

# Definability in fields

## Lecture 3:

### Finding structure through transcendence

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# Finitely generated fields

A **finitely generated field** is simply a field  $K$  in which there is some **finite** set  $A = \{a_1, \dots, a_n\} \subseteq K$  for which every element of  $K$  may be written as  $P(a_1, \dots, a_n)/Q(a_1, \dots, a_n)$  for some polynomials  $P$  and  $Q$  with integer coefficients.

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- Number fields
- Fields of rational functions  $k(t)$  over finitely generated fields
- Finite extensions of finitely generated fields
- **Function fields**

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# Algebraic varieties over $\mathbb{C}$

## Proposition

*If  $V$  is a projective algebraic variety over  $\mathbb{C}$ , then the set of meromorphic functions on  $V$  is naturally a finitely generated field over  $\mathbb{C}$  and every field finitely generated over  $\mathbb{C}$  has this form.*

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More algebraically, if  $V \subseteq \mathbb{C}^n$  is an irreducible affine complex algebraic variety, then

$I(V) := \{f \in \mathbb{C}[X_1, \dots, X_n] \mid (\forall \mathbf{a} \in V) f(\mathbf{a}) = 0\}$  is a prime ideal and the field of rational functions on  $V$ ,  $\mathbb{C}(V)$ , may be expressed as the field of fractions of  $\mathbb{C}[X_1, \dots, X_n]/I(V)$ .

# Finitely generated fields as function fields

In general, if  $K$  is finitely generated over the field  $k$ , then  $K$  may be expressed as a function field over  $k$ .

- Choose generators  $a_1, \dots, a_n \in K$ .
- The ideal  $\mathfrak{p} := I(\mathbf{a}/k) := \{f \in k[X_1, \dots, X_n] \mid f(\mathbf{a}) = 0\}$  is prime.
- If  $V$  is the algebraic variety defined by  $\mathfrak{p}$ , then  $K \cong k(V)$ .

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# How much does $k(V)$ know about $V$ ?

- The association from an algebraic variety  $V$  over the field  $k$  to the finitely generated (over  $k$ ) field  $k(V)$  is well-defined.
- Our construction of an inverse requires the choice of generators.
- In general, non-isomorphic varieties may yield the same field.
- In fact, there need not even be a “best choice” of a variety.

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# Function fields of smooth curves

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## Theorem

*The association  $C \mapsto k(C)$  is an equivalence of categories between the category of **smooth, projective, absolutely irreducible, curves** over the field  $k$  and the category of finitely generated (over  $k$ ) fields of transcendence degree one over  $k$ .*

# Changing the ground field

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## Answer

Change the ground field.

# A construction

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If  $a_1, \dots, a_{n-1} \in K$  are algebraically independent, then  $k := \mathbb{Q}(a_1, \dots, a_{n-1})^{\text{alg}} \cap K$  is a relatively algebraically closed subfield of  $K$  with  $\text{tr. deg}_k(K) = 1$ .



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Thus, there is a unique smooth, projective, absolutely irreducible curve over  $k$  for which  $K = k(C)$ .

# Observations about the construction

- Whilst  $C$  is determined from  $K$  and  $k$ ,  $k$  depends on a choice.
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# Natural $\mathcal{L}_{\omega_1, \omega}$ definition of algebraic dependence

The elements  $a_1, \dots, a_n$  in a field  $K$  are algebraically dependent if and only if

$$K \models \bigvee_{F \in \mathbb{Z}[X_1, \dots, X_n]} (\exists x_1, \dots, x_n)(F(\mathbf{a}) = 0 \ \& \ F(\mathbf{x}) \neq 0)$$

# Poonen's definition of algebraic dependence

## Theorem (Poonen, after Pop)

*For each positive integer  $n$  there is a formula  $\delta_n(x_1, \dots, x_n)$  in the language of rings having  $n$  free variables for which for any finitely generated field  $K$  and  $n$ -tuple  $\mathbf{a} \in K^n$  one has  $K \models \delta_n(\mathbf{a})$  if and only if the tuple is algebraically dependent.*

# Basics of quadratic forms

For  $K$  a field and  $\mathbf{b} = (b_1, \dots, b_d) \in (K^\times)^d$ , the Pfister form associated to  $\mathbf{b}$  is

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A quadratic form  $q$  **represents zero** if there is a nontrivial solution to the equation  $q(\mathbf{a}) = 0$ . It is **universal** if for every  $r \in K^\times$  there is a solution to the equation  $q(\mathbf{a}) = r$ .



# Pop's definition of transcendence degree

## Theorem (Pop)

*If  $K$  is a finitely generated field of characteristic zero, then  $\text{tr. deg}(K) = d$  if and only if for every  $d + 2$ -tuple  $(b_1, \dots, b_{d+2}) \in (K[\sqrt{-1}]^\times)^{d+2}$  the form  $q_{\mathbf{b}}$  is universal while for some choice the form does not represent zero.*

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The proof uses Voevodsky's theorem on the Milnor conjecture relating Galois cohomology groups, Milnor  $K$ -groups, and Witt groups.

# Pop's near definition of $\delta$

## Theorem (Pop)

If  $K$  is a finitely generated field of characteristic zero and transcendence degree  $d$ , then if  $c$  and  $d$  are algebraic numbers for which  $q_{(t_1, \dots, t_d, c, d)}$  does not represent zero over  $K[\sqrt{-1}]$ , then  $(t_1, \dots, t_d)$  are algebraically independent. **Almost** conversely, if  $(t_1, \dots, t_d)$  is a transcendence basis, then for **many** choices of  $(a_1, \dots, a_d, c, d) \in \mathbb{Z}^{d+2}$ , the form  $q_{(t_1 - a_1, \dots, t_d - a_d, c, d)}$  does not represent zero.

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Poonen's refinement is proven by showing that the relevant algebraic and integer points may be recognized as the coördinates of points on certain elliptic curves.

- The algebraic closure of the prime field is definable as  $k := \delta_1(K)$ .
- If  $a_1, \dots, a_d \in K$  are algebraically independent, then the relative algebraic closure of the field generated by  $a_1, \dots, a_d$  is (parametrically) definable as  $\delta_{d+1}(K; \mathbf{a})$ .
- Consequently, every infinite finitely generated field has an undecidable theory.
- If  $K$  is a finitely generated field of positive transcendence degree, then there is a (parametrically) definable relatively algebraically closed subfield  $k \subseteq K$  for which  $\text{tr. deg}_k(K) = 1$ .

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# Know thyself

If  $K$  is a finitely generated field of transcendence degree at least one, then we can express  $K$  as  $K = k(C)$  where  $k$  is a (parametrically) definable relatively algebraically closed subfield and  $C$  is a smooth projective curve over  $k$ .

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Working by induction on the transcendence degree and using Rumely's theorem in the base case, we may assume that  $k$  is (parametrically) biinterpretable with  $\mathbb{Z}$ . In particular, every arithmetic (relative to the standard recursive presentation of  $k$ ) set in  $k$  is definable.

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To conclude that  $K$  is biinterpretable with  $\mathbb{Z}$  it would suffice for it to recognize itself as a field of functions.

# Valuations on curves

On the field  $k(t)$ , for any  $a \in k$  there is a valuation  $\text{ord}_a : k(t) \rightarrow \mathbb{Z} \cup \{\infty\}$  given by the order of vanishing at  $a$ .

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More generally, for any smooth projective curve  $C$  over  $k$  and closed point  $P \in C$  there is a valuation  $\text{ord}_P$  on  $k(C)$  given by order of vanishing at  $P$  and  $f(P) = b$  just in case  $\text{ord}_P(f - b) > 0$ .

# Which local-global principle?

To define the valuations on function fields of curves we use a local-global principle for Brauer groups (proven by Auslander and Brumer), though if one ignores characteristic two the Witt index theorem would suffice.

# Central simple algebras

Let  $\ell$  be a prime,  $K$  a field of characteristic different from  $\ell$ ,  $\omega \in K^\times$  an  $\ell^{\text{th}}$  root of unity, and  $A$  and  $B$  two nonzero elements of  $K$ . Then  $D(A, B, \omega; K)$  is the noncommutative  $K$ -algebra generated by  $\alpha$  and  $\beta$  subject to the relations  $\alpha^\ell = A$ ,  $\beta^\ell = B$ , and  $\beta\alpha = \omega\alpha\beta$ .



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## Theorem

*$D(A, B, \omega; K)$  is a division ring if and only if  $A$  is not an  $\ell^{\text{th}}$  power in  $K$  and  $B$  is not a norm from  $K(\sqrt[\ell]{A})$ .*

# Auslander-Brumer local-global principle

## Theorem

*Let  $\ell$  be a prime,  $k$  be a field,  $\omega \in k$  an  $\ell^{\text{th}}$  root of unity,  $A, B \in k(t)^{\times}$  two nonzero rational functions over  $k$ . Then  $D(A, B, \omega; k(t))$  is a division ring if and only if there is some completion  $K$  of  $k(t)$  with respect to a valuation trivial on  $k$  for which  $D(A, B, \omega; K)$  is a division ring.*

# Defining valuations

- Using the local-global principle, we cook up a couple of rings depending on the parameter  $f \in k(t)$  which will be division rings just in case  $\text{ord}_a(f) \equiv 0 \pmod{\ell}$ .
- To be honest, the formulas in question only work when  $k$  is replaced by some kind of complete field and a density argument is required to encode everything in  $k(t)$ .
- One extends to  $k(C)$  for more general curves analogously to J. Robinson's method of studying number rings.

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# Geometric problems

- Is  $\text{Th}(\mathbb{C}(t))$  decidable?
- Is  $\mathbb{C}[t]_{(t)}$  parametrically definable in  $\mathbb{C}(t)$ ?
- If  $K$  and  $L$  are finitely generated over the same algebraically closed field  $k$  and  $K \equiv L$ , must we have  $K \cong L$ ?

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# Arithmetic problems

- Is there a straightforward way to deduce the existence of Gödel coding in finitely generated fields from the Pop/Poonen transcendence definition?
- Is there an alternate way to demonstrate undecidability of the theory of  $\mathbb{Q}$  without directly interpreting  $\mathbb{Z}$ ?
- Is there a uniform interpretation of infinite finitely generated fields with  $\mathbb{Z}$ ?
- $\mathbb{Q}^{\text{alg}}(t, s)$  and  $\mathbb{Z}$  interpret each other.

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