

## Lecture 2-28

Wednesday, February 28, 2018 9:30 AM

Let  $x \in \mathbb{R}^n$ ,  $x = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$ .

We know in the **standard basis**,

$$E = \{e_1, \dots, e_n\}$$

We can write  $x$  as a **linear combination**,

$$x = x_1 e_1 + \dots + x_n e_n$$

Thus we say,  $[x]_E = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} = x$ .

Given an alternative basis  $B = \{b_1, \dots, b_n\}$  of  $\mathbb{R}^n$

we can write  $x$  as a linear combination of these vectors

$$x = \lambda_1 b_1 + \dots + \lambda_n b_n$$

And so we say,  $[x]_B = \begin{bmatrix} \lambda_1 \\ \vdots \\ \lambda_n \end{bmatrix}$ . Note that

$[x]_B = x$  only if  $B$  is the standard basis.

Ex. Given  $B = \{b_1, \dots, b_n\}$  and  $x = \lambda_1 b_1 + \dots + \lambda_n b_n$ ,

determine  $[x]_B$  and  $x$ .

$$(1) B = \left\{ \begin{bmatrix} 1 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ -1 \end{bmatrix} \right\}, \quad x = 3b_1 - 3b_2$$

$$\Rightarrow x = 5 \begin{bmatrix} 1 \\ 1 \end{bmatrix} - 2 \begin{bmatrix} 1 \\ -1 \end{bmatrix}$$

$$(2) \mathcal{B} = \{ [1, 3], [1, 0] \}, \quad \underline{x} = 2\underline{b}_1 + \underline{b}_2$$

$$(3) \mathcal{B} = \left\{ \begin{bmatrix} 1 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ 1 \end{bmatrix} \right\}, \quad \underline{x} = 2\underline{b}_1 + \underline{b}_2 - 3\underline{b}_3$$

Now let's say we have a vector  $\underline{x}$  and we want to know how to represent it in a basis  $\mathcal{B}$ , that is find  $[\underline{x}]_{\mathcal{B}}$ .

Solve one of the two:

$$(a) \underset{\mathcal{B} \leftarrow \mathcal{E}}{P} \underline{x} = [\underline{x}]_{\mathcal{B}} \quad \text{matrix is harder to find, straight forward multiplication}$$

$$(b) \underset{\mathcal{E} \leftarrow \mathcal{B}}{P} [\underline{x}]_{\mathcal{B}} = \underline{x} \quad \text{matrix is easy to find, solve augmented system}$$

Ex.

$$(1) \text{ Let } \underline{x} = \begin{bmatrix} 2 \\ 5 \end{bmatrix} \text{ and } \mathcal{B} = \left\{ \begin{bmatrix} 1 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ -1 \end{bmatrix} \right\}$$

compute  $[\underline{x}]_{\mathcal{B}}$  using (b).

$$\underset{\mathcal{E} \leftarrow \mathcal{B}}{P} = \left[ \begin{array}{c|c} [\underline{b}_1]_{\mathcal{E}} & [\underline{b}_2]_{\mathcal{E}} \\ \hline \end{array} \right] = \begin{bmatrix} \underline{b}_1 & \underline{b}_2 \end{bmatrix}$$

$$(2) \text{ Let } \underline{x} = \begin{bmatrix} 3 \\ 1 \end{bmatrix} \text{ and } \mathcal{B} = \left\{ \begin{bmatrix} 2 \\ 1 \end{bmatrix}, \begin{bmatrix} -1 \\ 3 \end{bmatrix} \right\} \text{ compute}$$

$[\underline{x}]_{\mathcal{B}}$  using (a).

$$P_{\mathcal{B} \leftarrow \mathcal{E}} = P_{\mathcal{E} \leftarrow \mathcal{B}}^{-1}$$

Changing between two bases that aren't standard.

Consider two bases,  $\mathcal{B} = \{b_1, \dots, b_n\}$  and  $\mathcal{C} = \{c_1, \dots, c_n\}$  in  $\mathbb{R}^n$

$$\begin{aligned} [c_1 \ \dots \ c_n \mid b_1 \ \dots \ b_n] &\sim [I \mid [b_1]_{\mathcal{C}} \ \dots \ [b_n]_{\mathcal{C}}] \\ &= [I \mid P_{\mathcal{C} \leftarrow \mathcal{B}}] \end{aligned}$$

$$\begin{aligned} [b_1 \ \dots \ b_n \mid c_1 \ \dots \ c_n] &\sim [I \mid [c_1]_{\mathcal{B}} \ \dots \ [c_n]_{\mathcal{B}}] \\ &= [I \mid P_{\mathcal{B} \leftarrow \mathcal{C}}] \end{aligned}$$

Ex. let  $\mathcal{B} = \left\{ \begin{bmatrix} 0 \\ 3 \end{bmatrix}, \begin{bmatrix} 1 \\ 1 \end{bmatrix} \right\}$  and  $\mathcal{C} = \left\{ \begin{bmatrix} 1 \\ 2 \end{bmatrix}, \begin{bmatrix} 2 \\ 1 \end{bmatrix} \right\}$ .

Suppose  $[x]_{\mathcal{B}} = \begin{bmatrix} 2 \\ 1 \end{bmatrix}$  compute  $[x]_{\mathcal{C}}$