families of (1.32) is the same as a family of maps $I_j \rightarrow \mathcal{P}(U)_{j \in J}$ and, by the universal property of disjoint union, there exists a unique map $L \rightarrow \mathcal{P}(U)$ whose ‘restrictions’ to I_j are the component-families

$$(I_j \rightarrow \mathcal{P}(U))_{j \in J}.$$

We shall refer to $L \rightarrow \mathcal{P}(U)$ as the *total* family.

1.3.1.3

Now we are ready to make an observation about iterated unions of families. The following sets are equal

$$\bigcup_{j \in J} \bigcup_{i_j \in I_j} X_{i_j} = \bigcup_{l \in L} X_l.$$

Exercise 8 Formulate the corresponding associativity laws for intersection of families.

1.3.1.4 The nonindexed families case

There are two natural maps

$$\bigcup: \mathcal{P}(\mathcal{P}(\mathcal{P}(U))) \longrightarrow \mathcal{P}(\mathcal{P}(U)) \quad (1.34)$$

and

$$\bigcup_*: \mathcal{P}(\mathcal{P}(\mathcal{P}(U))) \longrightarrow \mathcal{P}(\mathcal{P}(U)). \quad (1.35)$$

The first one is the familiar *union-of-a-family* map applied to $\mathcal{P}(U)$ instead of U . It sends $\mathcal{X} \in \mathcal{P}(\mathcal{P}(\mathcal{P}(U)))$, i.e., a family of families of subsets of U , to the family of subsets of U which belong to at least one member family $\mathcal{X} \in \mathcal{X}$

$$\bigcup \mathcal{X} = \bigcup_{\mathcal{X} \in \mathcal{X}} \mathcal{X}.$$

The other one is induced by the map $\bigcup: \mathcal{P}(\mathcal{P}(U)) \rightarrow \mathcal{P}(U)$. It is formed by the *unions* of member families $\mathcal{X} \in \mathcal{X}$,

$$\bigcup_*(\mathcal{X}) := \{Y \subseteq U \mid Y = \bigcup \mathcal{X} = \bigcup_{X \in \mathcal{X}} X \text{ for some } \mathcal{X} \in \mathcal{X}\}.$$

Exercise 9 Show that the following diagram commutes

$$\begin{array}{ccc} \mathcal{P}(\mathcal{P}(\mathcal{P}(U))) & \xrightarrow{\bigcup_*} & \mathcal{P}(\mathcal{P}(U)) \\ \bigcup \downarrow & & \downarrow \bigcup \\ \mathcal{P}(\mathcal{P}(U)) & \xrightarrow{\bigcup} & \mathcal{P}(U) \end{array}, \quad (1.36)$$

i.e., show that the following two subsets of U ,

$$\bigcup \left(\bigcup \mathcal{X} \right) = \bigcup_{X \in \bigcup \mathcal{X}} X$$

and

$$\bigcup \bigcup_*(\mathcal{X}) = \bigcup_{\mathcal{X} \in \mathcal{X}} \left(\bigcup \mathcal{X} \right) = \bigcup_{\mathcal{X} \in \mathcal{X}} \left(\bigcup_{X \in \mathcal{X}} X \right),$$

are equal for any family of families $\mathcal{X} \subseteq \mathcal{P}(\mathcal{P}(U))$.

1.3.1.5 Terminology

If one is going to deal with “families of families of subsets of a set U ,” et caetera, on an extended basis, then one perhaps should use a less cumbersome terminology. One could, for example, call subsets of the n times iterated power set

$$\mathcal{P}^n(U) = \underbrace{\mathcal{P}(\dots \mathcal{P}(U) \dots)}_{n \text{ times}} \quad (n \geq 0)$$

n -families in a set U . In particular, subsets of $U = \mathcal{P}^0(U)$ are 0-families, families of subsets of U are 1-families, families of families of subsets of U are 2-families, etc.

1.3.2 Cartesian product

1.3.2.1 Associativity of Cartesian product

Instead of equality, we have only a canonical identification between the iterated Cartesian product of a family of families of sets and the Cartesian product of the total family.

Let us consider first the case of a pair of families of sets

$$(X_i)_{i \in I} \quad \text{and} \quad (X_k)_{k \in K}.$$

The natural correspondence

$$((x_i)_{i \in I}, (x_k)_{k \in K}) \leftrightarrow (x_l)_{l \in I \sqcup K}$$

identifies the iterated Cartesian product

$$\prod_{i \in I} X_i \times \prod_{k \in K} X_k$$

with

$$\prod_{l \in I \sqcup K} X_l.$$

In the case of a pair of finite families (A_1, \dots, A_m) and (B_1, \dots, B_n) , this identification acquires the form

$$((a_1, \dots, a_m), (b_1, \dots, b_n)) \leftrightarrow (a_1, \dots, a_m, b_1, \dots, b_n).$$

1.3.2.2 Iterated Cartesian product

In general, given any family of families of sets (1.32), the iterated Cartesian product and the Cartesian product of the total family are naturally identified

$$\prod_{j \in J} \prod_{i_j \in I_j} X_{i_j} \longleftrightarrow \prod_{l \in L} X_l \quad \text{where} \quad L = \coprod_{j \in J} I_j. \quad (1.37)$$

Indeed, elements of $\prod_{i_j \in I_j} X_{i_j}$ are maps

$$\xi_j : I_j \rightarrow \bigcup_{i_j \in I_j} X_{i_j}$$

such that $\xi_j(i_j) \in X_{i_j}$. By composing maps ξ_j with the inclusions

$$\bigcup_{i_j \in I_j} X_{i_j} \hookrightarrow U := \bigcup_{l \in L} X_l,$$

we can consider all ξ_j as being maps with the common target U . Thus, elements of

$$\prod_{j \in J} \prod_{i_j \in I_j} X_{i_j}$$

become families $(\xi_j)_{j \in J}$ of maps $\xi_j : I_j \rightarrow U$. By the universal property of the disjoint union, there exists a unique map $\tilde{\xi} : L \rightarrow U$ whose ‘restrictions’ to I_j are families maps $\xi_j : I_j \rightarrow U$.

This map $\tilde{\xi}$ is an element of $\prod_{l \in L} X_l$. Since the correspondence between families $(\xi_j)_{j \in J}$ and maps $\tilde{\xi}$ is bijective, the correspondence in (1.37) is bijective.

Exercise 10 *Formulate and prove the corresponding associativity laws for disjoint union.*

1.3.2.3 Calculus of Cartesian powers of a set

For any sets A , B , and C , one has natural identifications

$$A^B \times A^C \longleftrightarrow A^{B \sqcup C} \quad (1.38)$$

and, more generally,

$$\prod_{j \in J} A^{B_j} \longleftrightarrow A^{\coprod_{j \in J} B_j} \quad (1.39)$$

which are special cases of identifications (1.37).

1.3.2.4

One has also the following natural identification

$$(A^B)^C \longleftrightarrow A^{B \times C} \quad (1.40)$$

given by the following pair of mutually inverse correspondences

$$(A^B)^C \ni f \mapsto \phi \in A^{B \times C}, \quad \text{where} \quad \phi(b, c) := (f(c))(b)$$

and

$$A^{B \times C} \ni \phi \mapsto f \in (A^B)^C, \quad \text{where} \quad f(c) := \phi(\cdot, c).$$

1.3.2.5

Using the *families-of-elements* notation instead of *maps* notation, we can describe identification (1.40) also in this form

$$(X^I)^J \longleftrightarrow X^{I \times J}, \quad ((x_{ij})_{i \in I})_{j \in J} \leftrightarrow (x_{ij})_{(i,j) \in I \times J}. \quad (1.41)$$

1.4 The language of categories and functors

1.4.1 Oriented graphs

1.4.1.1

An *oriented graph* \mathcal{C} consists of two classes, \mathcal{C}_0 (the class of *vertices*) and \mathcal{C}_1 (the class of *arrows*) which are related by a pair of correspondences:

$$\begin{array}{ccc} & \mathcal{C}_1 & \\ t \swarrow & & \searrow s \\ \mathcal{C}_0 & & \mathcal{C}_0 \end{array} . \quad (1.42)$$

1.4.1.2

For any arrow α we shall refer to $s(\alpha)$ as its *source*, and to $t(\alpha)$ as its *target*.

1.4.1.3

Note that we are saying *classes*—not *sets*. Basic concepts of Category Theory impose on the foundations on which the edifice of Mathematics rests, that one is allowed to talk about classes that are not sets, like the class of all sets, the class of all singleton sets, the class of all vector spaces over a given field of coefficients, etc, and one is likewise allowed to talk about correspondences between classes as if they were mappings between sets.

We henceforth will be cautiously extending to classes certain terminology and notation usually associated with sets. For example, we may indicate that a is a vertex of a graph \mathcal{C} by writing either $a \in \mathcal{C}_0$ or $a \in \text{Vert } \mathcal{C}$. Similarly, we may indicate that α is an arrow of a graph \mathcal{C} by writing either $\alpha \in \mathcal{C}_1$ or $\alpha \in \text{Arr } \mathcal{C}$.

1.4.2 Categories

1.4.2.1 The class of composable arrows

Consider the class \mathcal{C}_2 of pairs (α_0, α_1) of arrows such that the source of α_1 is the target of α_0 . This class fits naturally into the diagram

$$\begin{array}{ccccc}
 & & \mathcal{C}_2 & & \\
 & p_0 \swarrow & & \searrow p_1 & \\
 & \mathcal{C}_1 & \circlearrowleft & \mathcal{C}_1 & \\
 t \swarrow & & & & \searrow s \\
 \mathcal{C}_0 & & \mathcal{C}_0 & & \mathcal{C}_0
 \end{array} \tag{1.43}$$

where p_i sends (α_0, α_1) to α_i .

1.4.2.2

A graph equipped with a correspondence

$$m: \mathcal{C}_2 \leftrightarrow \mathcal{C}_1 \tag{1.44}$$

is said to be a *category* if (1.44) is *associative*, i.e.,

$$(\alpha_0 \circ \alpha_1) \circ \alpha_2 = \alpha_0 \circ (\alpha_1 \circ \alpha_2) \tag{1.45}$$

for any *composable* triple of arrows. The latter means that

$$s(\alpha_0) = t(\alpha_1) \quad \text{and} \quad s(\alpha_1) = t(\alpha_2). \tag{1.46}$$

1.4.2.3 Objects and morphisms

1.4.2.4 $\text{Hom}_{\mathcal{C}}(a, b)$

It was observed early that if one requires in the definition of a category that, for any pair of objects $a, b \in \mathcal{C}_0$, morphisms with a as their source and with b as their target form a *set* and not just a class, then one can avoid essentially all the potential dangers arising from presence of classes in foundations of Category Theory.

This set is usually denoted $\text{Hom}_{\mathcal{C}}(a, b)$ and its elements are referred as morphisms from a to b .

1.4.2.5 The class of composable pairs of morphisms

We say that a pair (α, β) of morphisms is *composable* if $s(\alpha) = t(\beta)$. Denote by \mathcal{C}_2 the class of composable pairs of morphisms. We assume that a correspondence

$$m: \mathcal{C}_2 \rightarrow \mathcal{C}_1, \quad (\alpha, \beta) \mapsto \alpha \circ \beta, \quad (1.47)$$

is given. It is referred to as *composition* of morphisms, and is possibly the single most important element of the structure of a category.

1.4.2.6 The class of composable triples of morphisms

We say that a triple (α, β, γ) of morphisms is *composable* if $s(\alpha) = t(\beta)$ and $s(\beta) = t(\gamma)$. As can be expected, we denote the class of composable triples of morphisms by \mathcal{C}_3 . (Binary) composition (1.47) induces two correspondences $\mathcal{C}_3 \rightarrow \mathcal{C}_2$

$$m_1: (\alpha, \beta, \gamma) \mapsto (\alpha \circ \beta, \gamma) \quad \text{and} \quad m_2: (\alpha, \beta, \gamma) \mapsto (\alpha, \beta \circ \gamma). \quad (1.48)$$

By applying correspondence (1.47), we obtain two correspondences $\mathcal{C}_3 \rightarrow \mathcal{C}_1$. We require them to be equal which means that

$$(\alpha \circ \beta) \circ \gamma = \alpha \circ (\beta \circ \gamma) \quad (1.49)$$

for any composable triple of morphisms. This condition is called *associativity* of the composition of morphisms.

1.4.2.7

Associativity identity (1.49) can be expressed as commutativity of the following diagram

$$\begin{array}{ccc}
 \mathcal{C}_3 & \xrightarrow{m_1} & \mathcal{C}_2 \\
 m_2 \downarrow & & \downarrow m \\
 \mathcal{C}_2 & \xrightarrow{m} & \mathcal{C}_1
 \end{array} \tag{1.50}$$

1.4.2.8 The identity morphisms

We could stop here and call the defined structures *categories*. The classical and still a ‘default’ definition of a category additionally requires presence of a correspondence

$$i: \mathcal{C}_0 \rightarrow \mathcal{C}_1, \quad a \mapsto \text{id}_a \in \text{Hom}_{\mathcal{C}}(a, a), \tag{1.51}$$

such that

$$\alpha \circ \text{id}_a = \alpha \quad \text{and} \quad \text{id}_b \circ \alpha = \alpha \tag{1.52}$$

for any $\alpha \in \text{Hom}_{\mathcal{C}}(a, b)$. Morphism id_a is referred to as the *identity* morphism of object a .

1.4.2.9

Each of the identities in (1.52) can be expressed as commutativity of a diagram of correspondences:

$$\begin{array}{ccc}
 \mathcal{C}_1 \times \mathcal{C}_0 & \xrightarrow{(\text{id}_{\mathcal{C}_1}, i)} & \mathcal{C}_2 \\
 (\text{id}_{\mathcal{C}_1}, s) \uparrow & & \downarrow m \\
 \mathcal{C}_1 & \xrightarrow{\text{id}_{\mathcal{C}_1}} & \mathcal{C}_1
 \end{array} \quad \text{and} \quad \begin{array}{ccc}
 \mathcal{C}_0 \times \mathcal{C}_1 & \xrightarrow{(i, \text{id}_{\mathcal{C}_1})} & \mathcal{C}_2 \\
 (t, \text{id}_{\mathcal{C}_1}) \uparrow & & \downarrow m \\
 \mathcal{C}_1 & \xrightarrow{\text{id}_{\mathcal{C}_1}} & \mathcal{C}_1
 \end{array} \tag{1.53}$$

1.4.2.10

There are very good reasons not to require presence of the identity morphisms in general, and to call the categories that possess such morphisms—*unital* categories.

1.4.2.11 Isomorphisms

We say that a morphism $\alpha \in \text{Hom}_{\mathcal{C}}(a, b)$ is an *isomorphism* if there exists $\beta \in \text{Hom}_{\mathcal{C}}(b, a)$ such that

$$\alpha \circ \beta = \text{id}_b \quad \text{and} \quad \beta \circ \alpha = \text{id}_a. \quad (1.54)$$

Exercise 11 Show that if there exist morphisms $\beta, \gamma \in \text{Hom}_{\mathcal{C}}(b, a)$ such that

$$\alpha \circ \beta = \text{id}_b \quad \text{and} \quad \gamma \circ \alpha = \text{id}_a.$$

then $\beta = \gamma$.

1.4.2.12

In view of the above exercise, if there exists at least one *right* inverse and at least one *left* inverse for a morphism α , then they are equal, which implies that the two-sided inverse, (1.54), is unique when it exists. It is denoted α^{-1} .

1.4.2.13 Endomorphisms of an object

Morphisms $\alpha: a \rightarrow a$ are called *endomorphisms* of object a . The set $\text{Hom}_{\mathcal{C}}(a, a)$ is often denoted $\text{End}_{\mathcal{C}}(a)$.

1.4.2.14 Automorphisms of an object

Isomorphisms $\alpha: a \rightarrow a$ are called *automorphisms* of object a . The set of automorphisms is denoted $\text{Aut}_{\mathcal{C}}(a)$.

1.4.2.15 Symmetries

Before categorical language was proposed and developed as means to describe and study underlying structure of numerous areas of Mathematics, automorphisms of various objects: geometric, physical systems, etc—were often called *symmetries*.

1.4.2.16 Subcategories

For a category \mathcal{C} , suppose that, a pair of subclasses $\mathcal{C}'_0 \subseteq \mathcal{C}_0$ and $\mathcal{C}'_1 \subseteq \mathcal{C}_1$ is given such that the source and the target of any morphism in \mathcal{C}'_1 is

a member of \mathcal{C}'_0 and the composition of any two such morphisms is a member of \mathcal{C}'_1 .

If we equip the pair of classes $(\mathcal{C}_0, \mathcal{C}_1)$ with the source, target, and multiplication correspondences induced from category \mathcal{C} , we obtain a category on its own. Denote it \mathcal{C}' .

This situation arises frequently. We say that \mathcal{C}' is a *subcategory* of \mathcal{C} .

1.4.2.17 Full subcategories

If

$$\text{Hom}_{\mathcal{C}'}(a, b) = \text{Hom}_{\mathcal{C}}(a, b) \quad (a, b \in \mathcal{C}'_0),$$

then we say that \mathcal{C}' is a *full* subcategory of category \mathcal{C} .

1.4.3 Natural definitions of a morphism between sets

1.4.3.1 Set

The category of sets usually takes pride of being presented as the first example of a category. The objects of this category are sets. There are, however, several natural candidates for the morphisms. The standard choice for morphisms $X \rightarrow Y$ is to take maps $f: X \rightarrow Y$:

$$\text{Hom}_{\text{Set}}(X, Y) = Y^X.$$

This category will be denoted Set and referred to as *the* category of sets.

Note that isomorphisms in the category of sets coincide with the class of bijections.

1.4.3.2 Multivalued maps

A *multivalued* map, $\phi: X \multimap Y$, from a set X to a set Y , is a map $\phi: X \rightarrow \mathcal{P}(Y)$. Multivalued maps will be also called *multimaps*.

1.4.3.3 Maps versus multimaps

Every map $f: X \rightarrow Y$ defines the multimap

$$x \mapsto \phi_f(x) := \{f(x)\} \quad (x \in X).$$

The correspondence $f \mapsto \phi_f$ identifies maps $f: X \rightarrow Y$ with multimaps $\phi: X \multimap Y$ satisfying the property

$$|\phi(x)| = 1 \quad (x \in X). \quad (1.55)$$

1.4.3.4 The image map for a multimap

Every multimap $\phi: X \multimap Y$ naturally extends to a map $\mathcal{P}(X) \rightarrow \mathcal{P}(Y)$,

$$A \longmapsto \phi(A) := \bigcup_{x \in A} \phi(x) \quad (A \subseteq X). \quad (1.56)$$

We will continue to denote it ϕ and will call it the *image map* associated with multimap ϕ .

1.4.3.5 The reverse of a multimap

Every multimap $\phi: X \multimap Y$ also defines a multimap $Y \multimap X$

$$\phi^{\text{rev}}(y) := \{x \in X \mid \phi(x) \ni y\}. \quad (1.57)$$

We shall refer to it as the *reverse* of ϕ . When ϕ is a map $f: X \rightarrow Y$, then $\phi^{\text{rev}}(x) = \{x \in X \mid f(x) = y\}$ is called the *fiber* of f at $y \in Y$.

1.4.3.6 The preimage map for a multimap

The image map for the reverse multimap, ϕ^{rev} , will be called the *preimage map* for ϕ .

Exercise 12 Show that

$$\phi^{\text{rev}}(B) = \{x \in X \mid \phi(x) \cap B \neq \emptyset\} \quad (B \subseteq Y). \quad (1.58)$$

1.4.3.7 Composition of multimaps

Given multimaps $\phi: X \multimap Y$ and $\chi: Y \multimap Z$, their *composition*,

$$\chi \circ \phi: x \longmapsto \chi(\phi(x)) \quad (x \in X), \quad (1.59)$$

is a multimap $X \multimap Z$.

Exercise 13 Given maps $f: X \rightarrow Y$ and $g: Y \rightarrow Z$, show that

$$\phi_g \circ \phi_f = \phi_{g \circ f}. \quad (1.60)$$

Exercise 14 Show that composition of multimaps is associative, i.e.,

$$(\chi \circ \phi) \circ v = \chi \circ (\phi \circ v),$$

for any $v: W \multimap X$, $\phi: X \multimap Y$, and $\chi: Y \multimap Z$.

1.4.3.8 Set_{mult}

Thus, the class of sets equipped with multimaps as morphisms forms a category. We shall denote it Set_{mult} .

Exercise 15 Show that the canonical embedding $\iota_X: X \hookrightarrow \mathcal{P}(X)$,

$$\iota_X: x \longmapsto \{x\} \quad (x \in X)$$

is the identity endomorphism of set X in Set_{mult} .

1.4.3.9 Submaps

Let us call a multimap $\phi: X \rightarrow Y$ satisfying the condition

$$|\phi(x)| \leq 1 \quad (x \in X), \quad (1.61)$$

a *submap* (compare it with (1.55)).

If multimaps satisfying (1.55) correspond to maps $F: X \rightarrow Y$, then submaps correspond to *partially defined* maps from X to Y , i.e., to maps $f: X' \rightarrow Y$ whose domain is a subset of X .

Exercise 16 Show that $\chi \circ \phi$ is a submap if both ϕ and χ are submaps.

1.4.3.10 Set_{sub}

The class of sets with submaps as morphisms defines another category whose objects are sets. We shall denote it Set_{sub} .

1.4.3.11 Set_{fin}

More generally, we shall say that $\phi: X \multimap Y$ is a *finitely-valued map*, if

$$|\phi(x)| < \infty \quad (x \in X). \quad (1.62)$$

Exercise 17 Show that $\chi \circ \phi$ is finitely-valued if both ϕ and χ are finitely-valued.

In particular, sets with finitely-valued maps as morphisms form a category. We shall denote it Set_{fin} .

1.4.3.12 $\text{Set}_{\text{count}}$

Another possibility is to consider *countably-valued maps* as morphisms,

$$\phi(x) \text{ countable for all } x \in X. \quad (1.63)$$

Let us denote the corresponding category by $\text{Set}_{\text{count}}$.

1.4.3.13

The above categories form an increasing chain of unital subcategories of the category of sets and multimaps

$$\text{Set} \subseteq \text{Set}_{\text{sub}} \subseteq \text{Set}_{\text{fin}} \subseteq \text{Set}_{\text{count}} \subseteq \text{Set}_{\text{mult}}.$$

Note that they share the same class of objects. They differ only in their morphisms.

1.4.3.14 Composition of binary relations

A different approach to defining morphisms from a set X to a set Y is to consider binary relations $R \subseteq X \times Y$. For $R \subseteq X \times Y$ and $S \subseteq Y \times Z$,

$$R \circ S := \{(x, z) \in X \times Z \mid \text{there exists } y \in Y \text{ such that } (x, y) \in R \text{ and } (y, z) \in S\}. \quad (1.64)$$

is a binary relation between elements of X and Z . If we use notation $x \sim_R y$ ("element $x \in X$ is in relation R with element $y \in Y$ ") to express the fact that $(x, y) \in R$, then we can rewrite Definition (1.64) as follows

$$R \circ S := \{(x, z) \in X \times Z \mid \text{there exists } y \in Y \text{ such that } x \sim_R y \text{ and } y \sim_S z\}. \quad (1.65)$$

Exercise 18 Show that composition of binary relations is associative, i.e.,

$$(Q \circ R) \circ S = Q \circ (R \circ S)$$

for any $Q \subseteq W \times X$, $R \subseteq X \times Y$, and $S \subseteq Y \times Z$.

1.4.3.15 The identity relation

For any set X we shall call the binary relation

$$\Delta_X := \{(x, x') \in X \times X \mid x = x'\} \quad (1.66)$$

the *identity* relation on X .

Exercise 19 Show that

$$\Delta_X \circ R = R = R \circ \Delta_Y$$

for any $R \subseteq X \times Y$.

1.4.3.16

Denote the category whose objects are sets and relations $R \subseteq X \times Y$ are morphisms $X \rightarrow Y$ by Set_{rel} .

1.4.4 Discrete categories

1.4.4.1

There are much simpler categories than the categories of sets. The simplest, are perhaps the categories with the *empty* class of morphisms. Such categories are referred to as *discrete*.

1.4.4.2 Discrete unital categories

Every unital category is supposed to have at least the identity morphisms for each object. For this reason, in the context of unital categories *discrete* means: *no morphisms besides the identity morphisms*.

1.4.5 Small categories

1.4.5.1

If the class of objects forms a set, such a category is called a *small* category. In this case, the class of morphisms is a set too. Indeed, it is the union

$$\mathcal{C}_1 = \bigcup_{(a,b) \in \mathcal{C}_0 \times \mathcal{C}_0} \text{Hom}_{\mathcal{C}}(a,b)$$

of the family of $\text{Hom}_{\mathcal{C}}(a,b)$ which is indexed by the Cartesian square of the set of objects.

1.4.5.2

Several fundamentally important structures in Mathematics can be interpreted as small categories. We give here just one yet very important example of such structures: a *preordered* set. Other examples will appear later.

1.4.5.3 Preordered sets

We say that a binary relation \rightarrow on a set X is a *preorder* (the term *quasiorder* is used too), if it is *reflexive*,

$$x \rightarrow x \quad (x \in X), \quad (1.67)$$

and *transitive*

$$\text{if } x \rightarrow y \text{ and } y \rightarrow z, \text{ then } x \rightarrow z \quad (x, y, z \in X). \quad (1.68)$$

Of these two properties transitivity is far more important.

A *preordered set*. i.e., a set equipped with a preorder gives rise to the category whose objects are elements of X , and $\text{Hom}(x, y)$ consists of a single element, if $x \rightarrow y$, and is empty otherwise. Since $\text{Hom}(x, y)$ has at most one element, it does not matter how does one denote it. One may use, for example, symbol \rightarrow or, to indicate its source and target, $x \rightarrow y$.

Note that in the associated category, objects x and y are isomorphic if and only if $x \rightarrow y$ and $y \rightarrow x$.

1.4.5.4

Vice-versa, any small category \mathcal{C} with the property that, for any $a, b \in \mathcal{C}_0$,

$$\text{Hom}_{\mathcal{C}}(x, y) \text{ has at most one element,} \quad (1.69)$$

is obtained this way.

Exercise 20 For a small category that satisfies (1.69), show that

$$x \rightarrow y \quad \text{if} \quad \text{Hom}_{\mathcal{C}}(x, y) \neq \emptyset$$

defines a preorder relation on $X := \mathcal{C}_0$.

1.4.5.5 Partially ordered sets

A *partial order* on a set X is a preorder which is *weakly antisymmetric*

$$\text{if } x \rightarrow y \text{ and } y \rightarrow x, \text{ then } x = y. \quad (1.70)$$

1.4.5.6

Small discrete categories correspond to discrete partially ordered sets, i.e., the sets equipped with the smallest order relation—the *identity* relation:

$$x \rightarrow_{\text{discr}} y \quad \text{if} \quad x = y.$$

1.4.6 Functors

1.4.6.1

A *functor* $F: \mathcal{C} \rightsquigarrow \mathcal{D}$ from a category \mathcal{C} to a category \mathcal{D} consists of two correspondences: between the classes of objects and between the classes of morphisms

$$F_0: \mathcal{C}_0 \rightarrow \mathcal{D}_0 \quad \text{and} \quad F_1: \mathcal{C}_1 \rightarrow \mathcal{D}_1$$

which are compatible with all the elements of the category structure. The latter means that the following diagrams of correspondences

$$\begin{array}{ccc}
 \mathcal{C}_0 & \xrightarrow{F_0} & \mathcal{D}_0 \\
 \uparrow s & & \uparrow s \\
 \mathcal{C}_1 & \xrightarrow{F_1} & \mathcal{D}_1 \\
 \downarrow t & & \downarrow t \\
 \mathcal{C}_0 & \xrightarrow{F_0} & \mathcal{D}_0
 \end{array} \tag{1.71}$$

and

$$\begin{array}{ccc}
 \mathcal{C}_2 & \xrightarrow{F_2} & \mathcal{D}_2 \\
 \downarrow m & & \downarrow m \\
 \mathcal{C}_1 & \xrightarrow{F_1} & \mathcal{D}_1
 \end{array} \tag{1.72}$$

are commutative. Here, F_2 denotes the correspondence induced by F_1 on the classes of composable pairs:

$$F_2: \mathcal{C}_2 \rightarrow \mathcal{D}_2, \quad (\alpha, \beta) \mapsto (F_1(\alpha), F_1(\beta)). \tag{1.73}$$

1.4.6.2 Unital functors

When the corresponding categories are *unital*, i.e., possess identity morphisms, then it is customary to require that a functor $F: \mathcal{C} \rightsquigarrow \mathcal{D}$ is compati-

ble also with the identities. This means that the diagram

$$\begin{array}{ccc}
 \mathcal{C}_0 & \xrightarrow{F_0} & \mathcal{D}_0 \\
 \text{id} \downarrow & & \downarrow \text{id} \\
 \mathcal{C}_1 & \xrightarrow{F_1} & \mathcal{D}_1
 \end{array} \tag{1.74}$$

is supposed to commute. We shall call such functors *unital*.

1.4.6.3

In the interest of keeping notation as transparent as possible it is customary to omit subscript indices and denote the correspondences between the objects, morphisms, composable pairs of morphisms, etc., using the same symbol F .

1.4.6.4

Commutativity of the two squares in diagram (1.71) then can be expressed as

$$s(F(\alpha)) = F(s(\alpha)) \quad \text{and} \quad t(F(\alpha)) = F(t(\alpha)) \quad (\alpha \in \mathcal{C}_1), \tag{1.75}$$

while commutativity of diagram (1.72) expresses the fact that

$$F(\alpha) \circ F(\beta) = F(\alpha \circ \beta) \tag{1.76}$$

for any pair of composable morphisms α and β in \mathcal{C} .

Finally, commutativity of diagram (1.74) means that

$$\text{id}_{F(a)} = F(\text{id}_a) \quad (a \in \mathcal{C}_0). \tag{1.77}$$

1.4.6.5 Contravariant functors

The functors we defined above are also called *covariant* functors. The *contravariant* variety is obtained if one requires instead

$$s(F(\alpha)) = F(t(\alpha)) \quad \text{and} \quad t(F(\alpha)) = F(s(\alpha)) \quad (\alpha \in \mathcal{C}_1), \tag{1.78}$$

and

$$F(\beta) \circ F(\alpha) = F(\alpha \circ \beta) \tag{1.79}$$

for any pair of composable morphisms α and β in \mathcal{C} .

Exercise 21 Express requirements (1.78) and (1.79) with help of diagrams analogous to (1.71) and (1.72).

1.4.6.6 An example: the graph functor

For a multimap $\phi: X \multimap Y$, the set

$$\Gamma_\phi := \{(x, y) \in X \times Y \mid \phi(x) \ni y\} \quad (1.80)$$

will be called the *graph* of ϕ . It can be naturally identified with the set

$$\bigcup_{x \in X} \{x\} \times \phi(x).$$

Exercise 22 Verify that $\Gamma_{\text{id}_X} = \Delta_X$ and, for any $\phi: X \multimap Y$ and $\chi: Y \multimap Z$, one has

$$\Gamma_{\chi \circ \phi} = \Gamma_\phi \circ \Gamma_\chi. \quad (1.81)$$

Thus, the double correspondence

$$X \mapsto X, \quad \phi \mapsto \Gamma_\phi \quad (X \in \mathbf{Ob}_{\mathbf{Set}}, \phi \in \mathbf{Arr}_{\mathbf{Set}_{\text{mult}}}), \quad (1.82)$$

defines a contravariant functor $\Gamma: \mathbf{Set}_{\text{mult}} \rightsquigarrow \mathbf{Set}_{\text{rel}}$. When ϕ satisfies condition (1.55), Γ_ϕ coincides with the graph of the corresponding map $f: X \rightarrow Y$.

1.4.6.7

Note that the correspondence

$$\mathbf{Hom}_{\mathbf{Set}_{\text{mult}}}(X, Y) \longrightarrow \mathcal{P}(X \times Y), \quad \phi \longmapsto \Gamma_\phi,$$

is bijective: for any $R \subseteq X \times Y$, one has $R = \Gamma_{\phi_R}$ where $\phi_R: X \multimap Y$ is the multimap

$$\phi_R(x) := \{y \in Y \mid (x, y) \in R\}.$$

1.4.6.8

Functors very often encode natural constructions in Mathematics. We have already encountered a few functors in Section 1.1.2 of the Introduction, all being functors $\mathbf{Set} \rightsquigarrow \mathbf{Set}$, the first and the third being covariant, the second and the fourth being contravariant.

1.4.6.9 The canonical inclusion functors

Given a subcategory \mathcal{C}' of a category \mathcal{C} , the natural inclusion correspondences $\iota_0: \mathcal{C}'_0 \rightarrow \mathcal{C}_0$ and $\iota_1: \mathcal{C}'_1 \rightarrow \mathcal{C}_1$ define the *inclusion* functor $\iota: \mathcal{C}' \rightsquigarrow \mathcal{C}$.

1.4.6.10 The category of small categories

The category whose objects are small categories and morphisms are *covariant* functors between small categories is itself a category. It is denoted Cat and is called the category of (small nonunital) categories.

1.4.6.11 The category of small unital categories

If we consider only unital small categories and unital functors, then we obtain the category of small unital categories. We shall denote it here Cat_1 . The reader should be warned that since categories are usually assumed to possess identity morphisms, the category of small unital categories is often denoted Cat .

1.4.6.12 The category of sets viewed as a subcategory of the category of small categories

Let us identify sets X with small discrete categories \mathcal{X} ,

$$\mathcal{X}_0 = X, \quad \mathcal{X}_1 = \emptyset.$$

Any map $f: X \rightarrow Y$ defines a functor $F: \mathcal{X} \rightsquigarrow \mathcal{Y}$,

$$F_0 = f, \quad F_1 = \text{id}_\emptyset,$$

and every functor $F: \mathcal{X} \rightsquigarrow \mathcal{Y}$ is necessarily of this form since id_\emptyset is the only map from \emptyset to \emptyset .

In particular, the category of sets can be viewed as a full subcategory of the category of small categories.

1.4.6.13 Set viewed as a subactory of Cat

In the unital case, we associate with any set X the category \mathcal{X}' ,

$$\mathcal{X}'_0 = X, \quad \mathcal{X}'_1 = X$$

with all the structural correspondences being id_X (note that $\mathcal{X}'_2 = \{(x, x) \mid x \in X\}$ is here naturally identified with set X).

Any map $f: X \rightarrow Y$ defines a functor $F: \mathcal{X}' \rightsquigarrow \mathcal{Y}'$,

$$F_0 = f, \quad F_1 = f, \tag{1.83}$$

Exercise 23 Show that any unital functor $F: \mathcal{X}' \rightarrow \mathcal{Y}'$ is of the form (1.83).

It follows that Set , the unital category of sets, is a full subcategory of Cat , the category of small unital categories.

1.4.6.14

Since functors between unital categories do not necessarily respect the identity morphisms (an example will be given below), Cat_1 is a subcategory of Cat yet not a full subcategory.

1.4.6.15 Natural transformations of functors

Given two (covariant) functors F and G from a category \mathcal{C} to a category \mathcal{D} , a natural transformation between them, denoted $\phi: F \Rightarrow G$, consists of a single correspondence $\phi: \mathcal{C}_0 \rightarrow \mathcal{D}_1$ which is compatible with all the present structures. The latter means that

$$\phi(a) \in \text{Hom}_{\mathcal{D}}(F(a), G(a)) \quad (a \in \mathcal{C}_0), \quad (1.84)$$

and, for any morphism $\alpha \in \text{Hom}_{\mathcal{C}}(a, b)$, the following square commutes

$$\begin{array}{ccc} F(a) & \xrightarrow{\phi(a)} & G(a) \\ F(\alpha) \downarrow & & \downarrow G(\alpha) \\ F(b) & \xrightarrow{\phi(b)} & G(b) \end{array} \quad (1.85)$$

1.4.6.16

In the language of correspondences, conditions (1.84) translates into commutativity of the following diagram

$$\begin{array}{ccc} & & \mathcal{D}_0 \\ & \nearrow F_0 & \uparrow s \\ \mathcal{C}_0 & \xrightarrow{\phi} & \mathcal{D}_1 \\ & \searrow G_0 & \downarrow t \\ & & \mathcal{D}_0 \end{array} \quad (1.86)$$

while conditions (1.85) expresses commutativity of the diagram

$$\begin{array}{ccc}
 \mathcal{C}_1 & \xrightarrow{(\phi \circ t, F_1)} & \mathcal{D}_2 \\
 (G_1, \phi \circ s) \downarrow & & \downarrow m \\
 \mathcal{D}_2 & \xrightarrow{m} & \mathcal{D}_1
 \end{array} \tag{1.87}$$

Exercise 24 Formulate the definition of a natural transformation of contravariant functors analogous to (1.84)-(1.85).

Exercise 25 Formulate the definition of a natural transformation of contravariant functors analogous to diagrams (1.86)-(1.87).

1.4.6.17

We have already encountered a natural transformation of contravariant functors $\chi: \mathcal{P}(\) \Rightarrow 2^{(\)}$ in Section 1.1.2.6.

1.4.6.18

Many properties normally expressed as identities involving objects, morphisms, sets, maps, elements of various sets, etc, can be often expressed as commutativity of certain diagrams. This leads to proliferation of what some call ‘diagrammatic thinking’ in modern Mathematics. Employing diagrams often can significantly clarify the picture.

On some occasions information conveyed by diagrams may be more difficult to understand than the same information expressed differently. I would say that it is probably easier to understand the meaning of conditions (1.85) than the meaning of the commutativity of diagram (1.87). That is probably due to the fact that the conditions (1.85) are themselves expressed in terms of commutativity of some easy-to-understand diagrams.

1.4.7 The opposite category

1.4.7.1

Note that if one retains the classes of objects and arrows, \mathcal{C}_0 and \mathcal{C}_1 , but exchanges the source and the target correspondences, $s: \mathcal{C}_1 \rightarrow \mathcal{C}_0$ and $t: \mathcal{C}_1 \rightarrow \mathcal{C}_0$, then one obtains a category again. This is the *opposite* category \mathcal{C}^{op} .

1.4.7.2

More precisely,

$$\mathcal{C}_0^{\text{op}} = \mathcal{C}_0, \quad \mathcal{C}_1^{\text{op}} = \mathcal{C}_1, \quad s^{\text{op}} = t, \quad \text{and} \quad t^{\text{op}} = s. \quad (1.88)$$

If an object a of \mathcal{C} is considered as an object of \mathcal{C}^{op} , then it should be denoted a^{op} . Similarly for morphisms: if $\alpha: a \rightarrow b$ is a morphism in \mathcal{C} , then α considered as a morphism of the opposite category is a morphism $b^{\text{op}} \rightarrow a^{\text{op}}$ and it should be denoted α^{op} .

1.4.7.3

The correspondences

$$a \mapsto a^{\text{op}} \quad \text{and} \quad \alpha \mapsto \alpha^{\text{op}} \quad (a \in \mathcal{C}_0; \alpha \in \mathcal{C}_1),$$

define a *contravariant* functor

$$(\)_{\mathcal{C}}^{\text{op}}: \mathcal{C} \rightsquigarrow \mathcal{C}^{\text{op}}.$$

1.4.7.4

Note that

$$(\)_{\mathcal{C}}^{\text{op}} \circ (\)_{\mathcal{C}^{\text{op}}}^{\text{op}} = \text{id}_{\mathcal{C}^{\text{op}}} \quad \text{and} \quad (\)_{\mathcal{C}^{\text{op}}}^{\text{op}} \circ (\)_{\mathcal{C}}^{\text{op}} = \text{id}_{\mathcal{C}}.$$

1.4.7.5 An example: a partially ordered set

If \mathcal{C} is the category that corresponds to a partially ordered set (X, \preceq) , cf. Section 1.4.5.6, then \mathcal{C}^{op} corresponds to set X equipped with the *reverse* order, \preceq^{rev} .

1.4.7.6

One of the uses of the concept of the opposite category is that it allows to consider any *contravariant* functor $F: \mathcal{C} \rightsquigarrow \mathcal{D}$ as a *covariant* functor either $\mathcal{C} \rightsquigarrow \mathcal{D}^{\text{op}}$ or $\mathcal{C}^{\text{op}} \rightsquigarrow \mathcal{D}$. Formally speaking, this is done by composing F with $(\)_{\mathcal{D}}^{\text{op}}$ or $(\)_{\mathcal{C}^{\text{op}}}^{\text{op}}$,

$$(\)_{\mathcal{D}}^{\text{op}} \circ F: \mathcal{C} \rightarrow \mathcal{D}^{\text{op}} \quad \text{or} \quad F \circ (\)_{\mathcal{C}^{\text{op}}}^{\text{op}}: \mathcal{C}^{\text{op}} \rightsquigarrow \mathcal{D}.$$

1.4.7.7

Any functor $F: \mathcal{C} \rightsquigarrow \mathcal{D}$, induces also a functor from \mathcal{C}^{op} to \mathcal{D}^{op}

$$F^{\text{op}} := ()_{\mathcal{D}}^{\text{op}} \circ F \circ ()_{\mathcal{C}^{\text{op}}}^{\text{op}}. \quad (1.89)$$

Note that F^{op} is covariant (respectively, contravariant) when F is covariant (respectively, contravariant).

1.4.7.8

Assigning to any category \mathcal{C} its opposite category \mathcal{C}^{op} is *natural* in \mathcal{C} , so one can expect that it gives rise to a functor on the category of (small) categories. This is so indeed, the correspondences

$$\mathcal{C} \mapsto \mathcal{C}^{\text{op}} \quad \text{and} \quad F \mapsto F^{\text{op}} \quad (\mathcal{C} \in \text{Cat}_0; F \in \text{Cat}_1), \quad (1.90)$$

defined by (1.88) and (1.89), yield a functor $()^{\text{op}}: \text{Cat} \rightsquigarrow \text{Cat}$.

Exercise 26 *Is functor (1.90) covariant or contravariant?*

1.4.7.9 $\text{Set}_{\text{mult}} \simeq (\text{Set}_{\text{rel}})^{\text{op}}$

The graph functor, $\Gamma: \text{Set}_{\text{mult}} \rightsquigarrow \text{Set}_{\text{rel}}$, which was defined in Section 1.4.6.6, identifies the category of sets with multimaps as morphisms with the category opposite to the category of sets with binary relations as morphisms. In other words, Set_{mult} is isomorphic to $(\text{Set}_{\text{rel}})^{\text{op}}$.

Isomorphisms between categories are, generally speaking, a rare occurrence.

1.4.7.10 Importance of the opposite category concept

Any diagram in a category \mathcal{C} can be interpreted as the same diagram—but with the *direction of all arrows reversed*—in the opposite category.

An immediate corollary of this simple observation yields the following *Duality Principle*:

For any categorical concept or construction involving one or more diagrams, there is a *dual* concept or construction.

1.4.8 Categories of arrows

1.4.8.1

For any category there are several naturally associated categories whose objects are morphisms. We shall mention here three.

1.4.8.2 The category of arrows

For a category \mathcal{C} , let \mathcal{C}^\rightarrow be the category whose objects are morphisms of \mathcal{C} ,

$$(\mathcal{C}^\rightarrow)_0 := \mathcal{C}_1, \quad (1.91)$$

and morphisms $\phi: \alpha \rightarrow \beta$ are pairs of morphisms $\phi = (\phi_s, \phi_t)$ in \mathcal{C} ,

$$\phi_s: s(\alpha) \rightarrow s(\beta), \quad \phi_t: t(\alpha) \rightarrow t(\beta), \quad (1.92)$$

such that the following diagram commutes

$$\begin{array}{ccc} \bullet & \xrightarrow{\phi_s} & \bullet \\ \alpha \downarrow & & \downarrow \beta \\ \bullet & \xrightarrow{\phi_t} & \bullet \end{array} \quad (1.93)$$

1.4.8.3

Category of arrows \mathcal{C}^\rightarrow is sometimes also denoted $\text{Arr } \mathcal{C}$. One should be advised however, that $\text{Arr } \mathcal{C}$ may also be used to denote the *class* of morphisms in \mathcal{C} .

1.4.8.4 Two comma categories

For any object a in a category \mathcal{C} , one can consider two categories: one, $\mathcal{C}^{a \rightarrow}$, whose objects are morphisms in \mathcal{C} with source a ,

$$(\mathcal{C}^{a \rightarrow})_0 := \{\alpha \in \mathcal{C}_1 \mid s(\alpha) = a\}, \quad (1.94)$$

and another one, $\mathcal{C}^{\rightarrow a}$, whose objects are morphisms with target a ,

$$(\mathcal{C}^{\rightarrow a})_0 := \{\alpha \in \mathcal{C}_1 \mid t(\alpha) = a\}. \quad (1.95)$$

1.4.8.5

Morphisms $\phi: \alpha \rightarrow \beta$ in $\mathcal{C}^{a \rightarrow}$ are morphisms $\phi: t(\alpha) \rightarrow t(\beta)$ such that the following diagram commutes

$$\begin{array}{ccc} & & \bullet \\ & \nearrow \alpha & \\ a & & \\ & \searrow \beta & \\ & & \bullet \end{array} \quad \begin{array}{c} \vdots \\ \phi \\ \vdots \end{array} \quad (1.96)$$

1.4.8.6

Morphisms $\phi: \alpha \rightarrow \beta$ in $\mathcal{C}^{\rightarrow a}$ are morphisms $\phi: s(\alpha) \rightarrow s(\beta)$ such that the following diagram commutes

$$\begin{array}{ccc} \bullet & & \\ & \searrow \alpha & \\ & & a \\ & \nearrow \beta & \\ \bullet & & \end{array} \quad \begin{array}{c} \vdots \\ \phi \\ \vdots \end{array} \quad (1.97)$$

1.4.9 Categories of diagrams

1.4.9.1

(Covariant) functors from a *small* category Γ to an arbitrary category \mathcal{C} form a category, denoted \mathcal{C}^Γ , with morphisms $\phi: F \rightarrow G$ being natural transformations of functors.

1.4.9.2 Diagrams as functors

Such functors are often called *diagrams in \mathcal{C}* and the reason will become clear when we look at a series of simple examples.

1.4.9.3 \mathcal{C}

Consider the category with a single object, \mathbf{o} , with empty class of morphisms. Denote this category by $\mathbf{1}$. Functors from $\mathbf{1}$ to \mathcal{C} correspond to single objects in \mathcal{C} , and $\mathcal{C}^{\mathbf{1}}$ becomes naturally identified with category \mathcal{C} itself.

1.4.9.4 \mathcal{C}^\rightarrow

Consider the category with two objects, 0 and 1, and a single morphism

$$0 \rightarrow 1.$$

Denote this category by **2**. Functors from **2** to \mathcal{C} correspond to single morphisms in \mathcal{C} , and \mathcal{C}^2 becomes naturally identified with the category of arrows, \mathcal{C}^\rightarrow .

1.4.9.5 The category of composable pairs of arrows

Consider the category with three objects, 0, 1 and 2, and just three morphisms, the following two

$$0 \rightarrow 1 \rightarrow 2,$$

and their composition. Denote this category by **3**. Functors from **3** to \mathcal{C} correspond to composable pairs of morphisms in \mathcal{C} , and \mathcal{C}^3 becomes naturally identified with the category of composable pairs of arrows in \mathcal{C} .

Exercise 27 *The category of composable pairs of arrows in \mathcal{C} has class \mathcal{C}_2 as its class of objects. Knowing that morphisms $\phi: (\alpha_0, \alpha_1) \rightarrow (\beta_0, \beta_1)$ are defined in a natural manner, give the definition of morphisms.*

1.4.9.6

Categories **1**, **2** and **3** correspond to the linearly ordered sets $\{0\}$, $\{0, 1\}$, $\{0, 1, 2\}$. Let **n** be the category with n objects

$$0, 1, \dots, n-1$$

which corresponds to the linearly ordered set $\{0, \dots, n-1\}$.

Exercise 28 *Find the number of morphisms in **n**.*

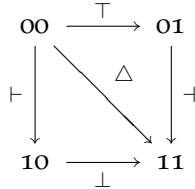
Exercise 29 *Provide a description of \mathcal{C}^n which generalizes to arbitrary n the descriptions given above for $n = 1, 2, 3$.*

1.4.9.7 The category of commuting squares

Consider the category with four objects

00 , 01 , 10 , and 11 ,

and just five morphisms



Denote this category by \square . Objects of \mathcal{C}^\square are commuting squares in \mathcal{C} .

Exercise 30 Describe morphisms in \mathcal{C}^\square .

1.4.9.8 The category of families of objects

Let \mathbf{I} be the category with a set I as its class of objects and empty class of morphisms. Objects of $\mathcal{C}^{\mathbf{I}}$ are families $(a_i)_{i \in I}$ of objects of category \mathcal{C} indexed by set I .

Exercise 31 Describe morphisms $\phi: (a_i)_{i \in I} \rightarrow (b_i)_{i \in I}$.

Chapter 2

Sets equipped with a family of subsets

2.1 Two categories of pairs

2.1.1

2.1.1.1

Pairs (X, \mathcal{A}) , where X is a set and $\mathcal{A} \subseteq \mathcal{P}(X)$ is a family of subsets of X , form a category in two natural ways. In both cases, pairs (X, \mathcal{A}) provide the objects.

The difference between these two categories is their morphisms: in the first case we consider the power-set functor $\mathcal{P}(\)$ as a covariant functor, in the second—as a contravariant functor.

2.1.1.2

A map $f: X \rightarrow Y$ induces the pair of maps

$$f_*: \mathcal{P}(X) \longrightarrow \mathcal{P}(Y), \quad f^*: \mathcal{P}(Y) \longrightarrow \mathcal{P}(X).$$

Each of these maps, in turn, induces the corresponding pair of maps between the iterated power sets

$$(f_*)_*: \mathcal{P}(\mathcal{P}(X)) \longrightarrow \mathcal{P}(\mathcal{P}(Y)), \quad (f_*)^*: \mathcal{P}(\mathcal{P}(Y)) \longrightarrow \mathcal{P}(\mathcal{P}(X)),$$

and

$$(f^*)_*: \mathcal{P}(\mathcal{P}(Y)) \longrightarrow \mathcal{P}(\mathcal{P}(X)), \quad (f^*)^*: \mathcal{P}(\mathcal{P}(X)) \longrightarrow \mathcal{P}(\mathcal{P}(Y)).$$

2.1.1.3 Notation

In order to simplify notation we shall drop the parentheses:

$$f_{**}\mathcal{A} := (f_*)_*(\mathcal{A}), \quad f^{**}\mathcal{A} := (f^*)^*(\mathcal{A}),$$

and

$$f_*^*\mathcal{B} := (f_*)^*(\mathcal{B}), \quad f_*^*\mathcal{B} := (f^*)_*(\mathcal{B}).$$

We shall, however, observe the correct placement of the asterisk sub- and superscripts, since

$$f_*^* \neq f_*^*.$$

2.1.1.4 The family of images

Exercise 32 Show that

$$f_{**}\mathcal{A} = \{f(A) \mid A \in \mathcal{A}\} = \{F \subseteq Y \mid F = f(A) \text{ for some } A \in \mathcal{A}\}. \quad (2.1)$$

We shall refer to $(f_*)_*(\mathcal{A})$ as the *family of f -images* for \mathcal{A} .

2.1.1.5 The inverse image of a family

Exercise 33 Show that

$$f_*^*\mathcal{B} = \{E \subseteq X \mid f(E) \in \mathcal{B}\} \quad (2.2)$$

We shall refer to $f_*^*\mathcal{B}$ as the *inverse image* of \mathcal{B} .

2.1.1.6 The family of preimages

Exercise 34 Show that

$$f_*^*\mathcal{B} = \{f^{-1}(B) \mid B \in \mathcal{B}\} = \{E \subseteq X \mid E = f^{-1}(B) \text{ for some } B \in \mathcal{B}\}. \quad (2.3)$$

We shall refer to $(f^*)_*(\mathcal{B})$ as the *family of f -preimages* for \mathcal{B} .

2.1.1.7 The direct image of a family

Exercise 35 Show that

$$f^{**}\mathcal{A} = \{F \subseteq Y \mid f^{-1}(F) \in \mathcal{A}\} \quad (2.4)$$

We shall refer to $(f^*)^*(\mathcal{A})$ as the *direct image* of \mathcal{A} .

2.1.1.8

All three pairs in 2.1.1.2 form Galois connections between the corresponding power sets ordered by the relation \subseteq . The iterated power sets are, however, equipped also with preorders \dashv and ε , and we have further pairs of Galois connections for two out of four mappings and for each of these preorders.

Exercise 36 *Show that*

$$f_{**}\mathcal{A} \varepsilon \mathcal{B} \quad \text{if and only if} \quad \mathcal{A} \varepsilon f_*^* \mathcal{B}$$

and

$$f_{**}\mathcal{A} \dashv \mathcal{B} \quad \text{if and only if} \quad \mathcal{A} \dashv f_*^* \mathcal{B}.$$

2.1.1.9

There are further relations involving the pairs of mappings that “go in the same direction”.

Exercise 37 *Show that*

$$f_{**}\mathcal{A} \varepsilon f^{**}\mathcal{A} \quad \text{and} \quad f_*^* \mathcal{B} \dashv f_*^* \mathcal{B}.$$

2.1.1.10

2.1.1.11 The first category of pairs (X, \mathcal{A})

Morphisms $(X, \mathcal{A}) \rightarrow (Y, \mathcal{B})$ are maps $f: X \rightarrow Y$ such that

$$f(A) \in \mathcal{B} \quad \text{for any} \quad A \in \mathcal{A}, \quad (2.5)$$

i.e., family $f_{**}\mathcal{A}$ is contained in family \mathcal{B} or, equivalently, family \mathcal{A} is contained in family $f_*^* \mathcal{B}$.

2.1.1.12 The second category of pairs (X, \mathcal{A})

Morphisms $(X, \mathcal{A}) \rightarrow (Y, \mathcal{B})$ are maps $f: X \rightarrow Y$ such that

$$f^{-1}(B) \in \mathcal{A} \quad \text{for any} \quad B \in \mathcal{B}, \quad (2.6)$$

i.e., family $f_*^* \mathcal{B}$ is contained in family \mathcal{A} or, equivalently, family \mathcal{B} is contained in family $f^{**}\mathcal{A}$.

2.1.1.13

One could profitably refer to the first as the *covariant* category of pairs (X, \mathcal{A}) , and to the second—as the *contravariant* category of pairs (X, \mathcal{A}) . Be forewarned however that the words ‘covariant’ and ‘contravariant’ are here used strictly as names that allow us to clearly indicate which of the two categories of pairs we mean. As *concepts*, ‘covariant’ and ‘contravariant’ apply to functors, not categories.

2.1.2 A variant: *rich* categories of pairs

2.1.2.1 Binary relations acting on power sets

Any binary relation $R \subseteq X \times Y$ naturally induces two order preserving maps,

$$\cdot R : \mathcal{P}(X) \rightarrow \mathcal{P}(Y) \quad \text{and} \quad R \cdot : \mathcal{P}(Y) \rightarrow \mathcal{P}(X), \quad (2.7)$$

where

$$A \mapsto A \cdot R := \{y \in Y \mid \text{there exists } x \in A \text{ such that } x \sim_R y\} \quad (2.8)$$

and

$$B \mapsto R \cdot B := \{x \in X \mid \text{there exists } y \in B \text{ such that } x \sim_R y\}. \quad (2.9)$$

To simplify notation we shall often omit the subscript if it is clear within which set the complement of a subset is formed.

2.1.2.2

If we identify X with $1 \times X$ and Y with $Y \times 1$,

$$x \leftrightarrow (1, x), \quad y \leftrightarrow (y, 1), \quad (x \in X, y \in Y), \quad (2.10)$$

then subsets $A \subseteq X$ and $B \subseteq Y$ become binary relations between elements of sets 1 and X , and elements of sets Y and 1 , respectively. Under these identifications, formulae (2.8) and (2.9) express composition of the corresponding relations, cf. Section 1.4.3.14

2.1.2.3

By conjugating mappings $\cdot R$ and $R \circ$ with the complement operations

$$(\cdot)_X^c: \mathcal{P}(X) \rightarrow \mathcal{P}(X), \quad A \mapsto A_X^c := X \setminus A,$$

and

$$(\cdot)_Y^c: \mathcal{P}(Y) \rightarrow \mathcal{P}(Y), \quad B \mapsto B_Y^c := Y \setminus B,$$

we obtain two additional order preserving maps

$$\diamond R: \mathcal{P}(X) \rightarrow \mathcal{P}(Y) \quad \text{and} \quad R \diamond: \mathcal{P}(Y) \rightarrow \mathcal{P}(X), \quad (2.11)$$

where

$$\begin{aligned} A \mapsto A \diamond R &:= (A_X^c \cdot R)_Y^c = Y \setminus ((X \setminus A) \cdot R) \\ &= \{y \in Y \mid x \approx_R y \text{ whenever } x \notin A\} \end{aligned} \quad (2.12)$$

and

$$\begin{aligned} B \mapsto R \diamond B &:= (R \cdot B_Y^c)_X^c = X \setminus (R \circ (Y \setminus B)) \\ &= \{x \in X \mid x \approx_R y \text{ whenever } y \notin B\}. \end{aligned} \quad (2.13)$$

Exercise 38 Let $A \subseteq X$ and $B \subseteq Y$. Show that

$$A \cdot R \subseteq B \quad \text{if and only if} \quad A \subseteq R \diamond B \quad (2.14)$$

and

$$R \cdot B \subseteq A \quad \text{if and only if} \quad B \subseteq A \diamond R. \quad (2.15)$$

2.1.2.4

Equivalence (2.14) expresses the fact that the pair of mappings $(\cdot R, R \diamond)$ forms a *Galois connection* between partially ordered sets $(\mathcal{P}(X), \subseteq)$ and $(\mathcal{P}(Y), \subseteq)$. One can show that any Galois connection between the power sets is of this form for a unique binary relation $R \subseteq X \times Y$ (cf. *Notes on Partially Ordered Sets*).

2.1.2.5

Dually, equivalence (2.15) expresses the fact that the pair of mappings $(R \diamond, \cdot R)$ forms a *Galois connection* between partially ordered sets $(\mathcal{P}(Y), \subseteq)$ and $(\mathcal{P}(X), \subseteq)$.

Exercise 39 Show that

$$A \cdot \Gamma_f = f(A) \quad \text{and} \quad \Gamma_f \cdot B = \Gamma_f \diamond B = f^{-1}(B) \quad (A \subseteq X, B \subseteq Y),$$

and

$$A \diamond \Gamma_f = (f(A^c))^c = Y \setminus f(X \setminus A), \quad (A \subseteq X),$$

where Γ_f is the graph of a map $f: X \rightarrow Y$.

2.1.2.6

Note that $\Gamma_f \diamond = \Gamma_f \circ$ while $\diamond \Gamma_f \neq \cdot \Gamma_f$, in general. This is due to the fact that while the preimage-of-a-map, $B \mapsto f^{-1}(B)$, commutes with the operation of taking the complement, the image-of-a-map, $A \mapsto f(A)$, does not.

Exercise 40 Let $\mathcal{A} \subseteq \mathcal{P}(X)$ be a family of subsets of X and $\mathcal{B} \subseteq \mathcal{P}(Y)$ be a family of subsets of Y . Show that $f\mathcal{A}$ coincides with the image of \mathcal{A} under mapping $\cdot \Gamma_f$ while $f_\bullet \mathcal{A}$ coincides with the preimage of \mathcal{A} under mapping $\Gamma_f \circ$,

$$f\mathcal{A} = (\cdot \Gamma_f)(\mathcal{A}) \quad \text{and} \quad f_\bullet \mathcal{A} = (\Gamma_f \cdot)^{-1}(\mathcal{A}).$$

Dually, show that $f^{-1}\mathcal{B}$ coincides with the preimage of \mathcal{B} under mapping $\cdot \Gamma_f$ while $f^\bullet \mathcal{B}$ coincides with the preimage of \mathcal{B} under mapping $\Gamma_f \circ$,

$$f^{-1}\mathcal{B} = (\Gamma_f \circ)(\mathcal{B}) \quad \text{and} \quad f^\bullet \mathcal{B} = (\circ \Gamma_f)^{-1}(\mathcal{B}).$$

2.1.2.7

In the “rich” variant of the *first* category of pairs, morphisms $(X, \mathcal{A}) \rightarrow (Y, \mathcal{B})$ are relations $R \subseteq X \times Y$ such that

$$A \cdot R \in \mathcal{B} \quad \text{for any} \quad A \in \mathcal{A}. \quad (2.16)$$

2.1.2.8

Like in the case of the category of sets, the *graph of the map* defines a functor that identifies the first category of pairs with a subcategory of the category opposite to the “rich” first category of pairs (cf. Sections 1.4.6.6 and 1.4.7.9).

2.1.2.9

In the “rich” variant of the *second* category of pairs, morphisms $(X, \mathcal{A}) \rightarrow (Y, \mathcal{B})$ are relations $R \subseteq X \times Y$ such that

$$R \cdot G \in \mathcal{A} \quad \text{for any} \quad B \in \mathcal{B}. \quad (2.17)$$

The *graph of the map* defines a functor that identifies the second category of pairs with a subcategory of the category opposite to the “rich” second category of pairs.

2.2 Topological spaces

2.2.1 Topologies

2.2.1.1

A family $\mathcal{T} \subseteq \mathcal{P}(X)$ is called a *topology* on a set X if it is closed under arbitrary unions and finite intersections, which means that, for any subfamily $\mathcal{V} \subseteq \mathcal{T}$, one has

$$\bigcup \mathcal{V} \in \mathcal{T} \quad (2.18)$$

and, for any *finite* subfamily $\mathcal{V} \subseteq \mathcal{T}$,

$$\bigcap \mathcal{V} \in \mathcal{T}. \quad (2.19)$$

It is also assumed that the smallest and the largest elements of $\mathcal{P}(X)$ belong to \mathcal{T} :

$$\emptyset \in \mathcal{T} \quad \text{and} \quad X \in \mathcal{T}. \quad (2.20)$$

Note, however, that \emptyset is the *union* of the empty subfamily $\emptyset \subseteq \mathcal{T}$ whereas X is the *intersection* of the empty subfamily, so, formally speaking, the two conditions of (2.20) follow from conditions (2.18)–(2.19).

2.2.1.2 The set of topologies on a set

The set $\text{Top}(X)$ of topologies on a set X is a subset of the set of all families of subsets of X , i.e., of $\mathcal{P}(\mathcal{P}(X))$. In particular, it is ordered by inclusion. It possesses the smallest element

$$\mathcal{T}^{\text{triv}} := \{\emptyset, X\} \quad (2.21)$$

which is called the *trivial* topology, and the largest element

$$\mathcal{T}^{\text{discr}} := \mathcal{P}(X) \quad (2.22)$$

which is called the *discrete* topology.

Exercise 41 Show that the intersection of any family of topologies \mathcal{T} on X ,

$$\bigcap_{\mathcal{T} \in \mathcal{T}} \mathcal{T} = \bigcap_{\mathcal{T} \in \mathcal{T}} \mathcal{T}, \quad (2.23)$$

is a topology on X .

2.2.1.3 The topology generated by a family of subsets

For any family of subsets $\mathcal{A} \subseteq \mathcal{P}(X)$ of a set X , the intersection of the family of all topologies \mathcal{T} containing \mathcal{A} ,

$$\mathcal{T}_{\mathcal{A}} := \bigcap_{\substack{\mathcal{T} \in \text{Top}(X) \\ \mathcal{T} \supseteq \mathcal{A}}} \mathcal{T},$$

is the smallest topology that contains \mathcal{A} . We shall call it the topology *generated* by \mathcal{A} .

2.2.1.4

Since (2.23) is the largest family of subsets of X , which is contained in every member \mathcal{T} of the family, it follows from Exercise 41 that any subset \mathcal{T} of partially ordered set $\text{Top}(X)$ has infimum, and that this infimum coincides with the infimum of \mathcal{T} when viewed as a subset of $\mathcal{P}(\mathcal{P}(X))$:

$$\inf_{\text{Top}(X)} \mathcal{T} = \bigcap \mathcal{T} = \inf_{\mathcal{P}(\mathcal{P}(X))} \mathcal{T}. \quad (2.24)$$

2.2.1.5

Recall that in any partially ordered set (S, \preceq) , if s is the infimum of the set $U(E)$ of upper bounds of a set $E \subseteq S$, then $s \in U(E)$ which means that

$$\inf U(E) = \min U(E),$$

and $\min U(E)$ is, by definition, $\sup E$. In particular, if every subset E has infimum in S , it has also supremum in S .

Applying this to $S = \text{Top}(X)$, we see that any family of topologies \mathcal{T} on X has the supremum. Unlike the corresponding infima, the supremum of \mathcal{T} in $\text{Top}(X)$ generally does not coincide with the supremum of \mathcal{T} in $\mathcal{P}(\mathcal{P}(X))$ because the union of a family of topologies is only rarely a topology.

Exercise 42 Show that $\sup_{\text{Top}(X)} \mathcal{T}$ is the topology generated by $\sup_{\mathcal{P}(\mathcal{P}(X))} \mathcal{T}$.

2.2.2 Topological spaces

2.2.2.1

Pairs (X, \mathcal{T}_X) , where \mathcal{T}_X is a topology on a set X , are called *topological spaces*. Topological spaces naturally form a subcategory of the category of pairs, and we have two possibilities: to consider topological spaces as a full subcategory of the *covariant* category of pairs, cf. 2.1.1.11, or of the *contravariant* category of pairs, cf. 2.1.1.12

2.2.2.2 Open maps

In the first case, morphisms $(X, \mathcal{T}_X) \rightarrow (Y, \mathcal{T}_Y)$ are called *open maps*.

2.2.2.3 Continuous maps

In the second case, morphisms $(X, \mathcal{T}_X) \rightarrow (Y, \mathcal{T}_Y)$ are called *continuous maps*.

2.2.2.4 The category of topological spaces

Since continuous maps are considered to be far more important than open maps, the established practice is to apply the name *the category of topological spaces* to the category whose morphisms are continuous maps. This category is usually denoted Top .

Exercise 43 Let $f: X \rightarrow Y$ be any mapping between sets, \mathcal{T} be a topology on X , and \mathcal{T}' be a topology on Y . Show that $f_{\bullet} \mathcal{T}$ is a topology on Y and $f^{-1} \mathcal{T}'$ is a topology on X .

Show that f is a continuous mapping $(X, \mathcal{T}) \rightarrow (Y, \mathcal{T}')$ if and only if

$$\mathcal{T}' \subseteq f_{\bullet} \mathcal{T}$$

if and only if

$$f^{-1} \mathcal{T}' \subseteq \mathcal{T}.$$

2.2.2.5

For obvious reasons, $f_{\bullet}\mathcal{T}$ is said to be the *strongest* topology on Y such that mapping f from (X, \mathcal{T}) is continuous, while $f^{-1}\mathcal{T}'$ is said to be the *weakest* topology on X such that mapping f to (Y, \mathcal{T}') is continuous.

2.2.2.6 The category of sets viewed as a subcategory of the category of topological spaces

Any map between discrete topological spaces is continuous:

$$\text{Hom}_{\text{Top}}\left((X, \mathcal{T}^{\text{discr}}), (Y, \mathcal{T}^{\text{discr}})\right) = \text{Hom}_{\text{Set}}(X, Y).$$

This observation allows us to consider Set as a subcategory of Top .

2.2.3 Measurable spaces

2.2.4 σ -algebras of subsets

2.2.4.1

A family $\mathcal{M} \subseteq \mathcal{P}(X)$ of subsets of a set X is called a σ -algebra if it is closed under *countable* unions and the operation of taking the complement

$$A \mapsto A^c := X \setminus A. \quad (2.25)$$

It is also assumed that $X \in \mathcal{M}$. Note, however, that X is the intersection of the empty family of subsets $\emptyset \subseteq \mathcal{M}$, hence $X \in \mathcal{M}$ in view of \mathcal{M} being closed under intersections of arbitrary countable subfamilies of \mathcal{M} .

2.2.4.2 The set of σ -algebras on a set

The set $\sigma\text{-alg}(X)$ of σ -algebras on a set X is a subset of the set of all families of subsets of X , i.e., of $\mathcal{P}(\mathcal{P}(X))$. In particular, it is ordered by inclusion. It possesses the smallest element

$$\mathcal{M}^{\text{triv}} := \{\emptyset, X\} \quad (2.26)$$

which is called the *trivial* σ -algebra, and the largest element

$$\mathcal{M}^{\text{discr}} := \mathcal{P}(X) \quad (2.27)$$

which will be called the *discrete* σ -algebra.

Exercise 44 Show that the intersection of any family of σ -algebras \mathcal{M} on X ,

$$\bigcap \mathcal{M} = \bigcap_{\mathcal{M} \in \mathcal{M}} \mathcal{M}, \quad (2.28)$$

is a σ -algebra on X .

2.2.4.3 The σ -algebra generated by a family of subsets

For any family of subsets $\mathcal{A} \subseteq \mathcal{P}(X)$ of a set X , the intersection of the family of all σ -algebras \mathcal{M} containing \mathcal{A} ,

$$\mathcal{A}^* := \bigcap_{\substack{\mathcal{M} \in \sigma\text{-alg}(X) \\ \mathcal{M} \supseteq \mathcal{A}}} \mathcal{M},$$

is the smallest σ -algebra on X which contains \mathcal{A} . We shall call it the σ -algebra generated by \mathcal{A} .

2.2.4.4 Measurable spaces

2.2.4.5

Pairs (X, \mathcal{M}) are referred to as *measurable spaces*.

2.2.4.6 Measurable maps

As in the case of topological spaces, we have two choices what to consider to be a morphism $(X, \mathcal{M}) \rightarrow (Y, \mathcal{N})$. And again, we condition (2.6) is the more important one. Maps $f: X \rightarrow Y$ such that

$$\text{for any } B \in \mathcal{N}, \text{ one has } f^{-1}(B) \in \mathcal{M}, \quad (2.29)$$

will be called *measurable*.

Exercise 45 Let $f: X \rightarrow Y$ be any mapping between sets, \mathcal{M} be a σ -algebra on X , and \mathcal{M}' be a σ -algebra on Y . Show that $f_{\bullet}\mathcal{M}$ is a σ -algebra on Y and $f^{-1}\mathcal{M}'$ is a σ -algebra on X .

Show that f is a measurable mapping $(X, \mathcal{M}) \rightarrow (Y, \mathcal{M}')$ if and only if

$$\mathcal{M}' \subseteq f_{\bullet}\mathcal{M}$$

if and only if

$$f^{-1}\mathcal{M}' \subseteq \mathcal{M}.$$

2.2.4.7 The category of measurable spaces

Below, the *category of measurable spaces* will always mean the full subcategory of the *contravariant* category of pairs, cf. 2.1.1.12. We shall denote it *Meas*.

2.2.5 Borel σ -algebra

2.2.5.1 Borel subsets of a topological space

For a topological space (X, \mathcal{T}) , the σ -algebra \mathcal{T}^* generated by the topology is called the *Borel σ -algebra*, and its members—Borel subsets of X .

2.2.5.2 Borel maps between topological spaces

A map $f: X \rightarrow Y$ between topological spaces is called a *Borel map* if it is a morphism of the corresponding Borel measurable spaces

$$(X, (\mathcal{T}_X)^*) \rightarrow (Y, (\mathcal{T}_Y)^*).$$

Any continuous map $f: (X, \mathcal{T}_X) \rightarrow (Y, \mathcal{T}_Y)$ is a Borel map.

2.2.5.3

A map $f: X \rightarrow Y$ from a measurable space to a topological space is said to be *measurable* if it is a morphism $(X, \mathcal{M}) \rightarrow (Y, (\mathcal{T}_Y)^*)$. Thus, we will be also talking of *measurable functions* $f: X \rightarrow \mathbf{R}$, $f: X \rightarrow [0, \infty]$, etc.

Chapter 3

Sets equipped with one or more relations

3.1 Introduction

3.1.1 Relations on a set

3.1.1.1 I -ary relations

Let I be a set. An I -ary relation on a set X is the same as a subset $R \subseteq X^I$.

3.1.1.2 Morphisms

A natural notion of a morphism $(X, R) \rightarrow (Y, S)$ is that it is a map $X \rightarrow Y$ such that the induced map $f_* : X^I \rightarrow Y^I$ sends R to S :

$$f_*(R) \subseteq S. \quad (3.1)$$

Explicitly, this means that if $(x_i)_{i \in I} \in R$, then $(f(x_i))_{i \in I} \in S$.

Exercise 46 *Formulate the notion of a set with two relations, and define the appropriate notion of a morphism.*

3.1.1.3 Notation: binary relations

If $R \subseteq X^2$ is a binary relation on a set X , an alternative notation may be used to denote the fact that $(x, x') \in R$:

$$x \sim_R x'$$

or simply

$$x \sim x',$$

when the relation is clear from the context. Here \sim is a generic symbol for a pair of elements 'being in relation'. In specific situations special symbols may be used. For example, when R is a *partial order* relation, then the symbols \preceq or \leq are generally used.

3.1.1.4

In this notation, a morphism $(X, \sim_X) \rightarrow (Y, \sim_Y)$ is a map $f: X \rightarrow Y$ such that

$$x \sim_X x' \quad \text{implies} \quad f(x) \sim_Y f(x') \quad (x, x' \in X).$$

3.1.1.5 Terminology: isotone maps

Morphisms $(X, \preceq_X) \rightarrow (Y, \preceq_Y)$ are referred to as *order-preserving*, or *isotone* maps. The latter is common in literature on partially ordered sets.

3.1.1.6 Restriction to a subset

If $Y \subseteq X$ is a subset, then Y^I can be naturally identified with the subset of X^I of those functions from I to X whose values belong to Y . In particular, $R \cap Y^I$ becomes an I -ary relation on Y . We shall call it the *restriction* of relation R to Y , and denote it $R|_Y$.

3.2 Sets equipped with an operation

3.2.1 I -ary operations

3.2.1.1

An I -ary operation on a set X is a map

$$\mu: X^I \rightarrow X. \tag{3.2}$$

3.2.1.2 Commutativity

We say that operation (3.2) is *commutative* if the following diagram commutes

$$\begin{array}{ccc} X^I & & \\ \rho^* \downarrow & \searrow \mu & \\ X^I & \nearrow \mu & X \end{array} \quad (3.3)$$

for any bijection $\rho: I \rightarrow I$.¹ Note that $\rho^*: X^I \rightarrow X^I$ is the induced map, introduced in (1.8).

3.2.2 n -ary operations

3.2.2.1

An n -ary operation on a set X is a map

$$\mu: X^n \rightarrow X. \quad (3.4)$$

It can be viewed as an $(n+1)$ -ary relation

$$R_\mu = \{(x_0, x_1, \dots, x_n) \in X^{n+1} \mid x_0 = \mu(x_1, \dots, x_n)\}. \quad (3.5)$$

Exercise 47 Let X be a set and $R \subseteq X^{n+1}$. Show that there exists an n -ary operation, (3.4), such that $R = R_\mu$ if and only if R satisfies the following property

$$\text{for any } x_1, \dots, x_n \in X, \text{ there exists a unique element } x_0 \in X, \text{ such that } (x_0, x_1, \dots, x_n) \in R. \quad (3.6)$$

3.2.2.2

Sets with an n -ary operation are sometimes called *n -ary structures*. They form a full subcategory of the category of sets with an $(n+1)$ -ary relation.

Exercise 48 Show that $f: (X, \mu) \rightarrow (Y, \nu)$ is a morphism if and only if

$$f(\mu(x_1, \dots, x_n)) = \nu(f(x_1), \dots, f(x_n)) \quad (x_1, \dots, x_n \in X). \quad (3.7)$$

¹Self-bijections of I are called *permutations* of elements of set I .

3.2.2.3

For a subset $Y \subseteq X$ of a set with an n -ary operation (X, μ) , the restriction of R_μ to Y is an n -ary relation on Y which does not need to satisfy property (3.6).

Exercise 49 Show that $(R_\mu)|_Y = R_\nu$ for some n -ary operation ν on Y if and only if

$$\text{for any } y_1, \dots, y_n \in Y, \text{ one has } \mu(y_1, \dots, y_n) \in Y. \quad (3.8)$$

Show that, for all $y_1, \dots, y_n \in Y$,

$$\nu(y_1, \dots, y_n) = \mu(y_1, \dots, y_n).$$

3.2.2.4

In this case, we shall denote ν by μ_Y , call it the operation on Y induced by μ , and (Y, μ_Y) , the subset-with-operation of (X, μ) .

3.2.3 The category of sets with an n -ary operation

3.2.3.1 Homomorphisms

Traditionally, maps $f: X \rightarrow Y$ between sets equipped with an n -ary operation which satisfy identity (3.7) are referred as *homomorphisms*. This is where the term *morphism* originated.

3.2.3.2

Identity (3.7) is equivalent to the commutativity of the following diagram

$$\begin{array}{ccc} X^n & \xrightarrow{f_*} & Y^n \\ \mu \downarrow & & \downarrow \nu \\ X & \xrightarrow{f} & Y \end{array} \quad (3.9)$$

where $f_*(x_1, \dots, x_n) := (f(x_1), \dots, f(x_n))$.

3.2.3.3 Induced operations

3.2.3.4

An n -ary operation on a set induces several other n -ary operations on related sets. We shall consider here just two examples.

3.2.3.5 The induced operation on Y^X

The set of maps from a set X to a set Y which is equipped with an n -ary operation ν , is itself naturally equipped with a n -ary operation that is induced by ν .

For maps f_1, \dots, f_n , we define $\nu(f_1, \dots, f_n)$ as the map $X \rightarrow Y$ whose value at $x \in X$ is calculated by applying ν to the values of f_1, \dots, f_n at x :

$$\nu(f_1, \dots, f_n)(x) := \nu(f_1(x), \dots, f_n(x)) \quad (x \in X). \quad (3.10)$$

3.2.3.6 The induced operation on $\mathcal{P}(X)$

The set of subsets of a set X which is equipped with an n -ary operation μ , is itself naturally equipped with an n -ary operation that is induced by μ .

For subsets A_1, \dots, A_n of X , we define $\mu(A_1, \dots, A_n)$ as the set obtained by applying μ to every n -tuple $(a_1, \dots, a_n) \in A_1 \times \dots \times A_n$:

$$\mu(A_1, \dots, A_n) := \{\mu(a_1, \dots, a_n) \mid (a_1, \dots, a_n) \in A_1 \times \dots \times A_n\} \quad (3.11)$$

Exercise 50 Suppose that subsets A_1, \dots, A_n are finite. Show that $\mu(A_1, \dots, A_n)$ is finite by demonstrating the inequality

$$|\mu(A_1, \dots, A_n)| \preceq |A_1| \cdots |A_n|. \quad (3.12)$$

3.2.4 0-ary operations

3.2.4.1

For any set X , there is just a single map $\emptyset \rightarrow X$, namely the canonical inclusion map ι that embeds the empty set into X . Thus, the zeroth Cartesian power of any set X has a single element, namely ι , and therefore any 0-ary operation on a set X ,

$$X^0 \rightarrow X, \quad (3.13)$$

is the same as selecting a single element $\xi \in X$, the latter being the only value of map (3.13).

3.2.4.2 The category of sets with a distinguished element

In particular, sets equipped with a 0-ary operation are just sets with a distinguished element. Morphisms $(X, \xi) \rightarrow (Y, v)$ are the maps $f: X \rightarrow Y$ which are compatible with the distinguished elements, i.e.,

$$f(\xi) = v. \quad (3.14)$$

3.2.5 Unary operations

3.2.5.1

A unary operation on a set X is the same as a map $\phi: X \rightarrow X$. Such maps are often referred to as *selfmaps* on X .

3.2.5.2 The category of sets with a self-map

Morphisms $(X, \phi) \rightarrow (Y, \psi)$ are the maps $f: X \rightarrow Y$ which are compatible with the selfmaps, i.e., such that the diagram

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ \phi \downarrow & & \downarrow \psi \\ X & \xrightarrow{f} & Y \end{array} \quad (3.15)$$

commutes which translates into the identity $f \circ \phi = \psi \circ f$.

3.2.5.3

Certain sets possess *natural* unary operations, e.g., $\mathcal{P}(X)$ comes equipped with the ‘complement-of-a-subset’ self-map, cf. (2.25).

3.3 Binary structures

3.3.1 General binary structures

3.3.1.1 Binary structures

Sets equipped with a single binary operation are sometimes called *binary structures*. They form a full subcategory of the category of sets equipped with a binary relation, cf. Section 3.2.3. We shall denote it Bin .

3.3.1.2 Notation

Traditionally, for a binary operation on a set X an alternative notation is used:

$$x * y \quad \text{instead of} \quad \mu(x, y)$$

where $*$ here stands for any symbol denoting the operation. You may see here $+$, \times , \cdot , \otimes , and many other symbols.

3.3.1.3 Simplified notation

A frequent practice is to omit the symbol for the operation altogether and to write xy for $\mu(x, y)$.

3.3.1.4 Identity elements

An element $e \in X$ is a *left identity* if

$$\mu(e, x) = x \quad (x \in X).$$

Exercise 51 Formulate the notion of a right identity in a set with a binary operation, and show that, if e is a left identity and e' is a right identity, then $e = e'$. In particular, any set with a binary operation has no more than one two-sided identity.

3.3.1.5 Sink elements

An element $z \in X$ is a *left sink* if

$$\mu(z, x) = z \quad (x \in X).$$

Exercise 52 Formulate the notion of a right sink in a set with a binary operation, and show that, if z is a left sink and z' is a right sink, then $z = z'$.

In particular, any set with a binary operation has no more than one two-sided sink.

3.3.1.6 Idempotents

An element x of a binary structure is called an *idempotent* if

$$\mu(x, x) = x. \quad (3.16)$$

3.3.1.7 Commutative binary operations

In the special case of a binary operation, the general notion of commutativity introduced in Section 3.2.1.2 takes on the following form. A binary operation μ is *commutative* if it satisfies the identity

$$\mu(x, y) = \mu(y, x) \quad (x, y \in X). \quad (3.17)$$

The same in simplified notation:

$$xy = yx \quad (x, y \in X).$$

3.3.1.8

Identity (3.17) can be expressed as commutativity of the following diagram

$$\begin{array}{ccc}
 X \times X & & \\
 \downarrow \tau & \searrow \mu & \\
 & X & \\
 X \times X & \nearrow \mu & \\
 & &
 \end{array}
 \quad (3.18)$$

where $\tau: X \times X \rightarrow X \times X$ is the *flip*

$$\tau: (x, y) \mapsto (y, x) \quad (x, y \in X). \quad (3.19)$$

3.3.1.9 Additive notation and terminology

A commutative binary operation is often referred to as *addition*. In that case *additive notation* $x + y$ is used rather than $\mu(x, y)$ or xy .

3.3.1.10 Additive maps

In additive notation homomorphisms between commutative binary structures $f: X \rightarrow Y$ are just *additive maps*

$$f(x + x') = f(x) + f(x') \quad (x, x' \in X).$$

3.3.1.11

Given a binary structure (X, μ) , define a binary relation \sim_μ on X by

$$x \sim_\mu y \quad \text{if} \quad \mu(x, y) = y \quad (x, y \in X). \quad (3.20)$$

Exercise 53 Show that relation \sim_μ defined in (3.20) is reflexive if and only if every element in (X, μ) is idempotent.

Exercise 54 Show that relation \sim_μ defined in (3.20) is weakly antisymmetric, cf. (1.70), if μ is commutative.

3.3.1.12 Associative binary operations

A binary operation is said to be *associative* if it satisfies the identity

$$\mu(\mu(x, y), z) = \mu(x, \mu(y, z)) \quad (x, y, z \in X). \quad (3.21)$$

The same in simplified notation:

$$(xy)z = x(yz) \quad (x, y, z \in X).$$

Exercise 55 Show that relation \sim_μ defined in (3.20) is transitive, cf. (1.68), if μ is associative.

3.3.1.13

Identity (1.49) can be expressed as commutativity of the following diagram

$$\begin{array}{ccc} X \times X \times X & \xrightarrow{\mu \times \text{id}_X} & X \times X \\ \text{id}_X \times \mu \downarrow & & \downarrow \mu \\ X \times X & \xrightarrow{\mu} & X \end{array} \quad (3.22)$$

3.3.2 Binary versus I -ary operations

3.3.2.1

A binary operation $\mu: X \times X \rightarrow X$ allows to convert 2 elements of a set into a single element. What about 3 or more elements? Given a list of n elements

$$x_1, \dots, x_n,$$

we have to apply it first to a single pair of consecutive elements, say x_i and x_{i+1} , and replace that pair with the result. The new list has length $n - 1$. By iterating this procedure $n - 1$ times we eventually obtain a list consisting of a single element. This is the final result.

3.3.2.2 Iterated n -ary operations

If we recorded the sequence of steps performed, we can use the same recipe to any n -tuple of elements of X . The resulting map $X^n \rightarrow X$ is what we call an *iterated n -ary operation* induced by a binary operation.

A binary operation induces two ternary operations,

$$(x_1, x_2, x_3) \mapsto \mu(\mu(x_1, x_2), x_3) \quad (x_1, x_2, x_3 \in X) \quad (3.23)$$

and

$$(x_1, x_2, x_3) \mapsto \mu(x_1, \mu(x_2, x_3)) \quad (x_1, x_2, x_3 \in X), \quad (3.24)$$

five quaternary operations $X^4 \rightarrow X$, etc. There are exactly

$$\frac{1}{n} \binom{2n-2}{n-1} \quad (3.25)$$

induced iterated n -ary operations in total, and all of them are different in general. The number of induced iterated n -ary operations coincides with the number of nested sequences of $n-1$ pairs of parentheses, and is known as the $(n-1)$ -st *Catalan* number.

3.3.2.3

Associativity of μ states that the two ternary operations above, (3.23)–(3.24), coincide. Using this fact, one can show by induction on n that all the iterated n -ary operations coincide. Thus, an associative binary operation induces exactly one n -ary operation for each $n \geq 2$.

3.3.2.4

Commutativity of an associative binary operation has one more advantage: it allows one to extend the operation to families of elements $(x_i)_{i \in I}$ of X indexed by *arbitrary* finite nonempty sets I , producing I -ary operations

$$\mu_I: X^I \rightarrow X. \quad (3.26)$$

For a general associative operation, performing the iterated operation requires that the indexing set be ordered, which amounts to providing a bijection between the set $\{1, \dots, n\}$ and I

$$I = \{i_1, \dots, i_n\}.$$

Then family $(x_i)_{i \in I}$ becomes a list

$$x_{i_1}, \dots, x_{i_n}$$

and we apply μ to that list as explained above.

If I has n elements, there are exactly $n!$ different orderings of I , and therefore an associative binary operation induces exactly $n!$ operations (3.26). They all coincide if μ is commutative and we denote this unique I -ary operation μ_I .

3.3.2.5 Associativity seen through the induced I -ary operations

Given a *finite* family $(\zeta_j)_{j \in J}$ of *finite* families $\zeta_j = (x_{i_j})_{i \in I_j}$ of elements of X , we can evaluate μ_{I_j} on each ζ_j to get

$$(\mu_{I_j}(\zeta_j))_{j \in J} \in X^J$$

and subsequently evaluate μ_J on it. We can also apply μ_L to the *total* family

$$(x_l)_{l \in L}$$

indexed by the disjoint sum

$$L = \coprod_{j \in J} I_j,$$

cf. Section 1.3.1.2. The results are equal

$$\mu_J \left((\mu_{I_j}((x_{i_j})_{i \in I_j}))_{j \in J} \right) = \mu_L((x_l)_{l \in L}). \quad (3.27)$$

If we simplify notation by omitting indexing sets, then identity (3.27) becomes a little easier to read

$$\mu_J \left((\mu_{I_j}((x_{i_j})))_{j \in J} \right) = \mu_L((x_l)). \quad (3.28)$$

3.3.2.6

Note that identity (3.27) holds even if some of the indexing sets I_j have a single element, provided that

$$\mu_I: X^I \rightarrow X$$

is to be understood as the canonical bijection that identifies $X^{\{\bullet\}}$ with X :

$$X^{\{\bullet\}} \ni \{\bullet \mapsto x\} \longleftrightarrow x \in X.$$

3.3.2.7 A comment on notation

Above, I is a set with a single element and we denoted that single element \bullet . Mathematical notation employs symbolic ‘names’ like a , b , c , etc. in order to distinguish between different elements of a set. There is no need to do that when the set has a single element. We denoted the single element of set I by \bullet to indicate the fact that we do not need to ‘name’ it first before we can make a reference to it.

3.3.2.8 A comment on the meaning of identity (3.27)

Identity (3.27) expresses compatibility of the system of induced operations μ_I and is a manifestation of associativity of the original binary operation.

3.3.2.9

'Associativity' identity (3.27) becomes more legible when we express it as commutativity of the diagram

$$\begin{array}{ccc}
 \prod_{j \in J} X^{I_j} & \xrightarrow{\prod_{j \in J} \mu_{I_j}} & \prod_{j \in J} X \\
 \uparrow (1.37) & & \downarrow \mu_I \\
 X^L & \xrightarrow{\mu_L} & X
 \end{array} \tag{3.29}$$

where the left vertical arrow is the canonical identification of Cartesian products discussed in Sections 1.3.2.2 and 1.3.2.3.

3.3.2.10 Iterated operations in additive notation

If we use additive notation and terminology, then

$$\mu_I((x_i)_{i \in I}) \quad \text{becomes} \quad \sum_{i \in I} x_i$$

and identity (3.27) becomes

$$\sum_{j \in J} \sum_{i_j \in I_j} x_{i_j} = \sum_{l \in L} x_l. \tag{3.30}$$

3.3.2.11

In this form associativity identity (3.27) seems much easier to comprehend than in the original form which employs functional notation for operations. This illustrates the fact that the notation we use indeed is either aiding or hindering our human comprehension.

It is one of the first duties of a professional mathematician to pay due respect to proper notation, and to always strive for notation that is simultaneously precise, clear, and suitable. One has to constantly negotiate between these goals which are at times not easy to reconcile.

3.3.3 Semigroups

3.3.3.1

An associative binary structure (X, μ) is called a *semigroup*.

3.3.3.2 The category of semigroups

Semigroups form a full subcategory of the category of sets with a binary operation, and therefore also a full subcategory of the category of sets with a ternary relation.² The category of semigroups will be denoted Semigrp .

3.3.3.3 Subsemigroups

Subsets-with-operation (Y, μ_Y) of a semigroup (X, μ) are called *subsemigroups*. The canonical inclusion of Y into X is then a homomorphism of semigroups.

Exercise 56 Let $(T_i)_{i \in I}$ be a family of subsemigroups of a semigroup S . Show that

$$\bigcap_{i \in I} T_i$$

is a subsemigroup of S .

3.3.3.4 The subsemigroup generated by a subset

The set of subsemigroups of a semigroup S is contained in $\mathcal{P}(S)$ and thus ordered by inclusion. It follows from Exercise 56 that, for any subset $X \subseteq S$,

$$\langle X \rangle := \bigcap_{\substack{T \text{ a subsemigroup of } S \\ T \supseteq X}} T, \quad (3.31)$$

is the *smallest* subsemigroup of A which contains X . We call it the subsemigroup *generated* by subset X .

Exercise 57 Show that $\langle X \rangle = X$ if and only if X is a subsemigroup.

Exercise 58 Show that $\langle \langle X \rangle \rangle = \langle X \rangle$.

²Ternary means $n = 3$.

3.3.3.5 A set of generators

We say that $X \subseteq A$ *generates* semigroup A , or *is a set of generators* for A , if $\langle X \rangle = A$.

3.3.3.6 Semigroups as categories with a single object

When a category \mathcal{C} has a single object, the structure of the category is uniquely determined by the set

$$\mathcal{C}_1 = \text{Hom}_{\mathcal{C}}(\bullet, \bullet)$$

where \bullet denotes the only object of \mathcal{C} , and the associative composition map

$$\mathcal{C}_2 = \mathcal{C}_1^2 \rightarrow \mathcal{C}_1.$$

In other words, the set of morphisms forms a semigroup under composition. Vice-versa, given any semigroup (X, μ) , one can associate with it the following category

$$\mathcal{C}_0 := \{\bullet\}, \quad \mathcal{C}_1 = \text{Hom}_{\mathcal{C}}(\bullet, \bullet) := X,$$

with μ playing the role of the composition map.

Exercise 59 Show that functors between categories with a single object are in one-to-one correspondence with homomorphisms of semigroups.

3.3.4 Examples of semigroups

3.3.4.1

Exercise 60 Let X be a set. Show that the canonical projections

$$p_1: (x, y) \mapsto x \quad (x, y \in Y) \tag{3.32}$$

and

$$p_2: (x, y) \mapsto y \quad (x, y \in Y), \tag{3.33}$$

are associative.

Exercise 61 Show that every element in X is a right identity for (X, p_1) and a left identity for (X, p_2) , and that (X, p_1) has a two-sided identity precisely when X has a single element.

3.3.4.2 Semilattices

A partially ordered set (S, \preceq) is called a *semilattice* if, for any $s, t \in S$, the set $\{s, t\}$ has supremum.

Exercise 62 Show that the operation

$$(s, t) \mapsto s \vee t := \sup\{s, t\} \quad (s, t \in S), \quad (3.34)$$

is associative.

Exercise 63 Show that the semigroup (S, \vee) has an identity element if and only if semilattice S has the smallest element.

Exercise 64 Show that the semigroup (S, \vee) has a sink element if and only if semilattice S has the largest element.

Exercise 65 Let f be a map between two semilattices (S, \preceq) and (S', \preceq') . Show that f is isotone, i.e.,

$$\text{if } s \preceq t, \text{ then } f(s) \preceq f(t) \quad (s, t \in S), \quad (3.35)$$

if and only if

$$f(s) \vee f(t) \preceq' f(s \vee t) \quad (s, t \in S). \quad (3.36)$$

3.3.4.3

The semigroup (S, \vee) associated with a semilattice (S, \preceq) is commutative and every element is an idempotent.

3.3.4.4

Vice-versa, if (X, μ) is a commutative semigroup with the property that every element is idempotent, then X equipped with relation \sim_μ introduced in Section (3.20), becomes a partially ordered set, cf. Exercises 53–55.

Exercise 66 Show that \sim_\vee coincides with the original partial order relation \preceq .

Exercise 67 Show that (X, \sim_μ) is a semilattice. More precisely, show that

$$\sup\{x, y\} = \mu(x, y) \quad (x, y \in X). \quad (3.37)$$

3.3.4.5

Identity (3.37) means that the \vee -operation corresponding to \sim_μ is the original μ -operation.

By combining everything together, we arrive at the following observation.

Proposition 3.3.1 *For any set X , there exists a natural correspondence between partial order relations which make X into a semilattice, and binary operations which make X into a commutative semigroup where every element is idempotent.*

3.3.4.6 A word of caution

Proposition 3.3.1 seems to say that there is a natural isomorphism between the category semilattices and the category commutative semigroups in which every element is idempotent.

This is indeed so if one properly understands what to consider to be a morphism of semilattices. Note that a map $f: X \rightarrow X'$ is a homomorphism of semigroups $(X, \vee) \rightarrow (X', \vee')$ if and only if f is a finitely sup-continuous morphism of partially ordered sets $(X, \preceq) \rightarrow (X', \preceq')$, i.e., if

$$f\left(\sup_{(X, \preceq)} A\right) = \sup_{(X', \preceq')} f(A) \quad (3.38)$$

for any nonempty finite subset $A \subseteq X$.

3.3.4.7 Right-exact maps

Let us call a map $f: X \rightarrow X'$ between partially ordered sets *right-exact* if it preserves the suprema of nonempty finite sets, i.e., it satisfies (3.38) whenever $\sup A$ exists and $A \subseteq X$ is nonempty and finite. According to Exercise 65 such a map is automatically isotone.

3.3.4.8 The category of semilattices

If by the category of semilattices we understand the subcategory of the category of partially ordered sets with morphisms being right-exact maps, then the category of semilattices is naturally isomorphic to the full subcategory of the category of commutative semigroups formed by semigroups where every element is idempotent.

3.3.4.9

Note that the identity

$$\mu(x, x) = x \quad (x \in X) \quad (3.39)$$

which expresses the fact that every element is idempotent, like many other such identities can be also expressed, without resorting to elements of X , as commutativity of a certain diagram, in this case:

$$\begin{array}{ccc} X & & \\ \parallel & \searrow \Delta & \\ X & & X \times X \\ & \swarrow \mu & \\ & X & \end{array} \quad (3.40)$$

where $\Delta: X \rightarrow X \times X$ denotes the *diagonal* embedding of X into $X \times X$:

$$\Delta: x \mapsto (x, x) \quad (x \in X). \quad (3.41)$$

3.3.4.10 Subsemilattices

A partially ordered subset $(A, \preceq|_A)$ of a semilattice (X, \preceq) does not need to be a semilattice. Indeed, the set $A = \{x, y\}$, for any two elements $x, y \in X$ which are not comparable, lacks both $\sup\{x, y\}$ and $\inf\{x, y\}$.

If, however, it is, we should consider it to be a *subsemilattice* of (X, \preceq) only if the inclusion map $A \hookrightarrow X$ is a morphism in the category of semilattices, i.e., is a right-exact map.

3.3.4.11

What we described above should be called sup-semilattices. By replacing sup with inf, one obtains the concept of an inf-semilattice. The theories are of course identical, since X with the reverse order,

$$x \preceq^{\text{rev}} x' \quad \text{if} \quad x' \preceq x \quad (x, x' \in X), \quad (3.42)$$

is a sup-semilattice precisely when (X, \preceq) is an inf-semilattice.³

Exercise 68 State for \wedge -semilattices the analog of the statement of Exercise 65.

³ sup-semilattices are also called *join-semilattices*, or \vee -semilattices; inf-semilattices are also called *meet-semilattices*, or \wedge -semilattices.

3.3.4.12 The semigroup of maps with values in a semigroup

The set of maps S^X from a set X into a semigroup S is naturally a semigroup: the binary operation is applied pointwise to the values, cf. Section 3.2.3.5, and associativity is an immediate consequence of associativity of the operation in S .

3.3.4.13

When X is equipped with a binary operation of its own, we can consider the subset of S^X formed by homomorphisms from X to S . In general, the product of two homomorphisms is not a homomorphism, unless they *commute*:

$$gf = fg. \quad (3.43)$$

Exercise 69 Show that the product fg of two homomorphisms from a binary structure X to a semigroup S is a homomorphism if f commutes with g .

3.3.4.14

It follows that if S is a commutative semigroup, then $\text{Hom}_{\text{Bin}}(X, S)$ is a subsemigroup of S^X .

3.3.4.15 The set of endomorphisms of an object

The set of endomorphisms $\text{End}_{\mathcal{C}}(a)$ of an object a in an arbitrary category \mathcal{C} is a monoid.

3.3.5 Monoids

3.3.5.1

Semigroups with a two-sided identity are called monoids. A homomorphism of semigroups does not necessarily send the identity element to the identity element, as the following simple example demonstrates:

$$X = M_2(\mathbf{Z}), \quad Y = \left\{ \begin{pmatrix} m & 0 \\ 0 & 0 \end{pmatrix} \mid m \in \mathbf{Z} \right\},$$

and the operation is the multiplication of 2×2 -matrices. Since Y is a subsemigroup of X , the inclusion of Y in X is a homomorphism of

semigroups. However, the identity element of Y ,

$$\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

is not the identity element of X ,

$$\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}.$$

3.3.5.2

In view of this, one additionally requires from a morphism of *monoids* that it respects the identity elements. In particular, the category of monoids is a subcategory of the category of semigroups, yet not a full subcategory. We shall denote it Mon .

3.3.5.3 An example: $\text{End}_{\mathcal{C}}(a)$

In a unital category, the set of endomorphisms $\text{End}_{\mathcal{C}}(a)$ of any object a is a monoid.

3.3.5.4 Submonoids

For the same reason, submonoids of (X, μ) are not just subsemigroups (Y, μ_Y) which happen to be monoids, but the subsemigroups which contain the identity element

3.3.5.5

Note that $\{e\}$ is the smallest submonoid (frequently referred to as a *trivial* submonoid) while \emptyset is the smallest subsemigroup.

3.3.5.6

Intersection of a family of submonoids is a submonoid. In particular, for any subset X of a monoid there exists the smallest submonoid that contains X . We call it the submonoid *generated by* X . e of (X, μ) .

3.3.5.7 Monoids as categories with a single object

Monoids correspond to unital categories with a single object, and homomorphisms between monoids correspond to unital functors.

3.3.5.8 Invertible elements

An element u of a monoid X is said to be a *left inverse* of an element x if

$$ux = e$$

where e is the identity element.

Exercise 70 *Formulate the notion of a right inverse and show that in a monoid, if u is a left inverse of x , and v is a right inverse of x , then $u = v$.*

3.3.5.9

In particular, every element x in a monoid has no more than one two-sided inverse. This unique element is denoted x^{-1} (if one uses *multiplicative* notation for the operation), and x is said to be *invertible*.

3.3.5.10

Invertible elements in a monoid correspond to isomorphisms in the associated category.

Exercise 71 *Show that any homomorphism of monoids $f: (X, \mu) \rightarrow (Y, \nu)$, sends invertible elements in X to invertible elements in Y . More precisely, show that for any such element, $(f(x))^{-1} = f(x^{-1})$.*

3.3.6 Groups

3.3.6.1

A monoid (X, μ) is called a *group*, if *every* element $x \in X$ is invertible. In view of the above exercise, it is natural to consider groups as the full subcategory of the category of monoids. This category is often denoted Grp .

In contrast to monoids, groups form a full subcategory of the category of sets with a binary operation. In particular, Grp is a full subcategory of the category of semigroups.

Exercise 72 *Let (G, μ) and (H, ν) be two groups, and $f: (G, \mu) \rightarrow (H, \nu)$ be a homomorphism of sets with a binary operation. Show that $f(e_G) = e_H$.*

3.3.6.2 Abelian groups

Commutative groups are called *abelian groups* in view of a long established tradition that predates nearly all the other terminology employed here.⁴ Abelian groups form, of course, a full subcategory of \mathbf{Grp} . It is denoted \mathbf{Ab} .

3.3.6.3 Groupoids

Groups correspond to categories with a single object and the property that any morphism is an isomorphism. For this reason, categories with the same property are called *groupoids*.

3.3.6.4 A comment about notation

If (X, μ) is a semigroup, monoid, or a group, it is customary to refer to X alone as a semigroup, monoid or, respectively, a group. This rarely leads to terminological confusion if the operation is clear from the context and often greatly simplifies notation. We shall follow this convention in the future.

3.3.6.5 A comment about terminology

The binary operation in a general semigroup, monoid, or a group, X , is often referred to as the *multiplication* in X .

Exercise 73 Show that in any monoid M , the set of invertible elements $G(M)$ is a group with respect to the operation induced by the multiplication in M .

3.3.6.6

Combined Exercises 73 and 71 show that associating with a monoid X the group of its invertible elements $G(X)$ defines a functor $\mathbf{Mon} \rightsquigarrow \mathbf{Grp}$.

⁴Named after Niels Henrik Abel (1802-1829), a Norwegian mathematician who proved impossibility of solving by radicals a general polynomial equation of degree greater than 4. Abel also proved that the equation was solvable by radicals if the group of automorphisms of the equation was commutative.

Introduced by Camille Jordan (1838-1922) in 1870, who used the term *groupe ablien* to denote some specific groups of matrices, this terminology is applied to general commutative groups later, first perhaps in a 1882 article by Heinrich Martin Weber (1842-1913).

Exercise 74 Let M be a monoid and $\iota: G(M) \hookrightarrow M$ denote the canonical inclusion of the group of invertible elements of M into M . Note that ι is a homomorphism of monoids.

Show that, for any group G and any homomorphism of monoids $f: G \rightarrow M$, there exists a unique homomorphism of groups $\tilde{f}: G \rightarrow G(M)$ such that $f = \iota \circ \tilde{f}$.

3.4 Sets with a pair of binary operations

3.4.1 Introduction

3.4.1.1

Structures involving a pair of binary operations on a given set are both very common and very important. An essential feature of such structures are ‘compatibility’ conditions that relate one of the two operations to the other one. These conditions are usually expressed in the form of identities.

3.4.1.2 Distributivity

The most important of all of these conditions is *distributivity*. Given two binary operations \bullet and \circ on a set X , we say that operation \circ is *left-distributive over* operation \bullet if the following identity holds

$$x \circ (y \bullet z) = (x \circ y) \bullet (x \circ z) \quad (x, y, z \in X). \quad (3.44)$$

Exercise 75 Formulate the definition of right-distributivity of \circ over \bullet .

Exercise 76 Consider union and intersection as binary operation on $\mathcal{P}(X)$. Show that \cap distributes over \cup and \cup distributes over \cap .⁵

3.4.1.3

Theory of *Lie algebras* is founded on *Jacobi identity*:

$$x \bullet (y \bullet z) + y \bullet (z \bullet x) + z \bullet (x \bullet y) = 0 \quad (x, y, z \in X) \quad (3.45)$$

One of the two operations is referred to as *addition*, the other—as the *Lie bracket* operation: the standard notation for $x \bullet y$ is $[x, y]$.

⁵This is a very rare situation when two binary operations distribute over each other.

3.4.1.4

In the context of lattices we shall encounter the pair of *absorption identities*, and the *modular identity*.

3.4.2 Semirings

3.4.2.1 Biadditive pairings

Suppose that commutative semigroups S , T , and U be given. We shall use additive notation and terminology throughout.

A map

$$\mu: S \times T \rightarrow U \quad (3.46)$$

is said to be *biadditive*, or a *biadditive pairing*, if it is additive in each argument:

$$\mu(s + s', t) = \mu(s, t) + \mu(s', t) \quad (s, s' \in S; t \in T), \quad (3.47)$$

and

$$\mu(s, t + t') = \mu(s, t) + \mu(s, t') \quad (s \in S; t, t' \in T). \quad (3.48)$$

3.4.2.2

Left-additivity of (3.47) expresses the fact that μ right-distributes over addition. Similarly, Right-additivity condition (3.48) expresses the fact that μ right-distributes over addition.

3.4.2.3 Semirings

A commutative semigroup S equipped with a biadditive binary operation

$$\mu: S \times S \rightarrow S \quad (3.49)$$

is called a *semiring*.

3.4.2.4 The additive semigroup of a semiring

In a semiring the original semigroup operation is referred to as *addition* and the corresponding semigroup as the *additive* semigroup of the semiring. We shall refer to semigroup $(S, +)$ as the *additive* semigroup of the semiring, and will denote it S^+ in order to distinguish it from S viewed as a semiring.

3.4.2.5 Multiplication

We will refer to biadditive operation (3.49) as the *multiplication*, and will usually denote $\mu(s, t)$ by $s \cdot t$ or st . Equipped with multiplication S is just a binary structure. We will denote it S^\times .

3.4.2.6 The category of (nonassociative) semirings

Morphisms $(S, +, \cdot) \rightarrow (T, +, \cdot)$ are maps $S \rightarrow T$ which are simultaneously homomorphisms of the additive semigroups $S^+ \rightarrow T^+$ and of multiplicative binary structures $S^\times \rightarrow T^\times$. Traditionally, such maps are called *homomorphisms* of semirings.

3.4.2.7

Terminology like an *associative* (resp. *commutative*, *unital*) semiring always refers to the corresponding properties of the multiplication. The identity element for multiplication is usually called *identity* or *unit*, and is most of the time denoted 1.

3.4.2.8 A comment about terminology

Semirings form a full subcategory of the category of sets with *two* binary operations. Associativity, however, is such an important property that a common practice is to tacitly assume it when speaking of semirings. From now on, the phrase *nonassociative* ring will refer to semirings that are *not assumed* to be associative. Note that such a reference does not preclude associativity.

3.4.2.9 The category of associative semirings

We shall denote the category of associative semirings by *Semiring* and will refer to its object simply as ‘semirings’. It is a full subcategory of the category of *nonassociative* semirings.

3.4.2.10 The multiplicative semigroup of an associative semiring

When S is an associative semiring, S^\times is a semigroup. We shall refer to it as the *multiplicative semigroup* of S .

3.4.2.11 Zero

If the additive semigroup of a semiring is a monoid, its identity element is denoted 0 and referred to as *zero*.

3.4.2.12 Semirings with zero

A *semiring with zero* is a semiring whose additive semigroup is a monoid and zero satisfies the following identity

$$0 \cdot s = 0 = s \cdot 0 \quad (s \in S). \quad (3.50)$$

Identity (3.50) means that 0 is a sink of the multiplicative semigroup, cf. Section 3.3.1.5.

3.4.3 Examples of semirings

3.4.3.1 $[0, \infty)$ and $[0, \infty]$

The set $[0, \infty)$ of nonnegative real numbers equipped with usual addition and multiplication of real numbers, forms an associative and commutative semiring with zero.

The set $[0, \infty] := [0, \infty) \cup \infty$ of *extended* nonnegative real numbers can be equipped with a semiring structure by extending addition and multiplication of real numbers as follows

$$a + \infty = \infty = \infty + a \quad \text{and} \quad a \cdot 0 = 0 = 0 \cdot a \quad (a \in [0, \infty]).$$

Note that ∞ is a sink of the additive monoid of $[0, \infty]$ while 0 is a sink of the multiplicative monoid of $[0, \infty]$.

3.4.3.2 The near-semiring S^S

The set of selfmaps $S \rightarrow S$ possesses two semigroup structures when S is a semigroup. The first one is obtained when we consider S^S as the set of all maps from *set* S to *semigroup* S : the operation is pointwise multiplication \cdot , as defined in Section 3.2.3.5. The other operation is composition of maps which endows S^S with a structure of a monoid.

Exercise 77 Show that \circ is right-distributive over \cdot .

3.4.3.3

Composition in S^S is practically never left-distributive over \cdot as even the simplest examples demonstrate.

3.4.3.4 Example demonstrating that S^S is not left-distributive

The two-element set $S = \{\pm 1\}$ equipped with usual multiplication of integers is a group. Let $f: S \rightarrow S$ be the constant map that sends both 1 and -1 to -1 . Then

$$f \circ (f \cdot f) = f$$

while

$$(f \circ f) \cdot (f \circ f) = f \cdot f$$

is the constant map that sends both 1 and -1 to 1.

3.4.3.5

Equipped with pointwise multiplication and composition, the S^S is an example of a *near-semiring*, a structure more general than a semiring. As we shall discover in a moment, under additional hypothesis that S is commutative, there is a true semiring inside of S^S .

3.4.3.6

For any semigroup S , the set $\text{End}_{\text{Semigrp}}(S)$ is a submonoid of (S^S, \circ) . As we noted in Section 3.3.4.14, it is also a subsemigroup of (S^S, \cdot) when S is commutative.

Exercise 78 Assuming $(S, +)$ to be a commutative semigroup, show that composition left-distributes over addition in $\text{End}_{\text{Semigrp}}(S)$.

3.4.3.7 The semiring of endomorphisms of a commutative semigroup

It follows from Exercises 77 and 78 that $(\text{End}_{\text{Semigrp}}(S), +, \circ)$ is a semiring. It is unital: the identity morphism id_S is its multiplicative identity. It is a semiring with zero precisely when $(S, +)$ is a monoid.

3.4.4 Rings

3.4.4.1

Semirings whose additive semigroup is a group are called *rings*.

3.4.4.2 One more comment about terminology

The remarks made in Section 3.4.2.8 apply here too: it is a common practice to tacitly assume associativity when speaking of rings, and to use the designation *nonassociative ring* when associativity is not assumed.

3.4.4.3 The category of associative rings

Nonassociative rings form a full subcategory of the category of nonassociative semirings. Similarly, associative rings form a full subcategory of the category of associative semirings.

The category of associative rings will be denoted \mathbf{Ring} and we will refer to its objects as ‘rings’.

3.4.5 Examples of rings

3.4.5.1 The ring of endomorphisms of an abelian group

For an abelian group $(A, +)$, the semiring $\text{End}_{\text{Ab}}(A)$ which was introduced in Section 3.4.3.7 is a unital ring.

3.4.6 Lattices

3.4.6.1

A partially ordered set (L, \preceq) is said to be a *lattice* if any nonempty finite subset $A \subseteq L$ has both supremum and infimum.

3.4.6.2

In particular, we obtain the binary operations on L ,

$$l \vee m := \sup\{l, m\} \quad (\text{‘join’}) \quad (3.51)$$

and

$$l \wedge m := \inf\{l, m\} \quad (\text{‘meet’}), \quad (3.52)$$

are commutative, associative, and every element $l \in L$ is an idempotent.

3.4.6.3 Absorption identities

The two operations are related to each other through the following pair of identities

$$l \wedge (l \vee m) = l \quad (l, m \in L) \quad (3.53)$$

and

$$l \vee (l \wedge m) = l \quad (l, m \in L). \quad (3.54)$$

Exercise 79 *Prove identities (3.53)–(3.54).*

3.4.6.4 Consequences of absorption identities

Let L be a set equipped with two binary operations \vee and \wedge which satisfy the above pair of absorption identities. Then, for any $l \in L$,

$$l \wedge l \stackrel{(3.53)}{=} l \wedge (l \vee (l \wedge l)) \stackrel{(3.54)}{=} l.$$

In other words, \wedge -idempotence of all elements is a consequence of the absorption identities.

Exercise 80 *Using just absorption identities show that*

$$l \vee l = l \quad (l \in L).$$

Exercise 81 *Using just absorption identities show that*

$$l = l \wedge m \quad \text{if and only if} \quad l \vee m = m \quad (l, m \in L). \quad (3.55)$$

3.4.6.5

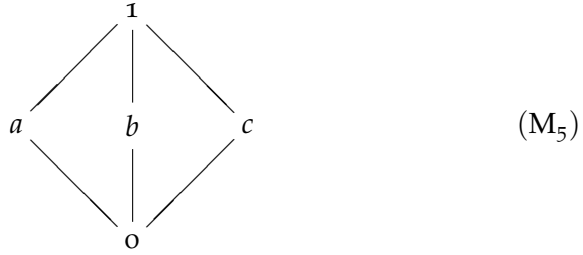
None of the two distributivity identities

$$l \wedge (m \vee n) = (l \wedge m) \vee (l \wedge n) \quad (l, m, n \in L) \quad (3.56)$$

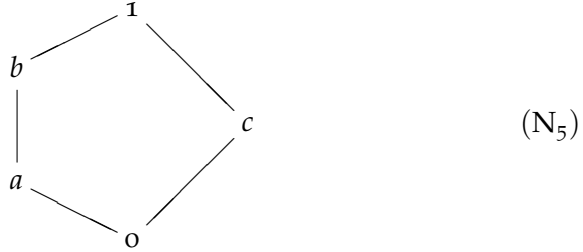
or

$$l \vee (m \wedge n) = (l \vee m) \wedge (l \vee n) \quad (l, m, n \in L) \quad (3.57)$$

holds in general. The minimal examples are two lattice structures on a set with 5-elements whose Hasse diagrams are



and



(we denoted $\inf L$ by 0 and $\sup L$ by 1).⁶

Exercise 82 Show that in any lattice the following two distributivity inequalities always hold

$$(l \wedge m) \vee (l \wedge n) \preceq l \wedge (m \vee n) \quad (l, m, n \in L) \quad (3.58)$$

and

$$(l \vee m) \wedge (l \vee n) \preceq l \vee (m \wedge n) \quad (l, m, n \in L). \quad (3.59)$$

3.4.6.6

Remarkably, if either one of identities (3.56) or (3.57) holds, the other one holds too. This is yet another consequence of the pair of absorption identities in combination with commutativity of \wedge and associativity of \vee .

⁶Hasse diagrams are visual presentations of partially ordered sets as graphs whose nodes correspond to the elements of the partially ordered set, and $a < b$ if there exists a 'path upwards' from node a to node b . Thus, in lattice N_5 , $a < b$ and neither one is comparable with c .

To derive identity (3.57) from (3.56), we note that

$$(l \vee m) \wedge (l \vee n) = ((l \vee m) \wedge l) \vee ((l \vee m) \wedge n) \quad \text{identity (3.56)} \quad (3.60)$$

$$= l \vee ((l \vee m) \wedge n) \quad (\vee)\wedge\text{-absorption} \quad (3.61)$$

$$= l \vee ((l \wedge n) \vee (m \wedge n)) \quad \text{identity (3.56)} \quad (3.62)$$

$$= (l \vee (l \wedge n)) \vee (m \wedge n) \quad \text{associativity of } \vee \quad (3.63)$$

$$= l \vee (m \wedge n) \quad \vee(\wedge)\text{-absorption} \quad (3.64)$$

Exercise 83 Derive identity (3.56) from (3.57).

3.4.6.7

We saw that to any lattice structure on a set L corresponds a pair of commutative and associative binary operations on L which satisfy the pair of absorption identities. In fact, this is a bijective correspondence: every such pair of binary operations arises from a unique lattice structure on L .

Proposition 3.4.1 *For any set L , there exists a natural correspondence between partial order relations that make L into a lattice, and pairs of commutative and associative binary operations which satisfy absorption identities.*

Proof. Suppose \wedge and \vee is a pair of commutative and associative operations on a set L , satisfying absorption identities. If we consider the associated relations \sim_\wedge and \sim_\vee , then (L, \sim_\wedge) and (L, \sim_\vee) are partially ordered sets. Moreover,

$$\sup_{(L, \sim_\wedge)} \{l, m\} = l \wedge m \quad \text{and} \quad \sup_{(L, \sim_\vee)} \{l, m\} = l \vee m \quad (l, m \in L),$$

cf. Exercise 67. Equivalence of equalities in (3.55) means that partial order \sim_\vee is the reverse of \sim_w , and in the reverse partial order infimum and supremum are switched. Thus,

$$\inf_{(L, \sim_\vee)} \{l, m\} = \sup_{(L, \sim_\wedge)} \{l, m\} = l \wedge m \quad (l, m \in L)$$

completing the proof that (L, \sim_\vee) is a lattice.

In the opposite direction, we already demonstrated that, for any lattice (L, \preceq) , operations (3.51)-(3.52) are commutative, associative, and satisfy absorption identities (Exercise 79).

The fact that the correspondences

$$\preceq \rightsquigarrow (\vee, \wedge) \quad \text{and} \quad (\vee, \wedge) \rightsquigarrow \sim_\vee$$

are mutually inverse follows from the corresponding fact established earlier for semilattices, combined with the fact that partial orders \sim_\wedge and \sim_\vee are reverse for each other. \square

3.4.6.8 The category of lattices

If we understand by the category of lattices the subcategory of the category of partially ordered sets whose objects are lattices and morphisms are supposed to be *exact maps*, i.e., to preserve suprema *and* infima of nonempty finite subsets, then it is a corollary of Proposition 3.4.1 that the category of lattices is isomorphic to the category of sets equipped with a pair of commutative and associative binary operations satisfying the absorption identities.

3.4.6.9 Sublattices

Comments made for semilattices apply here as well: A partially ordered subset $(X, \preceq|_X)$ of a lattice (L, \preceq) will be called a *sublattice* of (L, \preceq) if the inclusion map $X \hookrightarrow L$ is a morphism in the category of lattices, i.e., is an exact map.

3.4.6.10 Bounded lattices

A *bounded lattice* is a lattice that contains the largest and the smallest elements.

3.4.6.11 Complete lattices

A partially ordered set is a *complete lattice* if *every* subset has supremum and infimum.

Exercise 84 Let f be a map between two complete lattices (S, \preceq) and (S', \preceq') . Show that f is isotone if and only if

$$\sup_{(S', \preceq')} f(E) \preceq' f(\sup_{(S, \preceq)} E) \quad (E \subseteq S), \quad (3.65)$$

if and only if

$$f(\inf_{(S, \preceq)} E) \preceq' \inf_{(S', \preceq')} f(E) \quad (E \subseteq S). \quad (3.66)$$

3.4.7 Distributive lattices

3.4.7.1

We saw above that if \wedge distributes over \vee , then \vee distributes over \wedge .

3.4.7.2

The class of distributive lattices constitutes an ‘intersection’ between the class of commutative semirings and the class of lattices.

3.4.8 Examples of distributive lattices

3.4.8.1 $\mathcal{P}(X)$

The set of subsets of a set X provides perhaps the most important example of a complete distributive lattice. Note that semirings $(\mathcal{P}(X), \cup, \cap)$ and $(\mathcal{P}(X), \cap, \cup)$ are isomorphic: the ‘complement-of-a-subset’ map, (2.25) provides an isomorphism between the two.

3.4.8.2 Linearly ordered sets

A partially ordered set (L, \preceq) is said to be *linearly ordered* if any two elements are comparable, i.e.,

$$l \preceq m \quad \text{or} \quad m \preceq l \quad (l, m \in L).$$

In a linearly ordered set, $l \vee m = \max\{l, m\}$ and $l \wedge m = \min\{l, m\}$.

3.4.8.3 The set of natural numbers ordered by divisibility

Consider the relation of divisibility on the set of natural numbers:

$$l \preceq m \quad \text{if} \quad l \mid m \quad (l, m \in \mathbf{N}). \quad (3.67)$$

Here $l \vee m$ is the *greatest common multiple* of l and m , while $l \wedge m$ is their *greatest common divisor*. Note that natural number 1 is the smallest element while 0 is the largest element, thus (\mathbf{N}, \mid) is an example of a bounded distributive lattice. In fact, lattice (\mathbf{N}, \mid) is complete.

Exercise 85 Show that (\mathbf{N}, \mid) is a complete lattice.

3.4.9 Algebraic structures

3.4.9.1

The general notion of an *algebraic structure* on a set X is usually formulated as a sequence of operations (μ_1, \dots, μ_l) on X satisfying an explicit list of

properties that can be expressed as identities involving any number of those operations and arbitrary elements of set X .

Associativity and commutativity of a single binary operation are examples of such properties, as is left- and right-distributivity of one binary operation over another one.

3.4.9.2

For every operation its place on the list of operations forming the structure does matter. For example, if both μ and ν are binary operations, then (X, μ, ν) is a different structure from (X, ν, μ) unless $\mu = \nu$.

3.4.9.3 The signature of an algebraic structure

We say that an algebraic structure (X, μ_1, \dots, μ_l) has signature (n_1, \dots, n_l) if μ_i is an n_i -ary operation, $1 \leq i \leq l$. The signature is a sequence of natural numbers.

For example, an algebraic structure of signature

$$\underbrace{(0, \dots, 0)}_{l \text{ times}}$$

is the same as a set with a *sequence* of l distinguished points (not all necessarily distinct).

3.4.9.4 Morphisms

A morphism between two structures of the *same* signature,

$$(X, \mu_1, \dots, \mu_l) \rightarrow (Y, \nu_1, \dots, \nu_l),$$

is a map $f: X \rightarrow Y$ such that

$$f: (X, \mu_i) \rightarrow (Y, \nu_i)$$

is a homomorphism for each $1 \leq i \leq l$.

In particular, algebraic structures of a given signature and satisfying a given set of identities, form a (unital) category.

3.4.9.5 Substructures of algebraic structures

We say that (Y, ν_1, \dots, ν_l) is a *substructure* of (X, μ_1, \dots, μ_l) , if each operation μ_i induces operation ν_i on Y , cf. Section 3.2.2.4. This is frequently if not entirely correctly expressed by saying that Y is closed under each μ_i and that ν_i is the restriction of μ_i to Y .

Note that any identities satisfied by operations μ_1, \dots, μ_l and elements of X are automatically satisfied by operations ν_1, \dots, ν_l and elements of Y .

3.4.9.6

The intersection of any family of substructures of an algebraic structure is a substructure itself. Thus, for any subset $A \subseteq X$, there exists the smallest substructure of (X, μ_1, \dots, μ_l) which contains A . We shall denote it $\langle A \rangle$. If $\langle A \rangle = (X, \mu_1, \dots, \mu_l)$, we shall say that subset A *generates* structure (X, μ_1, \dots, μ_l) .

3.4.9.7

Properties of an algebraic structure that ascertain existence of certain elements can often be expressed as identities, if one introduces appropriate operations.

For example, existence of a left identity for a binary operation μ on a set X can be expressed as a 0-ary operation $e: X^0 \rightarrow X$, i.e., a distinguished element $e \in X^7$ such that

$$\mu(e, x) = x \quad (x \in X). \quad (3.68)$$

3.4.9.8

Identity (3.68) can be also expressed as commutativity of the following diagram

$$\begin{array}{ccc} X^0 \times X & \xrightarrow{e \times \text{id}_X} & X \times X \\ \uparrow (1.38) & & \downarrow \mu \\ X & \xlongequal{\quad} & X \end{array} \quad (3.69)$$

where the left vertical arrow is the canonical identification of $X^0 \times X^{\{1\}}$ with $X^{0 \cup \{1\}}$ which itself is identified with X .

⁷We identify a map $X^0 \rightarrow X$ with its single value.

3.4.9.9

Thus, one can define a monoid as a set X equipped with two operations (μ, e) , one binary, the other 0-ary, which satisfy two identities: (3.21), (3.68), and the right analog of (3.68)

$$\mu(x, e) = x \quad (x \in X). \quad (3.70)$$

3.4.9.10

In the similar vain, one can define a group as a set X equipped with three operations (μ, e, ι) , a binary, 0-ary, and unary, which satisfy identities (3.21), (3.68), (3.70), and the identities

$$\mu(\iota(x), x) = x \quad (x \in X) \quad (3.71)$$

and

$$\mu(x, \iota(x)) = x \quad (x \in X), \quad (3.72)$$

the meaning of which should be obvious.

Exercise 86 Express identity (3.71) as commutativity of a certain diagram.

3.4.9.11

Existence of a left sink in a binary structure (X, μ) can be expressed as a 0-ary operation $z: X^0 \rightarrow X$ such that

$$\mu(z, x) = z \quad (x \in X). \quad (3.73)$$

3.4.9.12

Identity (3.73) is expressed also by commutativity of the following diagram

$$\begin{array}{ccc} X^0 \times X & \xrightarrow{z \times \text{id}_X} & X \times X \\ \uparrow (1.38) & & \downarrow \mu \\ X^0 & \xrightarrow{z} & X \end{array} \quad (3.74)$$

3.4.9.13

I will leave it to you to describe semirings with zero and rings as algebraic structures.

3.4.10 Fields

3.4.10.1 Domains

A ring R is a *domain* if the subset of non-zero elements $R \setminus \{0\}$ forms a subsemigroup of R^\times . This is usually expressed by saying that $R \setminus \{0\}$ is closed under multiplication.

3.4.10.2 Division rings

A unital ring R is a *division ring* if $R \setminus \{0\}$ is a group.

3.4.10.3

A commutative division ring is called a *field*.

3.4.10.4

Domains, division rings, fields are all special kinds of rings. They differ from all the previously encountered kinds of algebraic structures: the property of being a domain, a division ring, or a field cannot be described in terms of a certain number of operations on a set which are supposed to obey a certain number of identities involving arbitrary elements of that set.

3.4.11 Sets with an infinitary operation

3.4.11.1

In Section 3.3.2.4 we saw how iterating a commutative and associative binary operation generates I -ary operations

$$\mu_I: X^I \rightarrow X$$

for arbitrary *finite* nonempty index sets I . This way however one cannot construct I -operations for infinite sets of indices.

3.4.11.2

This leads to an idea of defining an *infinitary* operation on a set X as a system of independent I -ary operations

$$\mu_I: X^I \rightarrow X,$$

where I is an arbitrary *nonempty* set no more assumed to be finite, which satisfy the compatibility conditions expressed through commutativity of diagram (3.29) and of the diagram

$$\begin{array}{ccc} X^{I'} & \xrightarrow{\mu_{I'}} & X \\ \rho^* \downarrow & & \uparrow \mu_I \\ X^I & & \end{array} \quad (3.75)$$

where $\rho: I \rightarrow I'$ is an arbitrary bijection between indexing sets, cf. diagram (3.3).

3.4.11.3

Any system of operations $\mu_I: X^I \rightarrow X$ for which diagrams (3.29) and (3.75) commute will be called an *infinitary operation* on a set X .

3.4.11.4

If diagram (3.29) expresses *associativity* of the infinitary operation, then diagram (3.75) expresses its *commutativity*.

3.4.11.5

An infinitary operation on X is not a set since its ‘components’ are indexed by the class of all nonempty sets. One should think of it as a functor from the category of nonempty sets with morphisms being bijections, to the comma category $\text{Set}^{\rightarrow X}$, cf. Section 1.4.8.4,

$$I \mapsto (\mu_I: X^I \rightarrow X) \quad (I \text{ any nonempty set}).$$

3.4.11.6 An infinitary operation with identity

If also μ_\emptyset is present and diagram (3.29) commutes for arbitrary sets, some of which may be empty, then we obtain a definition of an *infinitary operation with identity*. Note that operation μ_\emptyset provides a distinguished element $e \in X$ which indeed is an identity element for the binary operation $\mu_{\{1,2\}}$, cf. Section 3.4.9.8 and commuting diagram (3.69).

3.4.11.7 An infinitary semigroup

A set equipped with an infinitary operation will be called an *infinitary semigroup*. One should think of it as a commutative semigroup whose operation can be performed on arbitrary families of elements.

3.4.11.8 An infinitary monoid

A set equipped with an infinitary operation with identity will be called an *infinitary monoid*.

3.4.11.9 Variants of the definition: restrict the size of I

Limiting oneself to *finite* indexing sets in the definition of an infinitary operation yields a system of operations induced from a single commutative and associative binary operation

$$\mu := \mu_{\{1,2\}}$$

as described in Section 3.3.2.4. Thus we obtain a structure which is equivalent to the structure of a commutative semigroup.

3.4.11.10 σ -operations

Limiting oneself to *countable* indexing sets in the definition of an infinitary operation, yields the definition of a σ -operation on X .

3.4.11.11 A σ -semigroup

A set equipped with a σ -operation will be called σ -semigroup.

3.4.11.12 A σ -monoid

A set equipped with an infinitary operation with identity will be called a σ -monoid.

3.4.12 Examples of infinitary semigroups and monoids

3.4.12.1 The semigroup of maps with values in an infinitary semigroup

The set of maps S^X from a set X to an infinitary semigroup S is naturally an infinitary semigroup. If $\mu_I: S^I \rightarrow S$ is the corresponding I -ary operation

on S , then

$$\nu_I: (S^X)^I \rightarrow S^X \quad (3.76)$$

is the composition of the canonical identifications

$$(S^X)^I \longleftrightarrow S^{X \times I} \longleftrightarrow (S^I)^X$$

with operation μ_I performed ‘pointwise’ at every $x \in X$

$$(S^I)^X \xrightarrow{\prod_{x \in X} \mu_I} S^X.$$

If $e \in S$ is the identity element for S , then the constant map

$$X \rightarrow S, \quad x \mapsto e \quad (x \in X),$$

is the identity element for S^X .

3.4.12.2 A complete lattice

Exercise 87 Let $\mathcal{F} \subseteq \mathcal{P}(S)$ be a family of subsets of a partially ordered set (S, \preceq) such that $\sup F$ exists for every $F \in \mathcal{F}$. Show that the following two sets of upper bounds coincide

$$U\left(\bigcup \mathcal{F}\right) = U\left(\{\sup F \mid F \in \mathcal{F}\}\right). \quad (3.77)$$

Deduce that $\sup \bigcup \mathcal{F}$ exists if and only if $\sup\{\sup F \mid F \in \mathcal{F}\}$ exists and the two are equal

$$\sup\{\sup F \mid F \in \mathcal{F}\} = \sup \bigcup \mathcal{F}. \quad (3.78)$$

3.4.12.3

Given an I -indexed family $(l_i)_{i \in I}$ of elements of a complete lattice (L, \preceq) , let

$$\mu_I((l_i)_{i \in I}) := \sup\{l_i \mid i \in I\}. \quad (3.79)$$

In view of identity (3.78), the system of maps μ_I defined in (3.79) forms an infinitary operation with identity on L . The identity is the supremum of the empty family of elements of L , i.e., the smallest element of L .

3.4.13 Ordered binary structures

3.4.13.1

A binary structure (X, μ) equipped with a partial order \preceq is said to be an *ordered binary structure* if the operation respects the order

$$\mu(x, y) \preceq \mu(x', y') \quad \text{whenever} \quad x \preceq x' \quad \text{and} \quad y \preceq y'.$$

Ordered binary structures naturally form a category.

Exercise 88 Let A and B be subsets of an ordered binary structure X . Suppose that $\sup A$, $\sup B$, and

$$\sup\{\mu(a, b) \mid a \in A, b \in B\}$$

exist. Show that

$$\sup\{\mu(a, b) \mid a \in A, b \in B\} \preceq \mu(\sup A, \sup B). \quad (3.80)$$

3.4.13.2

In multiplicative notation inequality (3.80) becomes

$$\sup AB \preceq (\sup A)(\sup B), \quad (3.81)$$

where

$$AB = \{x \in X \mid x = \mu(a, b) \text{ for some } a \in A \text{ and } b \in B\}, \quad (3.82)$$

cf. Section 3.2.3.6.

3.4.13.3 Distributivity properties of binary structures

We say that an ordered binary structure is *left-distributive* if

$$\sup\{\mu(a, b) \mid b \in B\} = \mu(a, \sup B) \quad (a \in X; \emptyset \neq B \subseteq X) \quad (3.83)$$

whenever $\sup\{\mu(a, b) \mid b \in B\}$ and $\sup B$ exist.

3.4.13.4

In multiplicative notation identity (3.83) becomes

$$\sup aB = a(\sup B) \quad (a \in X). \quad (3.84)$$

Exercise 89 Formulate the definition of a right-distributive ordered binary structure.

3.4.13.5

We shall say that a binary structure is *distributive* if it is left- and right-distributive.

3.4.13.6 An example: $[0, \infty]$

Lemma 3.4.2 Both $([0, \infty], +, \preceq)$ and $([0, \infty], \cdot, \preceq)$ are distributive ordered monoids.

Proof. The obvious inequality

$$\sup B \preceq \sup(a + B) \quad (a \in [0, \infty])$$

shows that $\sup(a + B) = \infty$ if $\sup B = \infty$.

Assuming $\beta = \sup B < \infty$ we note that, for any $0 \preceq \epsilon \preceq \beta$, there exists $b \in B$ such that $\beta - \epsilon \preceq b$, and therefore

$$a + (\beta - \epsilon) \preceq a + b \preceq \sup(a + B).$$

It follows that

$$a + \sup B = \sup\{a + (\beta - \epsilon) \mid 0 \preceq \epsilon \preceq \beta\} \preceq \sup(a + B)$$

which completes the proof of distributivity of $([0, \infty], +, \preceq)$.

Turning to $([0, \infty], \cdot, \preceq)$, we note that, for $a \in (0, \infty)$, inequality (3.81) yields the following inequality

$$\sup B = \sup\left(\frac{1}{a} \cdot aB\right) \preceq \frac{1}{a} \sup aB.$$

After multiplying both sides by a we obtain the desired inequality.

For $a = \infty$, we note that

$$\sup \infty B = \begin{cases} 0 & \text{if } B = \{0\} \\ \infty & \text{if } B \text{ contains } b > 0 \end{cases}$$

and equality (3.84) holds in either case. For $a = 0$ equality (3.84) is trivially satisfied. \square

Exercise 90 Show that the product of any family of distributive ordered binary structures (X_i, μ_i, \preceq_i) is distributive.

3.4.13.7

In particular, the set of maps X^Y from *any* set Y to a distributive ordered binary structure is distributive when equipped with the ‘pointwise’ multiplication and pointwise order.

3.4.13.8

If we limit identity (3.83) to *finite* (respectively, *countable*) nonempty subsets $B \subseteq X$, then we obtain the definition of a *finitely* (respectively, *countably*) left-distributive ordered binary structure.

3.4.13.9

We shall say that a binary structure is *finitely distributive* (respectively, *countably distributive*) if it is finitely left- and right-distributive, (respectively, countably left- and right-distributive).⁸

Lemma 3.4.3 *Suppose an ordered binary structure (X, μ, \preceq) be distributive, and the ordered set (X, \preceq) be a complete lattice. Then inequality (3.80) becomes equality*

$$\sup\{\mu(a, b) \mid a \in A, b \in B\} = \mu(\sup A, \sup B) \quad (\emptyset \neq A, B \subseteq X), \quad (3.85)$$

valid for any pair of nonempty subsets of X .

3.4.13.10

In multiplicative notation identity (3.85) becomes

$$\sup AB = (\sup A)(\sup B) \quad (\emptyset \neq A, B \subseteq X). \quad (3.86)$$

Proof. To ease comprehension, we will employ multiplicative notation in the proof. Note that

$$AB = \bigcup_{a \in A} aB$$

and therefore

$$\sup AB = \sup \left(\bigcup_{a \in A} aB \right) = \sup \{ \sup aB \mid a \in A \} \quad (3.87)$$

⁸We could also refer to distributive binary structures as *completely* distributive, and use the term *distributive* without any adjective to describe *finitely* distributive binary structures. The latter would be in exact agreement with terminology prevalent in Lattice Theory.

by identity (3.78). Denote $\sup B$ by \bar{b} . Complete left-distributivity identity (3.84) yields

$$\{\sup aB \mid a \in A\} = \{a \sup B \mid a \in A\} = \{a\bar{b} \mid a \in A\} = A\bar{b}. \quad (3.88)$$

By combining (3.87) with (3.88) we obtain

$$\sup AB = \sup A\bar{b} = (\sup A)\bar{b} = (\sup A)(\sup B)$$

aided by right-distributivity of (X, μ, \preceq) . \square

3.4.13.11

Note that identity (3.85) holds when $B = \emptyset$ if and only if (X, \preceq) has the smallest element and that element is a left sink of binary structure (X, μ) .

3.4.13.12

By replacing two sets in identity (3.86) by n sets, we obtain the following corollary of Lemma 3.4.3.

Corollary 3.4.4 *Under hypotheses of Lemma 3.4.3 one has*

$$\sup(A_1 \cdots A_n) = (\sup A_1) \cdots (\sup A_n) \quad (\emptyset \neq A_1, \dots, A_n \subseteq X). \quad (3.89)$$

3.4.13.13 A positively ordered binary structure

An ordered binary structure (X, μ, \preceq) is said to be *positively ordered*, if it also satisfies the following inequality

$$x \preceq \mu(x, y) \quad \text{and} \quad y \preceq \mu(x, y) \quad (x, y \in X).$$

3.4.13.14 Example: $[0, \infty]$

Additive monoid $([0, \infty], +)$ is positively ordered while multiplicative monoid $([0, \infty], \cdot)$ is not. The latter, however, contains two positively ordered submonoids: $[1, \infty]$ and $[0, 1]$ except that the latter is positively ordered with respect to the *reverse order* \preceq^{rev} .

Exercise 91 *Show that if e is a left identity element in a positively ordered binary structure (X, μ, \preceq) , then e is the smallest element of X .*

Exercise 92 *Let (X, μ) be a binary structure. Show that $\mathcal{P}(X)$ is a positively ordered binary structure when equipped with the induced multiplication, cf. (3.82) above, and ordered by \subseteq .*

3.4.13.15 A construction of an infinitary operation

We shall now extend the binary operation on a commutative positively ordered semigroup to an infinitary operation if the former is a complete lattice under the partial order. The exposition will be easier to follow if we adopt additive notation for the binary operation and for operations μ_I .

Suppose $(S, +, \preceq)$ is a commutative positively ordered semigroup and assume that (S, \preceq) is a complete lattice. The iterated operations $\sum_{i \in I}$, introduced for finite nonempty sets I in Section 3.3.2.4, satisfy the following formula

$$\sum_{i \in I} s_i = \sup_{I' \subseteq I} \sum_{i \in I'} s_i \quad (I \text{ finite nonempty}). \quad (3.90)$$

Exercise 93 Let $(S, +, \preceq)$ be any commutative positively ordered semigroup (not necessarily a complete lattice when viewed as a partially ordered set). Demonstrate identity (3.90)

3.4.13.16

Formula (3.90) suggests a natural method to extend $\sum_{i \in I}$ from finite to arbitrary nonempty sets of indices:

$$\sum_{i \in I} s_i := \sup_{\substack{I' \subseteq I \\ I' \text{ finite}}} \sum_{i \in I'} s_i. \quad (3.91)$$

Exercise 94 Given two I -indexed families $(s_i)_{i \in I}$ and $(t_i)_{i \in I}$ of elements of S show that

$$\sum_{i \in I} s_i \preceq \sum_{i \in I} t_i \quad (3.92)$$

whenever $s_i \preceq t_i$ for all $i \in I$.

Lemma 3.4.5 Operations $\sum_{i \in I}$ defined in (3.91) satisfy inequality

$$\sum_{l \in L} s_l \preceq \sum_{j \in J} \sum_{i_j \in I_j} s_{i_j} \quad \text{where} \quad L = \coprod_{j \in J} I_j. \quad (3.93)$$

Proof. Any subset L' of the disjoint union of family of indexing sets $(I_j)_{j \in J}$ is the disjoint union

$$L' = \coprod_{j \in J} I'_j$$

of sets $I'_j = \{i_j \in I_j \mid (j, i_j) \in L'\}$ consisting of elements contributed by I_j to L' .

Let $J' \subseteq J$ be the set of $j \in J$ such that $I'_j \neq \emptyset$. If L' is finite, then each I'_j is finite, J' is finite, and

$$\sum_{l \in L'} s_l = \sum_{j \in J'} \sum_{i_j \in I'_j} s_{i_j} \preceq \sum_{j \in J'} \sum_{i_j \in I_j} s_{i_j} \preceq \sum_{j \in J} \sum_{i_j \in I_j} s_{i_j}. \quad (3.94)$$

Inequality (3.94) holds for any finite subset $L' \subseteq L$, hence it holds if we replace the sum over L' by

$$\sum_{l \in L} s_l = \sup_{\substack{L' \subseteq L \\ L' \text{ finite}}} \sum_{l \in L'} s_l.$$

□

Lemma 3.4.6 *Inequality in (3.93) becomes equality if ordered semigroup $(S, +, \preceq)$ is distributive. In particular, operations defined in (3.91) transform S into an infinitary semigroup.*

Proof. Denote $\sum_{l \in L} s_l$ by u . For any $J' \subseteq J$ and any $I'_j \subseteq I_j$ let

$$L' := \coprod_{j \in J'} I'_j.$$

Associativity and commutativity of addition in S mean that, when L' is a finite nonempty set, then

$$\sum_{j \in J'} \sum_{i_j \in I'_j} s_{i_j} = \sum_{l \in L'} s_l \preceq u. \quad (3.95)$$

Let A_j be the set formed by the sums

$$\sum_{i_j \in I'_j} s_{i_j}$$

where I'_j are arbitrary finite nonempty subsets of I_j . Then the set $\sum_{j \in J'} A_j$ is formed by the sums

$$\sum_{j \in J'} \sum_{i_j \in I'_j} s_{i_j}.$$

where finite set J' is fixed and I'_j are arbitrary finite nonempty subsets of I_j .

Distributivity of $(S, +, \preceq)$ yields with help of Corollary 3.4.4

$$\sum_{j \in J'} \sum_{i_j \in I_j} s_{i_j} = \sum_{j \in J'} \sup A_j = \sup \left(\sum_{j \in J'} A_j \right)$$

where

$$\sum_{j \in J'} A_j = \left\{ \sum_{j \in J'} \sum_{i_j \in I'_j} s_{i_j} \mid I'_j \subseteq I_j \text{ finite nonempty} \right\}.$$

Inequality (3.95) implies that

$$\sup \left(\sum_{j \in J'} A_j \right) \preceq u \quad (3.96)$$

which, combined with (3.96), implies

$$\sum_{j \in J'} \sum_{i_j \in I_j} s_{i_j} \preceq u$$

for all finite nonempty $J' \subseteq J$. Passing to the supremum over J' we obtain the inequality

$$\sum_{j \in J} \sum_{i_j \in I'_j} s_{i_j} = \sum_{l \in L} s_l.$$

The reverse inequality holds without any distributivity hypotheses, cf. Lemma 3.4.5. \square

3.4.14 Example: $[0, \infty]^X$

The additive monoid of semiring $[0, \infty]^X$ is distributive and positively ordered. According to Lemma 3.4.6 it becomes an infinitary monoid.

Exercise 95 Show that if $\sum_{i \in I} a_i < \infty$, then the support of family $(a_i)_{i \in I}$,

$$\text{supp}(a_i)_{i \in I} := \{i \in I \mid a_i \neq 0\} \quad (3.97)$$

is countable.

Exercise 96 Show that, for a sequence $(a_n)_{n \in \mathbb{N}}$,

$$\sum_{n \in \mathbb{N}} a_n = \sum_{n=0}^{\infty} a_n. \quad (3.98)$$

In other words, show that the sum of elements of family $(a_n)_{n \in \mathbb{N}}$ coincides with the value of the corresponding infinite series.

3.4.14.1 A variant: construction of a σ -operation

By replacing *arbitrary* sets of indices in (3.91) by *countable* ones we obtain a variant of the previous construction which produces a σ -operation.

Countable distributivity of an ordered semigroup $(S, +, \preceq)$ guarantees that identities S becomes a σ -semigroup. The proof of this fact is word by word the same

3.4.15 Infinitary semirings

3.4.15.1

3.5 Sets with an action

3.5.1 Sets with an action of another set

3.5.1.1

We say that a set G *acts* on a set X if we associate with every element $g \in G$, a selfmap $\lambda_g: X \rightarrow X$. The family of selfmaps $(\lambda_g)_{g \in G}$ is a map

$$\lambda: G \rightarrow X^X, \quad g \mapsto \lambda_g \quad (g \in G). \quad (3.99)$$

3.5.1.2

The action of G on X can be also given in the form of a pairing

$$\tilde{\lambda}: G \times X \rightarrow X, \quad (3.100)$$

where $\tilde{\lambda}$ and λ are linked by the identity

$$\tilde{\lambda}(g, x) = \lambda_g(x) \quad (g \in G; x \in X). \quad (3.101)$$

Using identity (3.101) one can recover λ from $\tilde{\lambda}$.

3.5.1.3 Simplified notation

A common practice is to denote $\lambda_g(x)$ by gx as if we are multiplying x by g on the left.

3.5.1.4 The category of G -sets

Sets equipped with an action of a given set G naturally form a category. Morphisms $(X, \lambda) \rightarrow (Y, \mu)$ are maps $f: X \rightarrow Y$ which are compatible with G -action. This translates into f satisfying the equalities

$$f \circ \lambda_g = \mu_g \circ f \quad (g \in G). \quad (3.102)$$

Explicitly,

$$f(\lambda_g(x)) = \mu_g(f(x)) \quad (g \in G; x \in X) \quad (3.103)$$

or, in simplified notation,

$$f(gx) = gf(x) \quad (g \in G; x \in X).$$

The category of G -sets will be denoted $G\text{-Set}$.

3.5.1.5

In the language of commuting diagrams identity (3.103) is equivalent to commutativity of the square diagram

$$\begin{array}{ccc} G \times X & \xrightarrow{\tilde{\lambda}} & X \\ \text{id}_G \times f \downarrow & & \downarrow f \\ G \times Y & \xrightarrow{\tilde{\mu}} & Y \end{array} \quad (3.104)$$

3.5.1.6 Equivariant maps

Traditionally, morphisms in the category of G -sets are called *equivariant maps*.

3.5.2 Objects with an action of a set

3.5.2.1

This is an obvious generalization of the previous structure. We say that a set G *acts* on an object a of a category \mathcal{C} if we associate with every element $g \in G$, an endomorphism $\lambda_g: a \rightarrow a$. The family of endomorphisms $(\lambda_g)_{g \in G}$ is a map

$$\lambda: G \rightarrow \text{End}_{\mathcal{C}}(a), \quad g \mapsto \lambda_g \quad (g \in G). \quad (3.105)$$

3.5.2.2 The category of G -objects

Objects of a category \mathcal{C} equipped with an action of a given set G form a category. Morphisms $(a, \lambda) \rightarrow (b, \mu)$ are morphisms $\alpha: a \rightarrow b$ which are compatible with G -action. This translates into α satisfying the identity

$$\alpha \circ \lambda_g = \mu_g \circ \alpha \quad (g \in G) \quad (3.106)$$

The category of G -objects for a category \mathcal{C} will be denoted $G\text{-}\mathcal{C}$.

3.5.3 Sets with an action of a semigroup

3.5.3.1

When a set G that acts on a set X , is equipped with a binary operation, it is natural to require that the operation and the action are compatible. This translates into saying that map (3.99) should be a homomorphism,

$$\lambda_{gh} = \lambda_g \circ \lambda_h \quad (g, h \in G) \quad (3.107)$$

or, using simplified notation, that the identity

$$(gh)x = g(hx) \quad (g, h \in G; x \in X). \quad (3.108)$$

3.5.3.2

Noting that (3.108) closely resembles the associativity condition, it is not surprising that this definition is particularly well suited to the case when multiplication in G is associative, i.e., when G is a semigroup.

3.5.3.3

In the context of semigroups, the phrase ‘a G -set’ always means

$$\text{a set equipped with an action of set } G \text{ such that the structural map, (3.99), is a homomorphism of semigroups.} \quad (3.109)$$

In this restricted sense, G -sets form a full subcategory of the category of G -sets where G is simply considered to be a set.

3.5.3.4

Similarly, we say that a semigroup G acts on an object a of a category \mathcal{C} , if a homomorphism of semigroups (3.105) is given.

3.5.3.5 Notation

Notation $G\text{-Set}$ and $G\text{-}\mathcal{C}$ will be used to denote the corresponding categories of G -sets and G -objects.

3.5.3.6 The case of a monoid

We have seen above that a homomorphism of semigroups does not preserve the identity elements, in general. A homomorphism of monoids is explicitly required to preserve the identity elements.

Note that X^X and $\text{End}_{\mathcal{C}}(a)$ are monoids. If G is a monoid G , we require that the structural maps, (3.99) and (3.105) are homomorphisms of monoids. In other words, we require them to be homomorphisms of the corresponding binary operations and, additionally to preserve the identity:

$$\lambda_e = \text{id}_X \quad (\text{in the case of an action on a set } X) \quad (3.110)$$

and

$$\lambda_e = \text{id}_a \quad (\text{in the case of an action on an object } a). \quad (3.111)$$

3.5.3.7 Group actions

This case is of particular importance. The role played by groups in Mathematics and its applications to Physics, Chemistry, and Engineering, is primarily as groups of symmetries of various objects.

3.5.4 Semimodules

3.5.4.1

The set of endomorphisms $\text{End}_{\text{Semigrp}}(A)$ of a commutative semigroup A is a unital semiring, cf. Section 3.4.3.7. We will say that a semiring R acts on a commutative group A if a homomorphism of semirings

$$\lambda: R \rightarrow \text{End}_{\text{Semigrp}}(A) \quad (3.112)$$

is given.

3.5.4.2 The action of a semiring analyzed

Let us translate into concrete identities the fact that (3.112) defines an action of semiring R on semigroup A . We will be using simplified notation throughout: $ra := \lambda_r(a)$.

3.5.4.3

Let us begin from the fact that, for each $r \in R$, map λ_r is supposed to be an endomorphism of semigroup A . This is expressed by the identity

$$r(a + b) = ra + rb \quad (r \in R; a, b \in A). \quad (3.113)$$

3.5.4.4

Map (3.112) is supposed to be a homomorphism of the additive semigroup R^+ of R into the additive semigroup of semiring $\text{End}_{\text{Semigrp}}(A)$. This is expressed by the identity

$$(r + s)a = ra + sa \quad (r, s \in R; a \in A). \quad (3.114)$$

3.5.4.5

Finally, map (3.112) is supposed to be also a homomorphism of the *multiplicative* semigroup R^\times of R into the multiplicative semigroup of semiring $\text{End}_{\text{Semigrp}}(A)$. This is expressed by the identity

$$(rs)a = r(sa) \quad (r, s \in R; a \in A). \quad (3.115)$$

3.5.4.6

If we disregard the fact that the left multiplier in the expression ra belongs to R while the right multiplier belongs to A , then we can interpret identity (3.113) as left-distributivity of multiplication by elements of R over addition in A .

Similarly, identity (3.114) can be interpreted as right-distributivity of multiplication by elements of A over addition in R .

Finally, identity (3.115) looks like associativity of multiplication, except that the left-hand-side of (3.115) involves two different ‘multiplications’: of two elements of R , and of an element of R and an element of A .

3.5.4.7 R -semimodules

A short name for a semiring R acting on a commutative semigroup is an R -*semimodule* or, to be precise, a *left* R -semimodule—since there is a version of the semimodule definition in which R acts from the right. An alternative way to say the same: a *semimodule over* R .

3.5.4.8 The category of R -semimodules

Given two R -semimodules, a morphism $A \rightarrow B$ is a morphism of semigroups $f: A \rightarrow B$, i.e., an additive map, which is compatible with actions of R on A and B . This last requirement is expressed as the identity

$$f(ra) = rf(a) \quad (r \in R; a \in A). \quad (3.116)$$

Maps between commutative semigroups $f: A \rightarrow B$ which satisfy identity (3.116) are said to be *homogeneous* (of degree 1).

Thus, morphisms between R -semimodules are maps that are additive and homogeneous of degree 1. Maps with these two properties are also called R -linear or, simply, linear, when the semiring of coefficients is clear from the context.

3.5.4.9 Terminology

Given an R -semimodule A , if we need to refer to the underlying structure of a semigroup forgetting the action of R , then we call it the *additive semigroup* of A .

If we need to refer to R , we call it the *semiring of coefficients*, or the *ground semiring*.

3.5.4.10 Subsemimodules

Let us look at the additive semigroup A^+ of an R -semimodule A . Suppose that a subsemigroup B be of A^+ satisfies the property

$$rb \in B \text{ for any } r \in R; \text{ and } b \in B. \quad (3.117)$$

Then one can consider B , equipped with the R -action induced from A , as an R -semimodule. Such a semimodule is called a *subsemimodule* of A .

3.5.4.11

Exercise 97 For any subset X of an R -semimodule A , let RX be the set formed by sums in A

$$\xi = \sum_{x \in X'} r_x x \quad (3.118)$$

where X' is any finite nonempty subset of X and $(r_x)_{x \in X'}$ is any family of elements of R , indexed by set X' . Show that RX is a subsemimodule of A .

3.5.4.12

Note that $R\emptyset = \emptyset$.

Exercise 98 Show that the intersection of any family $(B_i)_{i \in I}$ of subsemimodules of A is a subsemimodule.

Exercise 99 Show that RX coincides with the intersection of the family of subsemimodules of A which contain subset X .

3.5.4.13 The subsemimodule generated by a subset

Subsemimodule RX is the smallest subsemimodule of A which contains subset X . We shall refer to it as the subsemimodule *generated* by $X \subseteq A$.

3.5.4.14 Sets of generators

If $RX = A$, we say that X is a *set of generators* for R -semimodule A .

Exercise 100 Suppose that $z \in R$ is a right zero, i.e.,

$$r + z = r \quad \text{and} \quad rz = z \quad (r \in R).$$

Show that the set

$$zA := \{za \mid a \in A\} \tag{3.119}$$

is a subsemimodule of A and every element in zA is an additive idempotent

$$b + b = b \quad (b \in zA).$$

3.5.4.15 Semimodules over a semiring with zero

When A is a commutative monoid, then $\text{End}_{\text{Mon}}(A)$ is a semiring with zero: the constant map $A \rightarrow A$ which sends every element of A to $0 \in A$ playing the role of the zero element.

If the ground semiring itself has zero, then in the definition of a semimodule we additionally request that $0 \in R$ acts on elements of A via the zero map

$$0_R \cdot a = 0_A \quad (a \in A)$$

or, in simplified notation,

$$0a = 0 \quad (a \in A).$$

3.5.4.16 The category of semimodules over a semiring with zero

This is a subcategory of the category of semimodules whose objects are commutative monoids instead of commutative semigroups, and morphisms are supposed to be homomorphisms of commutative monoids, i.e., be additive maps and additionally send 0 to 0.

This is an example of a not full subcategory.

3.5.4.17 Unitary semimodules

Suppose that the ground ring is unital. If $1 \in R$ acts on A as the identity endomorphism, then A is said to be a *unitary* R -semimodule.

3.5.4.18 Example: unitary \mathbf{Z}_+ -semimodules

Consider the set of positive integers

$$\mathbf{Z}_+ := \{1, 2, \dots\} \quad (3.120)$$

equipped with usual addition and multiplication. It is a unital semiring.

For any unitary semimodule over \mathbf{Z}_+ , one has

$$na = \underbrace{(1 + \dots + 1)}_{n \text{ times}} a = \underbrace{a + \dots + a}_{n \text{ times}} \quad (a \in A) \quad (3.121)$$

which means that a structure of a unitary \mathbf{Z}_+ -semimodule on a semigroup A is completely determined by the structure of A as a semigroup. In particular, there is only one structure of a unitary \mathbf{Z}_+ -semimodule on any given commutative semigroup.

Vice-versa, for any commutative semigroup A , formula (3.121) defines an action of the semiring of positive integers on A making it a unitary \mathbf{Z}_+ -semimodule.

Exercise 101 *Show that any homomorphism of commutative semigroups is automatically a homomorphism of \mathbf{Z}_+ -semimodules.*

3.5.4.19

It follows that the category of unitary \mathbf{Z}_+ -semimodules is isomorphic to the category of commutative semigroups.

3.5.4.20 Example: unitary \mathbf{N} -semimodules with zero

Consider the set of natural numbers

$$\mathbf{N} := \{0, 1, 2, \dots\} \quad (3.122)$$

equipped with usual addition and multiplication. It is a unital semiring with zero which contains \mathbf{Z}_+ as a subsemiring.

Let A be a unitary \mathbf{N} -semimodule with zero. Thus, $0 \in \mathbf{N}$ acts by sending any element $a \in A$ to $0 \in A$ while any positive integer acts on A by formula (3.121). Accordingly, a structure of a unitary \mathbf{N} -semimodule with zero on a monoid A is completely determined by the structure of A as a monoid. In particular, there is only one structure of a unitary \mathbf{N} -semimodule on any given commutative monoid.

Vice-versa, for any commutative monoid A , formulae (3.121) and

$$0a = 0$$

define an action of the semiring of natural numbers on A making it a unitary \mathbf{N} -semimodule with zero.

Any homomorphism of commutative monoids is automatically a homomorphism of \mathbf{N} -semimodules. It follows that the category of unitary \mathbf{N} -semimodules with zero is isomorphic to the category of commutative monoids.

3.5.4.21 Modules over a ring

When A is an abelian group, then $\text{End}_{\text{Grp}}(A)$ is a unital ring. If the ground semiring is a ring, then any homomorphism of semirings (3.112) is automatically also a homomorphism of rings (rings form a full subcategory in the category of semirings). In this case we speak of R -modules, or modules over R , rather than semimodules.

3.5.4.22 The category of R -modules

The category of R -modules is a full subcategory of the category of R -semimodules.

3.5.4.23 The category of unitary R -modules

The most frequently encountered is the category of unitary modules over a unital ring R . It is this category that is usually denoted $R\text{-mod}$.

3.5.4.24 Vector spaces

Unitary modules over a field F are called F -vector spaces or vector spaces over F .

3.5.4.25

As we noted above, a semimodule A over a ring R is a module precisely when the additive semigroup of A is a group. When the ground ring is unital and A is a unitary R -semimodule, then the additive semigroup of A is forced to be a group.

Observation 3.5.1 *Any unitary semimodule is automatically a module. More precisely, for any $a \in A$, the element $(-1)a$ is the additive inverse to a .*

Indeed, for any $a \in A$, one has

$$0 = 0a = (1 + (-1))a = 1a + (-1)a = a + (-1)a,$$

i.e., $(-1)a$ is the right inverse of a in the additive semigroup of A (it is automatically a two-sided inverse since addition in A is commutative).

Exercise 102 *Let A be an R -semimodule over a nonunital ring R . Show that, for any $a \in A$, the following subset of A*

$$Ra := \{ra \mid r \in R\}$$

is a subgroup of the additive group of A .

3.5.4.26 Example: unitary \mathbf{Z} -modules

Consider the set of integers

$$\mathbf{N} := \{0, \pm 1, \pm 2, \dots\} \tag{3.123}$$

equipped with usual addition and multiplication. It is a unital ring which contains \mathbf{N} as a subsemiring with zero.

Let A be a unitary \mathbf{Z} -module. Action of positive integers is governed by identity (3.121).

3.5.5 Semialgebras

3.5.5.1

We defined semirings as commutative semigroups equipped with a biadditive binary operation. It happens very often that the semigroup is a semimodule over certain semiring, and that the operation is *bilinear*.

3.5.5.2 Bilinear pairings

Suppose that R -semimodules A , B , and C be given. A biadditive pairing

$$\mu: A \times B \rightarrow C \quad (3.124)$$

is said to be R -*bilinear*, or a R -*biadditive pairing*, if it is homogeneous of (degree 1) in each argument:

$$\mu(ra, b) = r\mu(a, b) \quad (r \in R; a \in A; b \in B), \quad (3.125)$$

and

$$\mu(a, rb) = r\mu(a, b) \quad (r \in R; a \in A; b \in B). \quad (3.126)$$

3.5.5.3

The notion of of a bilinear pairing makes sense for semimodules over any semiring. When R is not commutative, however, its usefulness is limited. A proper context for ‘bilinear’ and, more generally, multilinear maps requires replacing semimodules by *semibimodules*. This will not be discussed here, so from now on we shall assume that the ground ring is *commutative*.

3.5.5.4 Semialgebras: terminology and notation

A semimodule equipped with a bilinear multiplication

$$\mu: A \times A \rightarrow A$$

is called a *semialgebra*. There is a tradition to denote by k the ground semiring which, as you remember, is assumed to be commutative. If one needs to be more specific, terms like ‘a k -semialgebra’ or ‘a semialgebra over k ’ are used as well.

3.5.5.5

All the terms applicable to semirings continue to be applicable to semialgebras: ‘associative’, ‘commutative’, ‘with zero’, ‘unital’, etc.

3.5.5.6 Example: semirings as \mathbf{Z}_+ -semialgebras

Every semiring is automatically a semialgebra over \mathbf{Z}_+ , cf. The category of semirings and the category of \mathbf{Z}_+ -semialgebras are isomorphic.

3.5.5.7 Example: k^X

The set of maps $X \rightarrow k$ with values in a commutative semiring, with pointwise addition and multiplication is naturally a k -semialgebra.

3.5.5.8 Example: semirings with zero as \mathbf{N} -semialgebras

Every semiring with zero is automatically a semialgebra over \mathbf{N} , cf. The category of semirings with zero and the category of \mathbf{N} -semialgebras are isomorphic.

3.5.5.9 Algebras

‘Semialgebras’ are called *algebras*, if k is a ring and A is a k -module.

3.5.5.10 Morphisms

Morphisms between semialgebras $A \rightarrow B$ are maps $f: A \rightarrow B$ which are simultaneously homomorphisms of the corresponding additive semimodules and of the multiplicative binary structures. Like for other algebraic structures, they are usually called *homomorphisms*.

3.5.5.11

Morphisms $A \rightarrow B$ between semialgebras with zero are of course expected to send 0_A to 0_B .

3.5.5.12

We said it already twice before: associativity is of such importance that it became a standard practice to tacitly assume associativity when talking about semialgebras and algebras.

The category of associative semialgebras over k will be denoted $k\text{-semialg}$, and the category of associative algebras will be denoted $k\text{-alg}$.

3.5.5.13 Terminology: a warning

The term ‘algebra’ is used in Mathematics in at least two different ways: as a special kind of algebraic structure, and as a branch of Mathematics. In the latter sense I suggest to always capitalize it: Algebra.

Term ‘algebra’ can be also used in a loose sense of anything that involves extensive manipulations of symbolic expressions.

You have to be aware that various structures, not necessarily strictly algebraic, were designated with term ‘algebra’ before the latter became attached to that particular algebraic structure we call now *an algebra*.

For example, neither Boolean algebras nor σ -algebras are algebras in the sense given above. Both, however, are semirings of special kind.

3.5.5.14 Example: $\mathcal{P}(X)$ as an \mathbf{F}_2 -algebra

The set of all subsets of any set X is only a semiring when considered with operations \cup and \cap . If one replaces *union* by *disjoint union*,

$$A \mid B := (A \cup B) \setminus (A \cap B), \quad (3.127)$$

then $(\mathcal{P}(X), \mid, \cup)$ becomes a commutative unital algebra over the field with two elements $\mathbf{F}_2 = \{0, 1\}$.

Exercise 103 Show that \cap distributes over \mid .

Chapter 4

Universal constructions

4.1 Universal properties

4.1.1 Initial and final objects

4.1.1.1

An object u in a category \mathcal{C} is *initial* if, for any object c , there is a unique morphism $u \rightarrow c$.

Exercise 104 *Show that any two initial objects u and u' are isomorphic by a unique isomorphism.*

4.1.1.2

Dually, an object v in a category \mathcal{C} is *final* if, for any object c , there is a unique morphism $c \rightarrow v$. Note that v is final precisely when v^{op} is initial in the opposite category, \mathcal{C}^{op} .

4.1.1.3

The empty set is a unique initial object in the usual category of sets, Set , while single-element sets are the final objects.

4.1.1.4

The empty set is simultaneously a unique initial and a unique final object in the category of sets and multimaps, Set_{mult} , and in the isomorphic category of sets and binary relations, Set_{rel} .

4.1.1.5

If we view a partially ordered set as a category, then the least element is an initial object, and the greatest element is a final object. Note that in this case an initial, as well as a final, objects are unique when they exist.

4.1.1.6

In the category of groups, Grp , a one-element group is simultaneously an initial and a final object.

4.1.1.7

In the category of unital rings, Ring_1 , the ring of integers, \mathbb{Z} , is an initial object.

Exercise 105 Show that there is no final object in Ring_1 .¹ (Hint. Show that for any unital ring R there are at least two different ring homomorphisms $R \times R \rightarrow R$.)

4.1.2 The quotient of a set by a relation

4.1.2.1

Let $R \subseteq E \times E$ be a binary relation on a set E . Let us consider the category whose objects are maps $f: E \rightarrow X$, where X denotes an arbitrary set, such that

$$f(a) = f(b) \quad \text{whenever} \quad a \sim_R b \quad (a, b \in E).$$

Morphisms from $E \xrightarrow{f} X$ to $E \xrightarrow{g} Y$ are maps $\phi: X \rightarrow Y$ such that $g = \phi \circ f$ which may be expressed by commutativity of the diagram

$$\begin{array}{ccc} E & \xrightarrow{g} & Y \\ & \searrow f & \nearrow \phi \\ & X & \end{array} .$$

4.1.3 Product

4.1.3.1

Let

$$(a_i)_{i \in I} \tag{4.1}$$

¹Note: a ring is *unital* if it possesses a nonzero identity element.

be a family of objects in a category \mathcal{C} . Given a morphism $\alpha: x \rightarrow y$, and a family of morphisms $g_i: y \rightarrow a_i$, we can form the family of morphisms

$$f_i = g_i \circ \alpha \quad (i \in I). \quad (4.2)$$

The family $(f_i)_{i \in I}$ is said to be *induced from family $(g_i)_{i \in I}$* by morphism α .

4.1.3.2 A universal family

We say that an object $p \in \text{Ob } \mathcal{C}$ equipped with a family of morphisms $(\pi_i: p \rightarrow a_i)_{i \in I}$, is a *product* of family (4.1) if any family of morphisms (4.2) can be induced from family $(\pi_i)_{i \in I}$ by a *unique* morphism $\alpha: x \rightarrow p$.

Exercise 106 Show that if $(\pi_i: p \rightarrow a_i)_{i \in I}$ and $(\pi'_i: p' \rightarrow a_i)_{i \in I}$ are products of family $(a_i)_{i \in I}$, then p and p' they are isomorphic.

4.1.3.3 Terminology and notation

Morphisms $\pi_i: p \rightarrow a_i$ are referred to as the *canonical projections*. Even though the term ‘product’ is often applied just to object p , the canonical projections form a part of the product structure.

4.1.3.4 Example: the category of fields

A product may fail to exist. This happens, in particular, when for a given family of objects (4.1) there is no object $x \in \text{Ob } \mathcal{C}$ such that

$$\text{Hom}_{\mathcal{C}}(x, a_i) \neq \emptyset.$$

This situation may occur, e.g., in the category of fields where any morphism $E \rightarrow F$ is an injective map between the underlying sets. Thus, a product of \mathbf{F}_2 and $\mathbf{F}_3 = \{0, 1, -1\}$ does not exist.

Exercise 107 Show that \mathbf{F}_2 is a product of \mathbf{F}_2 and \mathbf{F}_2 in the category of fields, with the ‘canonical projections’ being the identity maps $\mathbf{F}_2 \rightarrow \mathbf{F}_2$.

4.1.3.5 Example: a partially ordered set viewed as a category

Here, a product of family (4.1) exists if and only if the set

$$\{a_i \mid i \in I\}$$

has infimum. In this case product is *unique*, namely

$$\inf \{a_i \mid i \in I\}.$$

Exercise 108 Prove the above two assertions.

4.1.3.6

When a product of (4.1) exists it is not unique if there exists $p' \in \text{Ob } \mathcal{C}$ and a non-identity isomorphism $p \simeq p'$. However, any two solutions to the problem of existence of a product of family $(a_i)_{i \in I}$ are isomorphic, and there exists only one such isomorphism which is compatible with all the projection morphisms.

4.1.3.7 Functorial products

When a product exists for any family of objects in a category \mathcal{C} , it frequently happens, that there is a *functorial* solution to the problem of existence of a product. This means that there exists a functor, denoted \prod , from the category of I -indexed families $(a_i)_{i \in I}$ of objects in \mathcal{C} to the category of I -indexed families of morphisms $(f_i: x \rightarrow a_i)_{i \in I}$ whose ‘values’ are products of the corresponding families.

There is no need to say more about it now. We will signal such functorial constructions of products when we encounter them.

4.1.3.8 Product of a family of sets

For a family of sets $(X_i)_{i \in I}$, let

$$X = \bigcup_{i \in I} X_i,$$

and let

$$\prod_{i \in I} X_i := \{\zeta: I \rightarrow X \mid \zeta(i) \in X_i\}. \quad (4.3)$$

Set (4.3) is called the Cartesian product of family $(X_i)_{i \in I}$. Usual interpretation of elements of the Cartesian product is as families $(x_i)_{i \in I}$ of elements of X such that $x_i \in X_i$. The maps that forget all but one component of $(x_i)_{i \in I}$ are the canonical projections.

4.1.3.9

Given any family of maps $f_i: W \rightarrow X_i$, define

$$\tilde{f}: W \rightarrow \prod_{i \in I} X_i \quad w \mapsto (f_i(w))_{i \in I}.$$

4.1.3.10 Functoriality of the Cartesian product

A morphism $(X_i)_{i \in I} \rightarrow (Y_i)_{i \in I}$ is a family of maps $\{f_i: X_i \rightarrow Y_i\}$. It induces the map

$$\prod_{i \in I} f_i: \prod_{i \in I} X_i \rightarrow \prod_{i \in I} Y_i, \quad (x_i)_{i \in I} \mapsto (f_i(x_i))_{i \in I}. \quad (4.4)$$

which is compatible with the composition of morphisms of families of sets.

4.1.3.11 Alternative notation

If I is a finite linearly ordered set like $\{1, \dots, n\}$, an alternative notation is frequently used

$$X_1 \times \cdots \times X_n \quad \text{and} \quad f_1 \times \cdots \times f_n$$

instead of

$$\prod_{i=1}^n X_i \quad \text{and} \quad \prod_{i=1}^n f_i.$$

4.1.3.12 Product of a family of algebraic structures

Consider a family of algebraic structures

$$((X_i, \mu_{i1}, \dots, \mu_{il}))_{i \in I}$$

of signature (n_1, \dots, n_l) . We shall equip the Cartesian product of sets $\prod_{i \in I} X_i$ with the structure of the same signature by applying corresponding operations *componentwise*:

$$v_j((x_i)_{i \in I}, (y_i)_{i \in I}) := (\mu_{ij}(x_i, y_i))_{i \in I} \quad (j = 1, \dots, l).$$

This construction depends functorially on the family of structures and one can easily verify that the

4.1.3.13 Product of a family of topological spaces

4.1.3.14 Product of a family of measurable spaces