# Quantum computation of Green's functions

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#### QuCQC, Useful Quantum Computation For Quantum Chemistry

February 2021

## Joint work with



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Fast inversion, preconditioned quantum linear system solvers, fast Green's function computation, and fast evaluation of matrix functions, (Tong, An, Wiebe, L., 2008.13295)

## Spectroscopic information and Green's function



Spectral function, 2D Hubbard model.  $A(\mathbf{k}, \omega) = -\frac{1}{\pi} \operatorname{Im}(G(\mathbf{k}, \omega))$ 

DMFT calculation: [Mejuto-Zaera, Zepeda-Nunez, Lindsey, Tubman, Whaley, L., 2020]

Lehmann representation of the single-particle Green's function

$$G(z) = \sum_{n} rac{f_n f_n^{\dagger}}{z - \varepsilon_n + \mathrm{i}\eta \operatorname{sgn}(\varepsilon_n - \mu)}, \quad \eta = 0^+.$$

 $\varepsilon_n$ : quasi-particle energy;  $f_n$ : quasi-particle wavefunction

- Poles: ionization potential, electron affinity.
- Many experiments: photoemission spectroscopy; inverse photoemission spectroscopy; ARPES...

### Quasi-particle and quasi-horse

"Quasi-horse": bare horse + response of dust (Mattuck, 1976)



Quasi-particle: bare particle + response of material

Quasi-electron: added electron + response

Quasi-hole: removed electron + response



#### Quantum many-body problem

- *N* sites (spin-orbitals)
- $\hat{a}_i, \hat{a}_i^{\dagger}, \hat{n}_i$ : annihilation, creation, number operator at site *i*.
- Many-body Hamiltonian (dimension: 2<sup>N</sup>)



•  $|\Psi_0\rangle$ : ground state with  $N_e$  electrons ( $N_e \le 2N$ )  $E_0$ : ground state energy.

#### Problem we consider: large spectral radius

- Planewave / real space refined spatial discretization:  $\|\hat{H}\| \approx \|\hat{H}_0\| \gg \|\hat{H}_1\|$
- Hubbard model, large U limit:  $\|\hat{H}\| \approx \|\hat{H}_1\| \gg \|\hat{H}_0\|$
- Schwinger model (Kogut-Susskind, 1975)
- Write  $\hat{H} = \hat{A} + \hat{B}$ ,  $\|\hat{A}\| \gg \|\hat{B}\|$ , but  $\hat{A}$  is somewhat easy to manipulate.

## Fast Green's function computation

Main result (informal), 2008.13295  $\hat{H} = \hat{A} + \hat{B}$ , with  $\tilde{\sigma}_{\min} = \Omega\left(\eta / \|\hat{B}\|\right)$ , and  $\|\hat{A}\| \gg \|\hat{B}\|$ ;  $|\operatorname{Im}(z)| \ge \eta > 0$  (broadening parameter)

Algorithm	Queries to block-	
	encodings	
HHL	$\widetilde{\mathcal{O}}(rac{ z +\ \hat{H}\ }{\eta^3\epsilon^2})$	
LCU/QSVT	$\widetilde{\mathcal{O}}(rac{ z +\left\Vert \widehat{oldsymbol{H}}  ightert}{\eta^{2}\epsilon})$	
Our work	$\widetilde{\mathcal{O}}(rac{\ m{m{m{m{m{m{m{m{m{m{m{m{m{$	

HHL,(Harrow-Hassidim-Lloyd, 2009); Linear combination of unitaries (LCU),(Childs-Kothari-Somma, 2017); Quantum singular value transformation (QSVT) (Gilyén-Su-Low-Wiebe, 2019)

#### Green's function

• Time-ordered single-particle Green's function (or Green's function for short) in the frequency domain: map  $\mathbb{C} \to \mathbb{C}^{N \times N}$ 

$$G(z) = G^{(+)}(z) + G^{(-)}(z).$$

• Advanced  $(G^{(+)})$  and retarded  $(G^{(-)})$  Green's functions

$$egin{aligned} G_{ij}^{(+)}(z) &:= \left\langle \Psi_0 \left| \hat{a}_i \left( z - \left[ \hat{H} - E_0 
ight] 
ight)^{-1} \hat{a}_j^\dagger \right| \Psi_0 
ight
angle \ G_{ij}^{(-)}(z) &:= \left\langle \Psi_0 \left| \hat{a}_j^\dagger \left( z + \left[ \hat{H} - E_0 
ight] 
ight)^{-1} \hat{a}_i \right| \Psi_0 
ight
angle. \end{aligned}$$

• Assume  $|Im(z)| \ge \eta > 0$  (broadening parameter)

Simplest setting: non-interacting system

• 
$$\hat{H} = \hat{H}_0 = \sum_{ij=1}^N T_{ij} \hat{a}_i^{\dagger} \hat{a}_j.$$

Very simple analytic solution via a (small) matrix inversion

$$G_0(z) = (z-T)^{-1} = \sum_n \frac{f_n f_n^{\dagger}}{z-\varepsilon_n}, \quad Tf_n = \varepsilon_n f_n.$$

- Bare Green's function (bare horse)
- With interaction  $\hat{H} = \hat{H}_0 + \hat{H}_1$ . G(z): quasi-horse
- Self energy (c.f. Frank Wilhelm-Mauch's talk on Monday)

$$\Sigma(z) := G^{-1}(z) - G_0^{-1}(z).$$

## Next simplest setting: quantum impurity

Example: Single-impurity Anderson model (SIAM)

$$\hat{H} = \underbrace{\sum_{\sigma} \epsilon_{f} \hat{t}_{\sigma}^{\dagger} \hat{t}_{\sigma} + \sum_{\langle j, j' \rangle \sigma} t_{jj'} \hat{c}_{j\sigma}^{\dagger} \hat{c}_{j'\sigma} + \sum_{j,\sigma} \left( V_{j} \hat{t}_{\sigma}^{\dagger} \hat{c}_{j\sigma} + V_{j'}^{*} \hat{c}_{j\sigma}^{\dagger} \hat{t}_{\sigma} \right)}_{\hat{H}_{0}} + \underbrace{U \hat{t}_{\uparrow}^{\dagger} \hat{t}_{\uparrow} \hat{t}_{\downarrow}^{\dagger} \hat{t}_{\downarrow}}_{\hat{H}_{1}}$$

- Perturbation to the Green's function is global.
- Self energy  $\Sigma(z)$  is only nonzero on the impurity.
- Foundation of DMFT / CT-QMC etc. "Folk theorem" at least since Feynman (with diagrammatic arguments)
- Non-perturbative proof (for general impurities): [L.-Lindsey, Ann. Henri Poincare 2020]



 $|\uparrow\rangle, |\downarrow\rangle$ 

## Computing Green's functions (with general interaction)

- With HF/DFT: essentially a non-interacting picture
- Small Ĥ<sub>1</sub>: many-body perturbation theory (MBPT). GF2, GW, SOSEX, GFCC..



- Large Ĥ<sub>1</sub>: exact diagonalization / CI, QMC, DMRG, quantum embedding (DMFT/DMET)..
- Quantum computer

#### Quantum strategy: direct computation of G

Brute-force "matrix-matrix-multiplication"

$$egin{aligned} G_{ij}^{(+)}(z) &:= \left\langle \Psi_0 \left| \hat{a}_i \left( z - \left[ \hat{H} - E_0 
ight] 
ight)^{-1} \hat{a}_j^\dagger \left| \Psi_0 
ight
angle &:= \left\langle \Psi_0 | \mathcal{A} | \Psi_0 
ight
angle \ \mathcal{A} &= \hat{a}_i \left( z - \left[ \hat{H} - E_0 
ight] 
ight)^{-1} \hat{a}_j^\dagger \end{aligned}$$

•  $(z - [\hat{H} - E_0])^{-1} \hat{a}_j^{\dagger} |\Psi_0\rangle$ : quantum linear system problem (QLSP)

## Compare the complexities of QLSP solvers

Significant progress in the past few years: Near-optimal complexity matching lower bounds. All

with the promise of poly(N) complexity for matrix of size  $2^N$ .

Algorithm	Query complexity	Remark	
HHL,(Harrow-Hassidim-Lloyd, 2009)	$\widetilde{\mathcal{O}}(\kappa^2/\epsilon)$	w. VTAA, complexity becomes $\widetilde{\mathcal{O}}(\kappa/\epsilon^3)$ (Ambainis 2010)	
Linear combination of unitaries (LCU),(Childs-Kothari-Somma, 2017)	$\widetilde{\mathcal{O}}(\kappa^2 \mathrm{polylog}(1/\epsilon))$	w. VTAA, complexity becomes $\widetilde{\mathcal{O}}(\kappa \operatorname{poly}\log(1/\epsilon))$	
Quantum singular value transfor- mation (QSVT) (Gilyén-Su-Low- Wiebe, 2019)	$\widetilde{\mathcal{O}}(\kappa^2\log(1/\epsilon))$	Queries the RHS only $\widetilde{\mathcal{O}}(\kappa)$ times	
Randomization method (RM) (Subasi-Somma-Orsucci, 2019)	$\widetilde{\mathcal{O}}(\kappa/\epsilon)$	Prepares a mixed state; w. repeated phase estimation, complexity becomes $\widetilde{\mathcal{O}}(\kappa \operatorname{poly}\log(1/\epsilon))$	
Time-optimal adiabatic quantum computing (AQC(exp)) (An-L., 2019, 1909.05500)	$\widetilde{\mathcal{O}}(\kappa \operatorname{poly} \log(1/\epsilon))$	No need for any amplitude amplifi- cation. Use time-dependent Hamil- tonian simulation.	
Eigenstate filtering (LTong, 1910.14596, Quantum 2020)	$\widetilde{\mathcal{O}}(\kappa \log(1/\epsilon))$	No need for any amplitude amplifi- cation. Does not rely on any com- plex subroutines.	

## Near-optimal algorithm for finding the ground energy

		Preparation	Ground energy	Preparation
		(bound known)		(bound unknown)
Uн	Our work	$\mathcal{O}\left(\frac{lpha}{\gamma\Delta}\log(\frac{1}{\epsilon}) ight)$	$\widetilde{\mathcal{O}}\left(rac{lpha}{\gamma h}\log(rac{1}{artheta}) ight)$	$\widetilde{\mathcal{O}}\left(rac{lpha}{\gamma\Delta}\log(rac{1}{artheta\epsilon}) ight)$
-11	GTC19	$\widetilde{\mathcal{O}}\left(\frac{\alpha}{\gamma\Delta}\right)$	$\widetilde{\mathcal{O}}\left(\frac{\alpha^{3/2}}{\gamma\hbar^{3/2}}\right)$	$\widetilde{\mathcal{O}}\left(\frac{\alpha^{3/2}}{\gamma\Delta^{3/2}}\right)$
U	Our work	$\mathcal{O}\left(\frac{1}{\gamma}\right)$	$\widetilde{\mathcal{O}}\left(rac{1}{\gamma}\log(rac{lpha}{\hbar})\log(rac{1}{artheta}) ight)$	$\widetilde{\mathcal{O}}\left(rac{1}{\gamma}\log(rac{lpha}{\Delta})\log(rac{1}{artheta}) ight)$
- 1	GTC19	$\widetilde{\mathcal{O}}\left(\frac{1}{\gamma}\right)$	$\widetilde{\mathcal{O}}\left(\frac{1}{\gamma}\sqrt{\frac{lpha}{\hbar}}\right)$	$\widetilde{\mathcal{O}}\left(\frac{1}{\gamma}\sqrt{\frac{lpha}{\Delta}}\right)$
Extra	Our work	<i>O</i> (1)	$\mathcal{O}(\log(\frac{1}{\gamma}))$	$\mathcal{O}(\log(\frac{1}{\gamma}))$
qubits	GTC19	$\mathcal{O}(\log(\frac{1}{\Delta}\log(\frac{1}{\epsilon})))$	$\mathcal{O}(\log(\frac{1}{h}))$	$\mathcal{O}(\log(\frac{1}{\Delta}\log(\frac{1}{\epsilon})))$

Well-known result: phase estimation; Previous best results: (Ge-Tura-Cirac, 2019) Our work: (L.-Tong, 2002.12508, Quantum 2020) h: precision of the ground energy estimate;  $1 - \vartheta$ : success probability Lower bound for the overlap:  $|\langle \phi_0 | \psi_0 \rangle| \ge \gamma$ , Bounds for the ground energy and spectral gap:  $\lambda_0 \le \mu - \Delta/2 < \mu + \Delta/2 \le \lambda_1$ .

Heisenberg-limited ground state energy estimation for early fault-tolerant quantum computers (L.-Tong, 2102.11340)

#### Block-encoding

• Quantum gates have to be unitary.

• 
$$A = \hat{a}_i \left( z - \left[ \hat{H} - E_0 \right] \right)^{-1} \hat{a}_j^{\dagger}$$
 is not unitary.

 Idea: extend *n*-qubit non-unitary matrix to a (*n* + *m*)-qubit unitary matrix (Low-Chuang, 2016; called "standard form" initially)

$$U_A pprox \left( egin{array}{cc} A/lpha & \cdot \ & \cdot & \cdot \end{array} 
ight)$$

 Many examples of block-encoding: density operators, POVM operators, *d*-sparse matrices, addition and multiplication of block-encoded matrices (Gilyén-Su-Low-Wiebe, 2019)

## Hadamard test for Green's function computation

• If we can block-encode the inverse:  $\left(z - \left[\hat{H} - E_0\right]\right)^{-1}$  $\Rightarrow$  Product of block-encoded matrices

$$\Rightarrow \text{ Product of block-encoded matrices} \\ A = \hat{a}_i \left( z - \left[ \hat{H} - E_0 \right] \right)^{-1} \hat{a}_j^{\dagger}, \text{ call it } U_A$$

Hadamard test circuit



- Success probability p(0) = ½(1 + Re ⟨φ|A|φ⟩). Similar circuit for the imaginary part.
- Amplitude estimation to improve dependence on  $\epsilon$ .

## Fast Green's functions computation

Algorithm	Queries to block-	
	encodings	
HHL	$\widetilde{\mathcal{O}}(rac{ z +lpha_{H}}{\eta^{3}\epsilon^{2}})$	
LCU/QSVT	$\widetilde{\mathcal{O}}(rac{ z +lpha_{H}}{\eta^{2}\epsilon})$	
Our work	$\widetilde{\mathcal{O}}(rac{lpha_{B}}{\widetilde{\sigma}_{\min}^{2}\epsilon})$	

• 
$$\hat{H} = \hat{A} + \hat{B}$$
, with  $\widetilde{\sigma}_{\min} = \Omega(\eta / \alpha_B)$ , and  $\left\| \hat{A} \right\| \gg \left\| \hat{B} \right\|$ .

• Block-encodings in our work involves fast inversion.

#### Fast inversion of diagonal matrices

- $D = \text{diag}(D_{ii})$ :  $||D^{-1}|| = \min |D_{ii}| = \Omega(1), ||D|| = \max |D_{ii}| \gg 1$
- Assume  $O_D \ket{i} \ket{0'} = \ket{i} \ket{D_{ii}}, \quad i \in [N]$
- Circuit  $U'_D$  for the block-encoding of  $D^{-1}$  (classical arithmetic)



• Circuit depth is independent of ||D||

#### Fast inversion beyond diagonal matrices

- 1-sparse matrices  $A = \Pi D$
- Normal matrices  $A = VDV^{\dagger}$

$$U_{\mathcal{A}}' = (V \otimes I_{l+1})U_{\mathcal{D}}'(V^{\dagger} \otimes I_{l+1}).$$

Example:

$$\sum_{\mathbf{x},\mathbf{y},\sigma} \mathcal{T}(\mathbf{x}-\mathbf{y}) \hat{a}^{\dagger}_{\mathbf{x},\sigma} \hat{a}_{\mathbf{y},\sigma} = \text{FFFT}\left(\sum_{\mathbf{G},\sigma} \hat{\mathcal{T}}(\mathbf{G}) \hat{c}^{\dagger}_{\mathbf{G},\sigma} \hat{c}_{\mathbf{G},\sigma}\right) \text{FFFT}^{\dagger}$$

## Preconditioned quantum linear system solver

Consider

$$\left( A+B
ight) \left| x
ight
angle \sim \left| b
ight
angle$$

Preconditioner: A<sup>-1</sup>

$$(I + A^{-1}B) \ket{x} \sim A^{-1} \ket{b}$$

- Condition number:  $\kappa(I + A^{-1}B) \le (1 + ||(A + B)^{-1}||||B||) (1 + ||A^{-1}|||B||)$
- Circuit depth: independent of ||A||

#### Green's function computation for fixed $N_e$

- $[\hat{H}, \hat{P}_{N_e}] = 0$ : Preserve the number of electrons  $N_e$  ( $N_e \ll N$ )
- Cost of preconditioned Hubbard solver in second quantization

$$\widetilde{\mathcal{O}}\left(rac{N^3(\min(|\mathcal{U}|,|t|)^3}{\eta^2\epsilon}\log\left(rac{1}{\delta}
ight)
ight),$$

 Cost of preconditioned Hubbard solver in second quantization, with N<sub>e</sub> scaling

$$\widetilde{\mathcal{O}}\left(\frac{\textit{\textstyle N_e^3}(\min(|\textit{\textit{U}}|,|\textit{t}|)^3}{\eta^2\epsilon}\log\left(\frac{1}{\delta}\right)\right),$$



#### Finite temperature effects (Gibbs state preparation)

• Prepare 
$$\rho_{\beta} = \frac{1}{Z_{\beta}} e^{-\beta H}$$
,  $Z_{\beta} = \text{Tr}(e^{-\beta H})$ .

- Purified Gibbs state  $|\Psi\rangle = \frac{1}{\sqrt{Z_{\beta}}} \sum_{x \in [N]} |x\rangle (e^{-\beta H/2} |x\rangle)$ : trace out first register  $\Rightarrow$  Obtain  $\rho_{\beta}$
- Two new approaches (convert to linear system problems):
  - Cauchy's Contour integral formula:

$$e^{-\beta H} = \frac{1}{2\pi i} \oint_{\Gamma} e^{-\beta z} (z - H)^{-1} dz$$

Inverse transform:

$$\boldsymbol{e}^{-\beta H} = \boldsymbol{e}^{-\beta (H^{-1})^{-1}}$$

## Fast algorithm for preparing $\propto e^{-H} \ket{b}$

	Algorithm	Query complexities
w.o. pre- conditioner	Phase estimation (Poulin- Wocjan, 2009)	$\widetilde{\mathcal{O}}(\frac{\alpha_H}{\xi\epsilon})$
	LCU (van Apeldoorn et al, 2020)	$\widetilde{\mathcal{O}}(rac{lpha_{H}}{\xi}\log(rac{1}{\epsilon}))$
w. precon- ditioner	This work (contour integral)	$\widetilde{\mathcal{O}}(rac{lpha_{\mathcal{B}}}{\xi\widetilde{\sigma}_{\min}^{\prime 2}}\log(rac{1}{\epsilon}))$
	This work (inverse transforma- tion)	$\widetilde{\mathcal{O}}\left( \tfrac{\alpha_{\mathcal{B}}}{\xi \widetilde{\sigma}_{\min}^2} \left[ \log\left( \tfrac{1}{\epsilon} \right) \right]^5 \right)$

 $\xi = \| e^{-H} | b \rangle \|, \widetilde{\sigma}'_{\min} = \Omega(1/\alpha_B), \widetilde{\sigma}_{\min} = \Omega(1/(1 + \| (A + B)^{-1} \| \| B \|))$ 



# Thank you for your attention!

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