

Thm [Division algorithm]

① If $f, g \in K[x]$, $g \neq 0$, then $\exists q, r \in K[x]$ such that $f = gq + r$, with $r = 0$ OR $0 \leq \deg r < \deg g$.

② If f, g have no common constant factors, \exists polynomials r, s such that $fr + gs = 1$.

③ If $p \in K[x]$ is irreducible and $p \mid fg$ then $p \mid f$ or $p \mid g$.

④ Every non-zero polynomial f can be written as $f = c p_1 \dots p_r$ with c a unit and p_i irreducible polynomials. This is unique up to reordering and multiplication by units.

① Follows from "long division" (Thm 23.1)

② Is the Euclidean algorithm of polynomials

④ I prove in class assuming ③.

Ex. $x^2 + 1$ is irreducible in $\mathbb{R}[x]$

• $x^3 + x + 1 = (x+1)(x^2 - x + 2) - 1$

• $x+1$ is a factor of $x^3 + 1$

Defn. Let $f \in K[x]$. An element $\alpha \in K[x]$ is called a root / zero of f if $f(\alpha) = 0$ in K .

• If $K' \supseteq K$ is a larger field and $\alpha \in K'$

satisfies $f(\alpha) = 0$, then we say α is a root of f in K' , (In general, we can evaluate polynomials in larger fields)

(Of course α is a root $\Leftrightarrow f \in \text{Ker } ev_\alpha$, where $ev_\alpha : K[x] \rightarrow K, f(x) \mapsto f(\alpha)$)

Ex i is a root of $x^2 + 1$ in \mathbb{C} , not in \mathbb{R} .

Prop Let $f \in K[x]$. Then f has a root in K
 $\Leftrightarrow f$ has a linear factor

PF If $f(x) = (x-c)g(x)$ for $c \in K, g(x) \in K[x]$
then $f(c) = (c-c)g(c) = 0$
 $\Rightarrow f$ has a root in K .

Conversely, $f \in \text{Ker } ev_\alpha$ for some $\alpha \in K$.

Apply division algorithm to f and $x-\alpha$
 $\therefore f(x) = (x-\alpha)q + r$ with $r=0$ or $\deg r < \deg(x-\alpha) = 1$. Thus r is a constant.

But $0 = f(\alpha) = (\alpha-\alpha)q + r = r$

$$\Rightarrow r = 0$$

$$\Rightarrow f(x) = (x-\alpha)q$$

$\Rightarrow f$ has a linear factor \blacksquare

Cor Let $f(x) \in K[x]$ be a polynomial of degree $n \geq 0$.
Then it has at most n distinct roots

Pf If α_1 is a zero of $f(x)$, then
 $f(x) = (x - \alpha_1) g_1(x)$ and $\deg g_1(x) = n - 1$
 If α_2 is a zero of $g_1(x)$ we have $f(x) = (x - \alpha_1)(x - \alpha_2) g_2(x)$.
 Proceeding in this way we obtain,
 $f(x) = (x - \alpha_1) \cdots (x - \alpha_r) g_r(x)$ with g_r having
 no zeroes in K .
 Since $\deg f = n \Rightarrow r \leq n$. Moreover, if
 $f(b) = 0 \Rightarrow (b - \alpha_1) \cdots (b - \alpha_r) g_r(b) = 0$
 $\Rightarrow b - \alpha_i = 0$ for some i as $g_r(b) \neq 0$ and K is a domain.
 $\Rightarrow \alpha_1, \dots, \alpha_r$ are all the zeroes of f .

Cor If $f, g \in K[x]$ of degree n and f, g agree for at least
 $n + 1$ values in K , then $f = g$.

Cor Let K be an infinite field. Let $f, g \in K[x]$
 s.t. $f(\alpha) = g(\alpha) \forall \alpha \in K$. Then $f = g$.

Exercise Show $f(x) = x^p - x \in (\mathbb{Z}/p\mathbb{Z})[x]$ is zero for
 every value of $\mathbb{Z}/p\mathbb{Z}$.

Conclusion Polynomials SHOULD NOT be thought of as
 functions