

① The Question

Question Which partial orders can be embedded in the Turing degrees in a Borel way?

More formally...

written $A \leq_B B$

Def If A and B are Borel binary relations on 2^ω , say A is Borel reducible to B if there is a Borel function $f: 2^\omega \rightarrow 2^\omega$ such that

$$\forall x, y \in 2^\omega \quad A(x, y) \Leftrightarrow B(f(x), f(y))$$

Question, restated Which Borel ~~partial orders~~ ^{quasi-orders} on 2^ω are Borel reducible to Turing reducibility?

Seems kind of unfair, since \leq_T is not a partial order on 2^ω , but just a quasi-order

→ reflexive, transitive relation
= "partial order" where some elements are equivalent

Question, restated Which Borel ~~partial orders~~ ^{quasi-orders} on 2^ω are Borel reducible to Turing reducibility?

In the Turing degrees, every element has at most countably many predecessors

The same must be true of any quasi-order which is reducible to \leq_T

Is that the only restriction?

Def A quasi-order (Q, \leq) is called **locally countable** if for any $q \in Q$, the set $\{p \in Q \mid p \leq q\}$ is countable

Question, final form Is every locally countable Borel quasi-order on 2^ω Borel reducible to Turing reducibility?

② Why?

Why should we care about this question?

Two reasons

① Sacks's question: which partial orders embed into the Turing degrees?
This is the Borel version

② Kechris's conjecture: Every countable Borel equivalence relation is Borel reducible to Turing equivalence
This is a strengthening of Kechris's conjecture

contradicts Martin's conjecture (Part 1)

Sacks's question & Kechris's conjecture/Martin's conjecture are two of the most long-standing open questions about the global structure of the Turing degrees

②.1 Sacks's Question

Question Does every locally countable, size-continuum partial order embed into the Turing degrees

↳ no req't to be Borel
↳ partial vs. quasi-order makes no difference here

locally countable partial order

What's known

① Every size ω_1 **lcpo** embeds
so answer is "yes" assuming CH

② Independent of ZF

③ Universal size continuum lcpos exist
"Hereditarily a column of"

What's not known

Independent of ZFC?

2.2 Martin's conjecture ← Part 1 only

Basic idea If $f: \mathcal{D}_T \rightarrow \mathcal{D}_T$ is a Borel function then either eventually f is constant or eventually $f(x) \geq_T x$

Def A cone of Turing degrees is a set of the form $\{y \mid y \geq_T x\}$ for some fixed x for some x ← what does this mean? ← called the "base" of the cone

Def $f: 2^\omega \rightarrow 2^\omega$ is Turing invariant if $\forall x, y \in 2^\omega \quad x \equiv_T y \Rightarrow f(x) \equiv_T f(y)$

Conjecture (Part 1 of Martin's conjecture) If $f: 2^\omega \rightarrow 2^\omega$ is a Turing invariant Borel function then either

① $\exists c$ s.t. $f(x) \equiv_T c$ for all x in some cone

or ② $f(x) \geq_T x$ for all x in some cone

super useful

Thm (Martin) Every Borel set of Turing degrees either contains a cone or is disjoint from a cone

← uses Borel determinacy

②.3 Kechris's conjecture

Def An equivalence relation is **countable** if all its equivalence classes are countable

Conjecture (Kechris) Every countable Borel equivalence relation on 2^ω is Borel reducible to Turing equivalence
This contradicts part 1 of Martin's conjecture

pf Assume Kechris & Martin both true

Have Borel reduction f of $\mathcal{D}_T \times 2$ to \mathcal{D}_T

Get two functions $\mathcal{D}_T \rightarrow \mathcal{D}_T$

ctbl-to-one $\left\{ \begin{array}{l} f_0: 1^{\text{st}} \text{ copy} \rightarrow \mathcal{D}_T \\ f_1: 2^{\text{nd}} \text{ copy} \rightarrow \mathcal{D}_T \end{array} \right.$

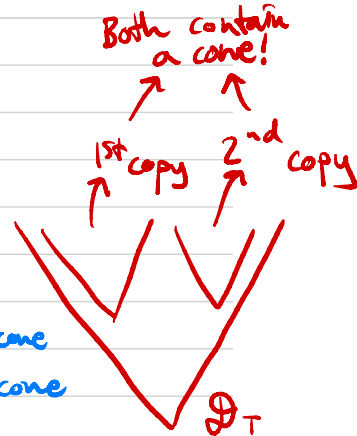
Martin: f_0 not constant on a cone $\Rightarrow f_0(x) \geq_T x$ on a cone

(conjecture) f_1 not constant on a cone $\Rightarrow f_1(x) \geq_T x$ on a cone

Martin: $\text{range}(f_0)$ contains a cone

(thm) $\text{range}(f_1)$ contains a cone

But any 2 cones have a cone in their intersection **Contradiction**

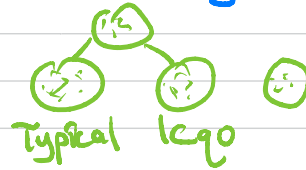
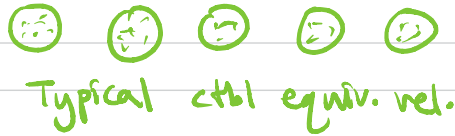


Conjecture (Kechris) Every countable Borel equivalence relation on 2^ω is Borel reducible to Turing equivalence

Original question Is every locally ctbl Borel quasi-order Borel reducible to Turing reducibility?

This is stronger than Kechris's conjecture

A countable equivalence relation is a locally ctbl quasi-order



A Borel reduction from $f: E \rightarrow \leq_T$ from a ctbl Borel equivalence relation (thought of as a Borel lcqo) to \leq_T is a Borel reduction of $E \rightarrow \equiv_T$

It sends different equivalence classes to incomparable Turing degrees rather than just different Turing degrees

③ The Answer

Question Is every locally ctbl Borel quasi-order Borel reducible to Turing reducibility?

Answer No!

Def $f: 2^\omega \rightarrow 2^\omega$ is Turing-order-preserving if
 $\forall x, y \in 2^\omega \quad x \leq_T y \Rightarrow f(x) \leq_T f(y)$

Thm (L.-Siskind) Part 1 of Martin's conjecture holds for all Borel Turing-order-preserving functions

Similar to Kechris $\Rightarrow \neg$ Martin, this shows 2 disjoint copies of \leq_T cannot be Borel reduced to one copy

Natural follow-up question So which ones are Borel reducible to \leq_T ?

③.1 A Generalization

Question Which locally countable Borel quasi-orders are Borel reducible to Turing reducibility?

Example 1 $\leq_T \times 2$ is not reducible

Def A quasi-order (Q, \leq) is countably directed if every countable subset of Q has an upper bound

Example Turing reducibility $\{x_0, x_1, x_2, \dots\}$ has upper bound $\bigoplus_{i \in \mathbb{N}} x_i = \underbrace{\begin{array}{c} | \\ | \\ | \\ \dots \end{array}}_{x_0 \ x_1 \ x_2}$

Example 2 $P_0 \cup P_1$ is not reducible where P_0, P_1 are any uncountable, countably directed Borel lco's

→ Proof is very similar to that on previous page

3.2 3 > 2

Question Which locally countable Borel quasi-orders are Borel reducible to Turing reducibility?

Example 3 (Higuchi-L.) Any locally countable Borel partial order of height 2 is reducible

Example 4 (Higuchi-L.) \exists locally countable Borel partial order of height 3 not reducible

Def A quasi-order (Q, \leq) has height n if the longest strict chain $x_1 \leq x_2 \leq \dots$ has length n
= Q can be divided into n "layers"



Warning Partial order of height 2 is much more restrictive than quasi-order of height 2
Proving example 3 for quasi-orders of height 1 would prove Kechris's conjecture

④ The Proof (sort of)

Thm (Higuchi-L.) Every locally countable Borel partial order of height 2 is Borel reducible to \leq_T

pf (sketch) (Q, \leq) is a locally ctbl Borel partial order of height 2

$Q_1 = 1^{\text{st}}$ level $Q_2 = 2^{\text{nd}}$ level

By Lusin-Novikov, Q_1 & Q_2 are Borel

① Embed 1st level

Classic thm: \exists perfect set of mutually 1-generic reals
Map each $x \in Q_1$ to a diff. generic

② Embed 2nd level

Map each $y \in Q_2$ to a generically chosen upper bound for images of its predecessors in Q_1

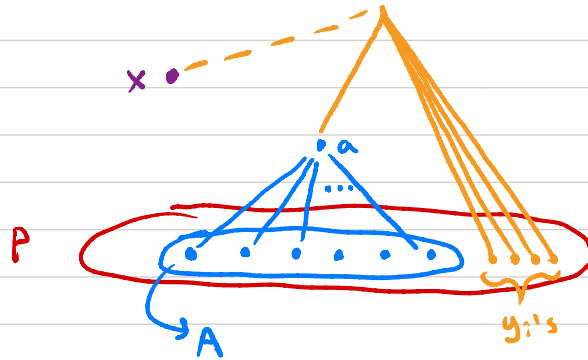
\hookrightarrow can enumerate using Lusin-Novikov

Why can't we just repeat step ② for a 3rd level?

Common ingredient in all
the non-reducibility proofs

A new "basis thm" for
perfect sets

Thm (L.-Siskind) If $P \subseteq 2^\omega$ is a perfect set, $A \subseteq P$ is a
countable dense subset of P and a is a real which
computes every element of A then
 $\forall x \exists y_1, y_2, y_3, y_4 \in P (x \leq_T a \oplus y_1 \oplus y_2 \oplus y_3 \oplus y_4)$



Apology I know this looks really technical 😞

Common ingredient in all
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$$\forall x \exists y_1, y_2, y_3, y_4 \in P (x \leq_T a \oplus y_1 \oplus y_2 \oplus y_3 \oplus y_4)$$

Used in combination with...

Thm (Sushin?) Every uncountable Σ_1^1 set contains a perfect set

& ctt-to-one

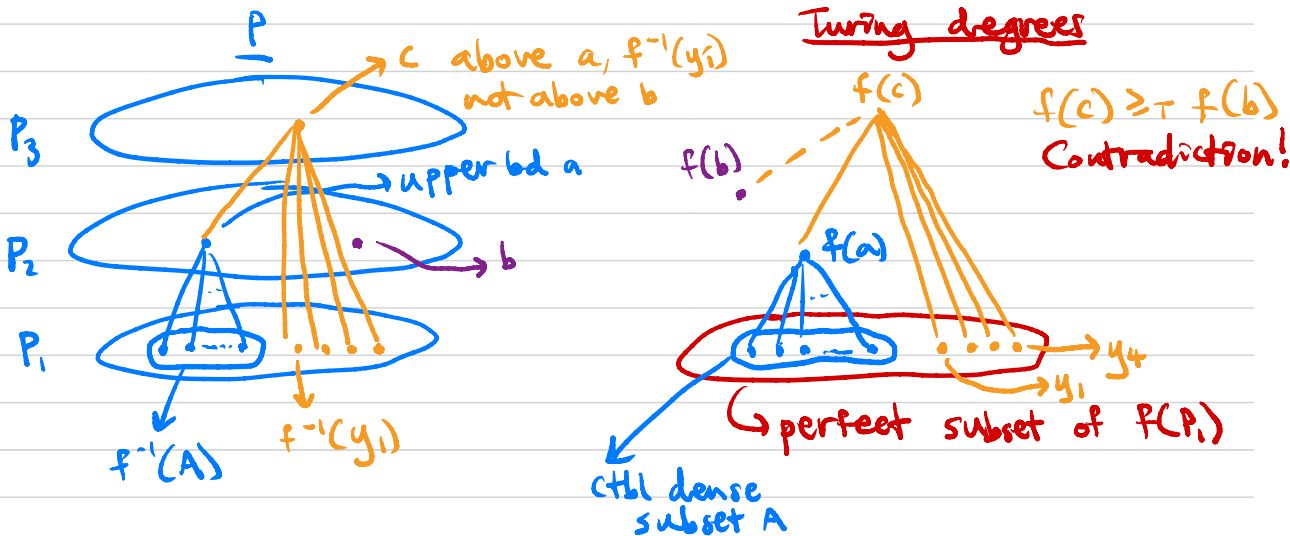
\Rightarrow if $f: 2^\omega \rightarrow 2^\omega$ Borel^v & $A \subseteq 2^\omega$ Borel & uncountable then $f(A)$ contains a perfect set

Thm (Higuchi-L.) There is a locally ctbl Borel partial order of height 3 that is not Borel reducible to \leq_T

pf (sketch) levels of P :

P_1	1st level
P_2	2nd level
P_3	3rd level

$$f: P \rightarrow \leq_T$$



Thm (Higuchi-L.) There is a locally ctbl Borel partial order of height 3 that is not Borel reducible to \leq_T

pf (sketch) Does a P with all necessary properties exist?

What did we use?

- ① Borel, locally countable, height 3
- ② Level 1 uncountable
- ③ Every ctbl subset of level 1 has an upper bound on level 2
- ④ For every finite subset of levels 1 & 2 & p in level 2 not in that subset, there is an upper bound for the finite set in level 3 that is not above the p

Can check that it is easy to construct P to satisfy these conditions

⑤ Lessons

One strategy to answer Sacks's question:

To embed (P, \leq) in \mathcal{D}_T , use transfinite recursion on P to define embedding } "one-by-one" approach

This works if $|P| = \omega_1$

Does not work in general 😞

Thm (Groszek-Slaman) It is consistent that there is a maximal independent set in \mathcal{D}_T of size less than $|\mathcal{Z}^{\omega_1}|$
↳ If we choose badly, we could get stuck

Another approach

Embed 1st level of P as a perfect set of generics, 2nd level as generic upper bds, etc } "all at once" approach

Maybe this works for P well-founded? No!

As soon as you let a perfect set into the range of your embedding, you have lost

⑥ Questions

Two open questions

① Is it provable in ZFC that every locally countable partial order on 2^{ω} of height \leq three embeds in \mathcal{D}_T ?
Probably!
But if so, must use new techniques

② Is every locally countable Borel quasi order of height 1 Borel reducible to Σ_T ?
I think not
Refuting this is easier than refuting Kechris's conjecture