

Turing Degrees and Definability of the Jump

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Outline

- ▶ Lecture 1
 - ▶ Forcing in arithmetic
 - ▶ Coding and decoding theorems
 - ▶ Automorphisms of countable ideals
 - ▶ Persistence
- ▶ Lecture 2
 - ▶ Absoluteness
 - ▶ Generic Persistence
 - ▶ Definable representations of automorphisms

Persistent Automorphisms

review

Definition

An automorphism ρ of a countable ideal \mathcal{I} is **persistent** if for every degree x there is a countable ideal \mathcal{I}_1 such that

- ▶ $x \in \mathcal{I}_1$ and $\mathcal{I} \subseteq \mathcal{I}_1$;
- ▶ there is an automorphism ρ_1 of \mathcal{I}_1 such that the restriction of ρ_1 to \mathcal{I} is equal to ρ .

We have shown that persistent automorphisms are locally presented and that persistent automorphisms of ideals which contain $0'$ have persistent extensions. Thus, given a persistent automorphism $\rho : \mathcal{I} \xrightarrow{\sim} \mathcal{I}$, a countable jump ideal \mathcal{J} extending \mathcal{I} , and a presentation J of \mathcal{J} , we can find an extension of ρ to \mathcal{J} which is arithmetic in J .

Persistent Automorphisms

absoluteness

Theorem

The property I is a representation of a countable ideal \mathcal{I} , $0' \in \mathcal{I}$, and R is a presentation of a persistent automorphism ρ of \mathcal{I} is Π_1^1 .

Proof

ρ is persistent if and only if for every presentation J of a jump ideal \mathcal{J} extending \mathcal{I} , there is an arithmetic in J extension of ρ to \mathcal{J} . This property is Π_1^1 . \square

Corollary

The properties R is a presentation of a persistent automorphism and There is a countable map $\rho: \mathcal{I} \xrightarrow{\sim} \mathcal{I}$ such that $0' \in \mathcal{I}$, ρ is persistent and not equal to the identity are absolute between well-founded models of ZFC.

Persistent Automorphisms

model theoretically

Let T be the fragment of ZFC in which we include only the instances of replacement and comprehension in which the defining formula is Σ_1 .

Definition

Suppose that $\mathcal{M} = (M, \in^{\mathcal{M}})$ is a model of T .

1. \mathcal{M} is an **ω -model** if $\mathbb{N}^{\mathcal{M}}$ is isomorphic to the standard model of arithmetic.
2. \mathcal{M} is **well-founded** if the binary relation $\in^{\mathcal{M}}$ is well-founded. That is to say that there is no infinite sequence $(m_i : i \in \mathbb{N})$ of elements of \mathcal{M} such that for all i , $m_{i+1} \in^{\mathcal{M}} m_i$.

Persistent Automorphisms

model theoretically

Theorem

Suppose that \mathcal{M} is an ω -model of T . Let \mathcal{I} be an element of \mathcal{M} such that

$\mathcal{M} \models \mathcal{I}$ is a countable ideal in \mathfrak{D} such that $0' \in \mathcal{I}$.

Then, every persistent automorphism of \mathcal{I} is also an element of \mathcal{M} .

Proof

\mathfrak{M} is closed under arithmetic definability. □

Persistent Automorphisms

model theoretically

Corollary

Suppose that \mathcal{M} is an ω -model of T and that ρ and \mathcal{I} are elements of \mathcal{M} such that $0' \in \mathcal{I}$, $\rho : \mathcal{I} \xrightarrow{\sim} \mathcal{I}$, and \mathcal{I} is countable in \mathcal{M} . Then,

$$\rho \text{ is persistent} \implies \mathcal{M} \models \rho \text{ is persistent.}$$

Proof

Persistent automorphisms extend persistently. Hence, their extensions belong to \mathfrak{M} . □

Generic Persistence

We now extend the notion of persistence to uncountable ideals. In what follows, V is the universe of sets and G is a V -generic filter for some partial order in V .

Definition

Suppose that \mathcal{I} is an ideal in \mathfrak{D} and ρ is an automorphism of \mathcal{I} . We say that ρ is **generically persistent** if there is a generic extension $V[G]$ of V in which \mathcal{I} is countable and ρ is persistent.

Generic Persistence

invariance under partial order

Theorem

Suppose that $\rho : \mathcal{I} \xrightarrow{\sim} \mathcal{I}$ is generically persistent. If $V[G]$ is a generic extension of V in which \mathcal{I} is countable then ρ is persistent in $V[G]$.

Proof

Generics for any two forcings can be realized simultaneously.
Apply absoluteness. □

Applications to $\text{Aut}(\mathcal{D})$

Theorem

Suppose that $\pi : \mathcal{D} \xrightarrow{\sim} \mathcal{D}$. Then, π is generically persistent.

Proof

If not, then the failure of π to be generically persistent would reflect to a countable well-founded model \mathfrak{M} . Then, we could add a generic counting of $\mathcal{D}^{\mathfrak{M}}$ to \mathfrak{M} to obtain $\mathfrak{M}[G]$ in which $\pi \upharpoonright \mathcal{D}^{\mathfrak{M}}$ is not persistent. Contradiction to the persistence of $\pi \upharpoonright \mathcal{D}^{\mathfrak{M}}$ and the absoluteness of persistence. \square

Applications to $\text{Aut}(\mathfrak{D})$

definability of automorphisms

Theorem

Suppose that $V[G]$ is a generic extension of V . Suppose that π is an element of $V[G]$ which maps Turing degrees in V automorphically to itself (that is, $\pi : \mathfrak{D}^V \xrightarrow{\sim} \mathfrak{D}^V$). If π is generically persistent in $V[G]$, then π is an element of $L(\mathbb{R}^V)$. That is, π is constructible from the set of reals in V .

Proof

π is generically persistent, so π is arithmetically definable in any generic counting of \mathfrak{D}^V . Consequently, π must belong to the ground model for such countings, namely $L(\mathbb{R}^V)$. \square

Applications to $\text{Aut}(\mathfrak{D})$

global extension of persistent automorphisms

Theorem

Suppose that $0'$ is an element of \mathcal{I} and $\rho : \mathcal{I} \xrightarrow{\sim} \mathcal{I}$ is persistent. Then ρ can be extended to a global automorphism $\pi : \mathfrak{D} \xrightarrow{\sim} \mathfrak{D}$.

Proof

ρ can be persistently extended to \mathfrak{D}^V in a generic extension of V . This extension belongs to $L[\mathbb{R}^V]$. \square

Corollary

*The statement **There is a non-trivial automorphism of the Turing degrees** is equivalent to a Σ_2^1 statement. It is therefore absolute between well-founded models of ZFC.*

Applications to $\text{Aut}(\mathcal{D})$

lifting automorphisms to generic extensions

Theorem

Let π be an automorphism of \mathcal{D} . Suppose that $V[G]$ is a generic extension of V . Then, there is an extension of π in $V[G]$ to an automorphism of $\mathcal{D}^{V[G]}$, the Turing degrees in $V[G]$.

Proof

There is a persistent extension π_1 of π in any generic extension of $V[G]$ in which $\mathcal{D}^{V[G]}$ is countable. This π_1 belongs to $V[G]$. □

Representing Automorphisms

Definition

Given two functions $\tau : \mathcal{D} \rightarrow \mathcal{D}$ and $t : 2^\omega \rightarrow 2^\omega$, we say that t **represents** τ if for every degree x and every set X in x , the Turing degree of $t(X)$ is equal $\tau(x)$.

We will analyze the behavior of an automorphism of \mathcal{D} in terms of the action of its extensions on the degrees of the generic reals.

Representing Automorphisms

continuity on sufficiently generic sets

Theorem

Suppose that $\pi : \mathcal{D} \xrightarrow{\sim} \mathcal{D}$. There is a countable family \vec{D} of dense open subsets of $2^{<\omega}$ such that π is represented by a continuous function f on the set \vec{D} -generic reals.

The proof has several steps, which we will sketch. Also, we introduce the notation $\Pi(Z)$ for a representative of $\pi(\text{degree}(Z))$.

1. Let $V[\mathcal{G}]$ be a generic extension of V obtained by adding ω_1 -many Cohen reals and let π_1 be an extension of π to $\mathcal{D}^{V[\mathcal{G}]}$.
2. Since $\pi_1 \in L[\mathbb{R}^{V[\mathcal{G}]}]$, fix $X \in \mathbb{R}^{V[\mathcal{G}]}$ so that π_1 is ordinal definable from X in $L[\mathbb{R}^{V[\mathcal{G}]}]$. Work in $V[X]$ and note that $V[\mathcal{G}]$ is a generic extension of $V[X]$ obtained by adding ω_1 -many Cohen reals. (The forcing factors.)

Representing Automorphisms

continuity on sufficiently generic sets

3. Consider a set G , of degree g , which is Cohen generic over $V[X]$. $\pi_1(g)$ is arithmetically definable relative to g and $\pi^{-1}(0')$. We can find an e and a k such that it is forced that $\pi_1(g)$ is represented by $\{e\}((G \oplus \Pi^{-1}(\emptyset'))^{(k)})$. Since G is Cohen generic, we can assume that e has the form $\{e\}(G \oplus \Pi^{-1}(\emptyset')^{(k)})$. Thus, π_1 is continuously represented on the set of $V[X]$ -generic reals.

4. We make an aside to exploit a phenomenon first observed by Jockusch and Posner, 1981: For any \vec{D} , the \vec{D} -generic degrees generate \mathfrak{D} under meet and join. We fix a mechanism by which this coding can be realized.

Join Coding

definition

Let G and Y be given. Define $\mathbb{C}(Y, G)$ by injecting the values of Y into G_{even} at those places where G_{odd} is not zero.

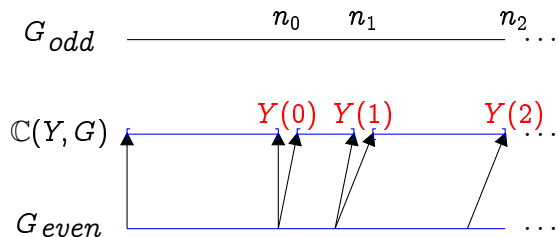
$$G_{\text{odd}} \quad \text{-----} \quad \begin{matrix} n_0 & n_1 & n_2 & \dots \end{matrix}$$

$$G_{\text{even}} \quad \text{-----} \quad \dots$$

Join Coding

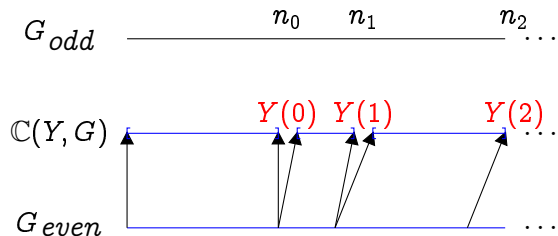
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Join Coding

degree equivalence and genericity

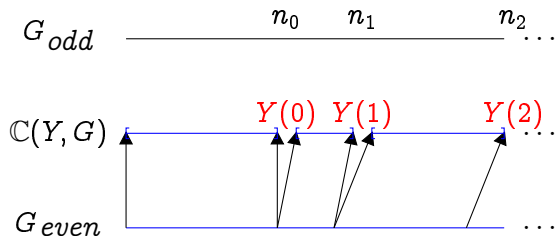


Lemma

If G_{odd} is infinite, then $\mathbb{C}(Y, G) \oplus G \equiv_T Y \oplus G$.

Join Coding

degree equivalence and genericity



Lemma

If G_{odd} is infinite, then $\mathbb{C}(Y, G) \oplus G \equiv_T Y \oplus G$.

Lemma

For any dense open subset of $2^{<\omega}$, D , there is a dense open set D^* , such that for all D^* -generic G and all Y , $\mathbb{C}(Y, G)$ is D -generic. In particular, for all G , Y , and Z , if G is generic over $V[Z]$, then so is $\mathbb{C}(Y, G)$.

Join Coding

generating \mathfrak{D}

Definition

For $Y \in 2^\omega$, let (Y) denote the set $\{Z : Z \leq_T Y\}$.

Join Coding

generating \mathcal{D}

Definition

For $Y \in 2^\omega$, let (Y) denote the set $\{Z : Z \leq_T Y\}$.

4. (continued) Let Y be given with Turing degree y , and let G_1 and G_2 be mutually Cohen generic over $V[X \oplus Y]$. We can write the ideal generated by Y as the meet of joins of generic ideals.

$$\begin{aligned}(\mathbb{C}(Y, G_1) \oplus G_1) \cap (\mathbb{C}(Y, G_2) \oplus G_2) &= (Y \oplus G_1) \cap (Y \oplus G_2) \\ &= (Y)\end{aligned}$$

Representing Automorphisms

pulling back to arithmetic generics

5. The previous equality is preserved by π_1 , as represented on generic reals.

$$\begin{aligned} \{Z : \text{the degree of } Z \text{ belongs to } (\pi^*(y))\} = \\ \left(\{e\}(\mathbb{C}(Y, G_1) \oplus \Pi^{-1}(\emptyset')^{(k)}) \oplus \{e\}(G_1 \oplus \Pi^{-1}(\emptyset')^{(k)}) \right) \\ \cap \left(\{e\}(\mathbb{C}(Y, G_2) \oplus \Pi^{-1}(\emptyset')^{(k)}) \oplus \{e\}(G_2 \oplus \Pi^{-1}(\emptyset')^{(k)}) \right) \end{aligned}$$

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When Y is also generic:

$$\begin{aligned} \left(\{e\}(Y \oplus \Pi^{-1}(\emptyset')^{(k)}) \right) = \\ \left(\{e\}(\mathbb{C}(Y, G_1) \oplus \Pi^{-1}(\emptyset')^{(k)}) \oplus \{e\}(G_1 \oplus \Pi^{-1}(\emptyset')^{(k)}) \right) \\ \cap \left(\{e\}(\mathbb{C}(Y, G_2) \oplus \Pi^{-1}(\emptyset')^{(k)}) \oplus \{e\}(G_2 \oplus \Pi^{-1}(\emptyset')^{(k)}) \right) \end{aligned}$$

Representing Automorphisms

pulling back to arithmetic generics

5. (continued) The previous equation

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expresses an arithmetic identity between generic reals. It holds when Y is arithmetically generic in V and G_1 and G_2 are $V[X]$ -generic. In this case, the right-hand-side represents $(\pi(\text{degree}(Y)))$. Thus, we have an arithmetic (hence continuous) representation of π on all arithmetically generic reals.

Representing Automorphisms

pulling back to arithmetic generics

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Corollary

If $\pi : \mathfrak{D} \xrightarrow{\sim} \mathfrak{D}$ then π has a Borel representation.

Arithmetic Representations

images of generic degrees

Theorem

Suppose that $\pi : \mathcal{D} \xrightarrow{\sim} \mathcal{D}$ and that \vec{D}^ is a countable collection of dense open subsets of $2^{<\omega}$. There is a countable collection of dense open subsets of $2^{<\omega}$, \vec{D} , such that for all \vec{D} -generic reals G there is a G^* such that G^* is \vec{D}^* -generic and the degree of G^* is less than or equal to $\pi(\text{degree}(G))$*

Arithmetic Representations

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- ▶ We chose \vec{D}_0 as above so that we have a continuous representation of π on the collection of \vec{D}_0 -generic reals.
- ▶ We show that for any sufficiently generic real G , $\pi(\text{degree}(G))$ computes a function which has no **a priori** bound on its rate of growth. This reflects the fact that arbitrarily much of $\mathbb{C}(Y, G)$ can be specified between injecting values of Y .

Arithmetic Representations

images of generic degrees

- ▶ We define a real G^* by comparing the values of several of these functions. Hence, for any sufficiently generic G , G^* will be defined at every argument and will be recursive in any representative of $\pi(\text{degree}(G))$.
- ▶ We show that for any dense open set D^* and any finite condition p , there is an extension q of p such that for any G extending q , G^* meets D^* .
- ▶ At the end of the analysis, we calculate the genericity needed on G to ensure that G^* is \vec{D} -generic.

Evaluating relative to a generic degree

Theorem

There is a family of dense open sets \vec{D} and a continuous function $F(G, X)$ such that for all \vec{D} -generic G ,

$$(\forall X)[\pi(\text{degree}(X \oplus G)) = \text{degree}(F(G, X) \oplus \Pi(G))]$$

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Proof

Apply the continuous representation on generic reals to $\mathbb{C}(X, G)$. □

Evaluating relative to a generic degree

Theorem

There is a family of dense open sets \vec{D} and an *injective* continuous function $F(G, X)$ such that for all \vec{D} -generic G , if $\Pi(G)$ is a representative of $\pi(\text{degree}(G))$, then

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Evaluating relative to a generic degree

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Proof

Use G_{odd} to obtain a perfect tree of mutually generic reals. Let X pick a path $T(X)$. Use the previous F on G_{even} and $\mathbb{C}(T(X), G_{\text{even}})$. □