

# Classical Shadows

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Charles Hadfield

July 7, 2021



# Outline

## Classical shadows for tomography

- Shadow tomography

- Classical shadows

## The shallow-circuit measurement problem

- Soln 0: Ell-1 sampling

- Soln 0: LDF grouping

- Soln 0: classical shadows

- Soln 1: **locally-biased classical shadows**

- Soln 3: **globally-biased classical shadows**

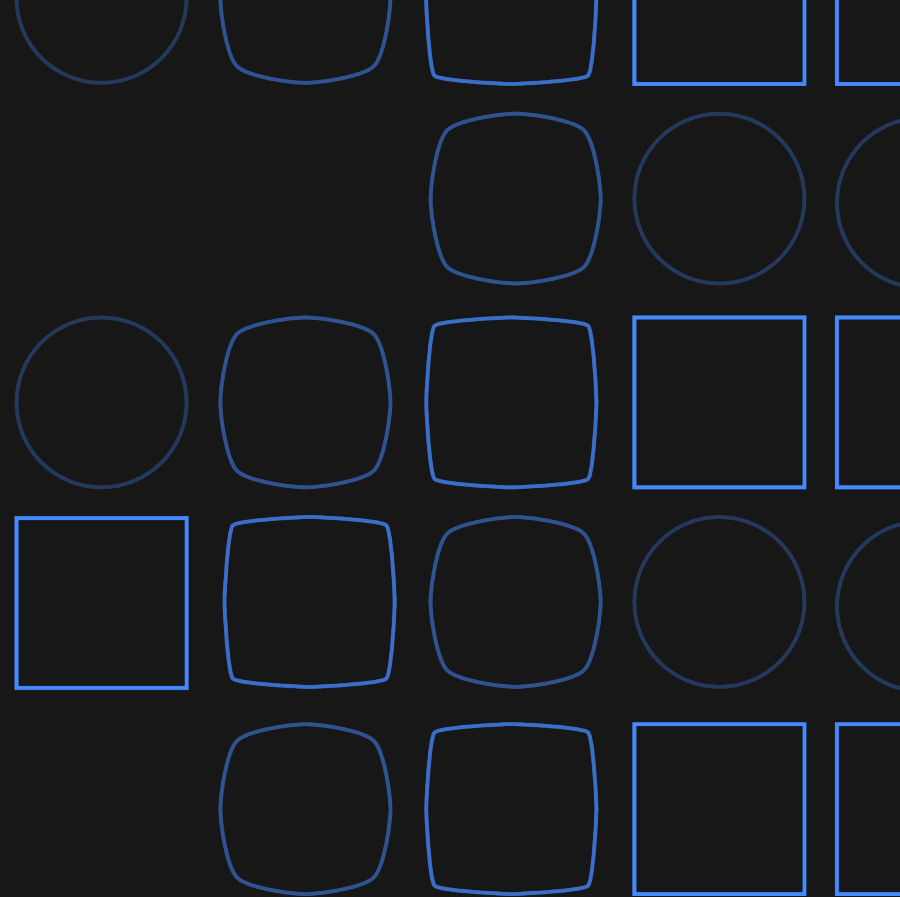
- Soln 2: derandomized classical shadows

- Soln 5: overlapped grouping method

- Soln 4: **adaptive Pauli shadows**

# Classical shadows

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# Predicting many properties of a quantum system from very few measurements

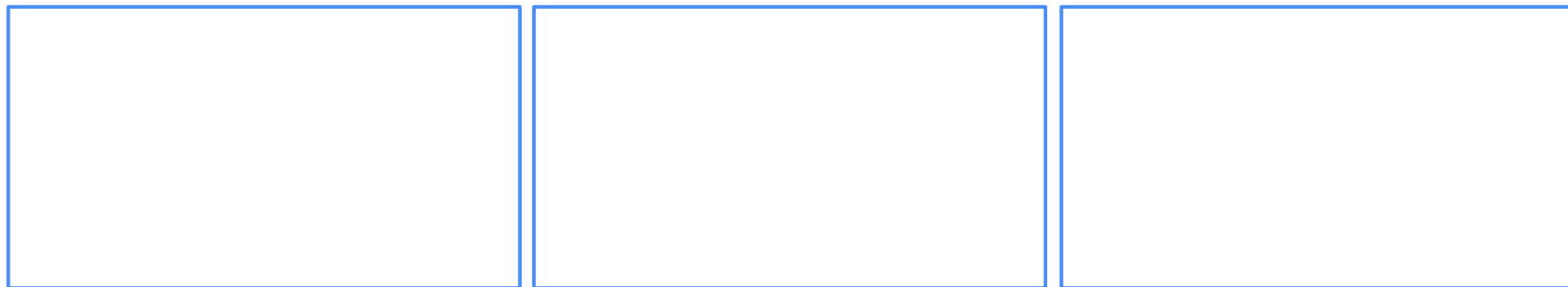
Hsin-Yuan Huang <sup>1,2</sup> , Richard Kueng<sup>1,2,3</sup> and John Preskill<sup>1,2,4</sup>

**Predicting the properties of complex, large-scale quantum systems is essential for developing quantum technologies. We present an efficient method for constructing an approximate classical description of a quantum state using very few measurements of the state. This description, called a 'classical shadow', can be used to predict many different properties; order  $\log(M)$  measurements suffice to accurately predict  $M$  different functions of the state with high success probability. The number of measurements is independent of the system size and saturates information-theoretic lower bounds. Moreover, target properties to predict can be selected after the measurements are completed. We support our theoretical findings with extensive numerical experiments. We apply classical shadows to predict quantum fidelities, entanglement entropies, two-point correlation functions, expectation values of local observables and the energy variance of many-body local Hamiltonians. The numerical results highlight the advantages of classical shadows relative to previously known methods.**

# Classical shadows

Given an  $n$  qubit state and  $M$  expectation values to estimate, how many copies of the state are required?

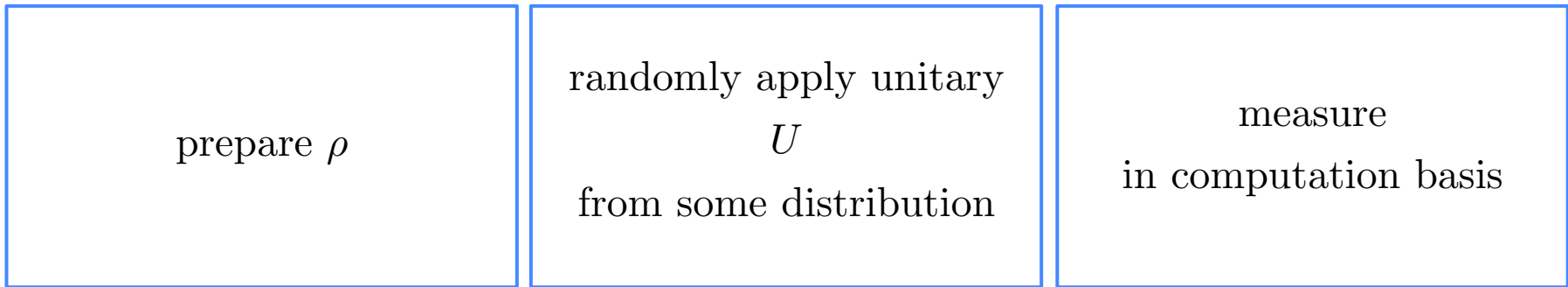
Today's notion of a “shadow”:



# Classical shadows

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And record the outcomes in a classically-efficient data structure.

# Random Clifford measurements

Given an  $n$  qubit state and  $M$  expectation values to estimate, how many copies of the state are required?

Today's notion of a “shadow”:

prepare  $\rho$

randomly apply  
Clifford circuit

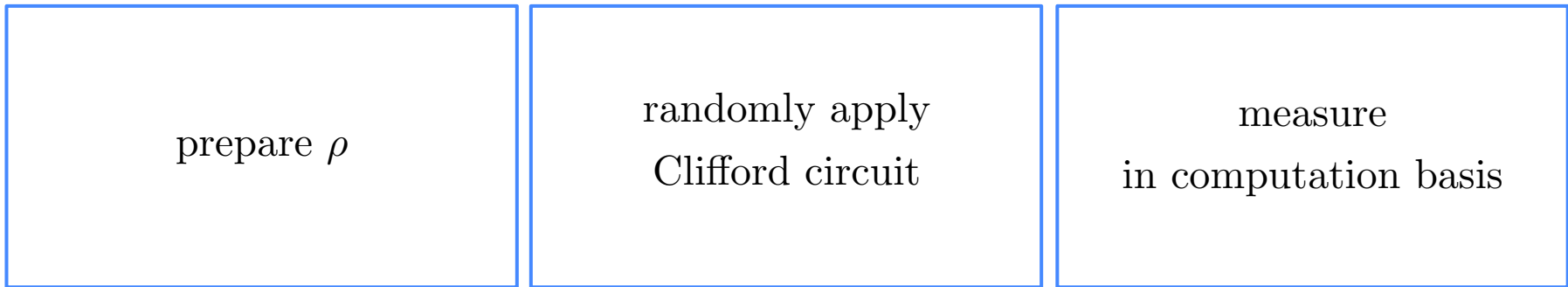
measure  
in computation basis

And record the outcomes in a classically-efficient data structure.

# Random Clifford measurements

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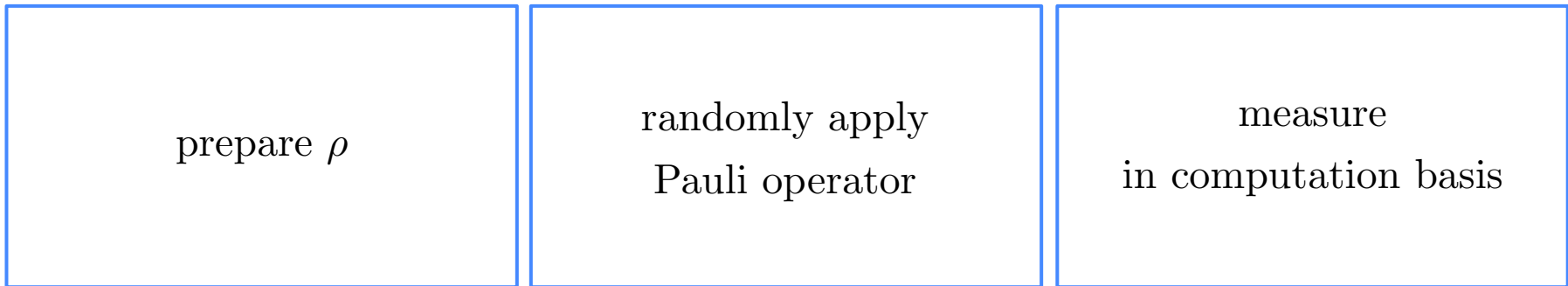
And record the outcomes in a classically-efficient data structure. (Symplectic representation over  $F_2$ )



# Random Pauli measurements

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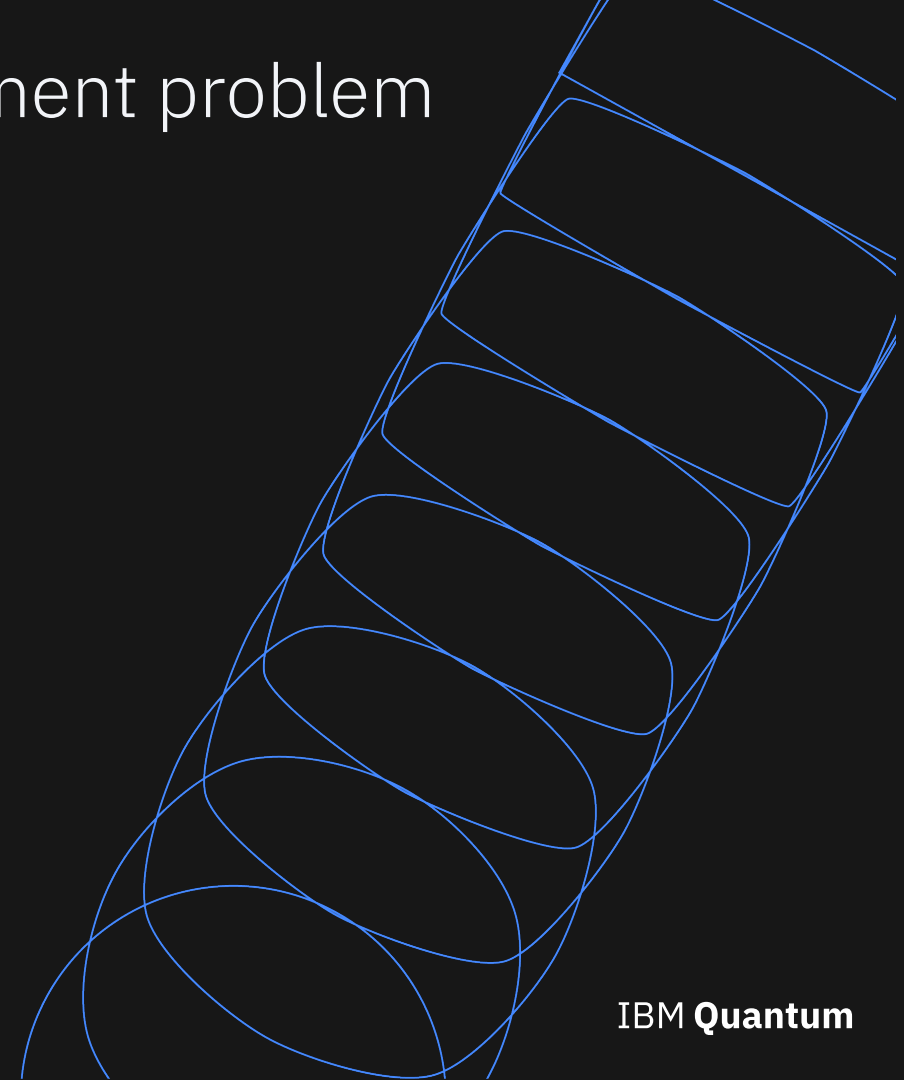
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And record the outcomes in a classically-efficient data structure. (Symplectic representation over  $F_2$ )

# The shallow-circuit measurement problem

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# Shallow circuits for VQE

Our Hamiltonian will be on  $n$  qubits

$$H = \sum_P \alpha_P P \quad P = \bigotimes_{i \in [n]} P_i \quad P_i \in \{I, X, Y, Z\}$$

Once a state  $\rho$  has been prepared, *no* entangling gates may be applied.

Measurement bases will be  $B = \bigotimes_{i \in [n]} B_i \quad B_i \in \{X, Y, Z\}$

How does one best choose the choice of measurement bases in order to estimate the energy, given the state, to accuracy  $\epsilon$ , with *as few as possible* preparations of the state?

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**Algorithm 1** Estimation of Energy

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**for** measurement  $m \in [M]$  **do**

    Prepare state  $\rho$

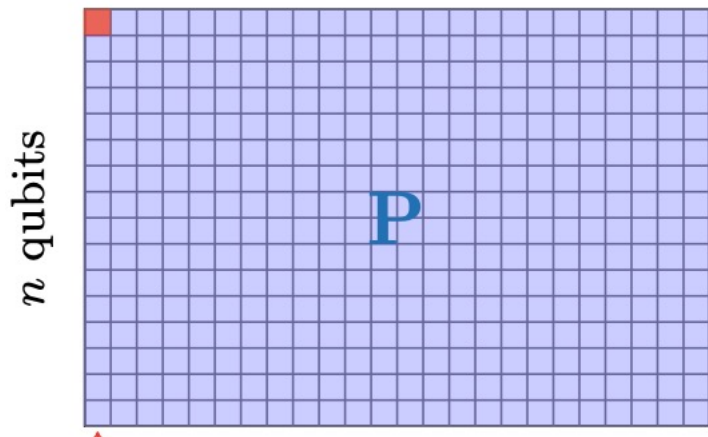
*Choose measurement basis  $B$  according to your algorithm*

    Measure  $\rho$  in basis and estimate Pauli operators

**return** energy estimate

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$M$  Pauli measurements



# Shallow circuits for VQE

Every algorithm's goal will be to build an *unbiased* estimator

$$\mathbb{E}(\nu) = \text{Tr}(H\rho)$$

whose single-shot variance is proportional to the required number of preparations of the state

$$\varepsilon = \sqrt{\frac{\text{Var}(\nu)}{M}}$$

# Shallow circuits for H<sub>2</sub>O on 14 qubits

A back of the envelope demonstration of time-advantages:

Chemical accuracy requires additive-error accuracy of  $\sim 1$  mHartree

So if I finally get a candidate density close to the true ground state, how many shots ( $\times 10^6$ ) are required to get an average error of 1 mHartree?

Ell-1 sampling (2015)	4 400
LDF grouping* (2017)	1 000
Classical shadows (2020)	2 800
<b>Locally-biased CS</b> (2020)	250
<b>Globally-biased CS</b> (2021)	300
Derandomized CS (2021)	15
<b>Adaptive Pauli shadows</b> (2021)	12
Overlapped grouping (2021)	16

Multiply this by how many different guesses in VQE are required before reaching an accurate representation of the ground state...

## Other molecules

		Qiskit	LBCS	GBCS	Caltech	APS
LiH	JW	55	15	10	1	1.6
	P	85	30	15	1	2.5
	BK	75	70	15	1.6	5
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NH <sub>3</sub>	JW	900	350		32	16
	P	2600	630		44	19
	BK	2150	380		14	12

# Solution 1: Locally-biased classical shadows

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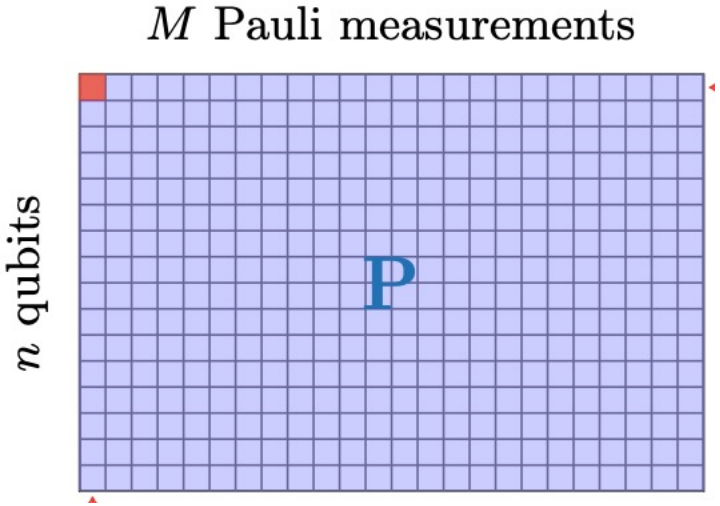
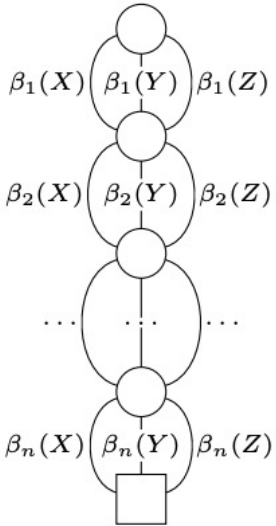
# Locally-biased classical shadows

## MEASUREMENTS OF QUANTUM HAMILTONIANS WITH LOCALLY-BIASED CLASSICAL SHADOWS

CHARLES HADFIELD, SERGEY BRAVYI, RUDY RAYMOND, AND ANTONIO MEZZACAPO

**ABSTRACT.** Obtaining precise estimates of quantum observables is a crucial step of variational quantum algorithms. We consider the problem of estimating expectation values of molecular Hamiltonians, obtained on states prepared on a quantum computer. We propose a novel estimator for this task, which is locally optimised with knowledge of the Hamiltonian and a classical approximation to the underlying quantum state. Our estimator is based on the concept of classical shadows of a quantum state, and has the important property of not adding to the circuit depth for the state preparation. We test its performance numerically for molecular Hamiltonians of increasing size, finding a sizable reduction in variance with respect to current measurement protocols that do not increase circuit depths.

# How do we choose our measurement bases?



# How do we choose our measurement bases?

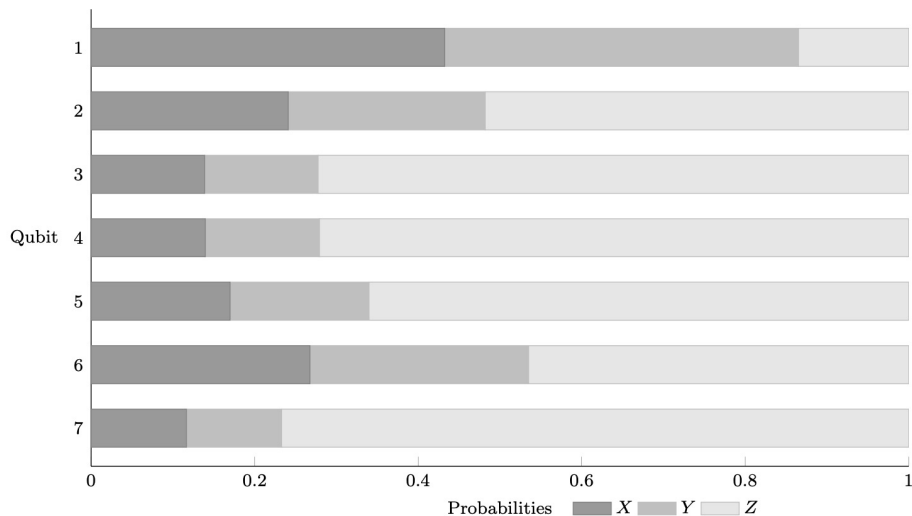
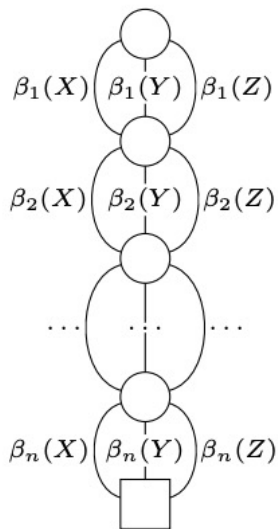


FIGURE 1. Probability distributions over the first 7 of 14 qubits for H<sub>2</sub>O Hamiltonian using the Jordan-Wigner encoding. The probability distributions have been optimised according to Eq. (16).

# Locally-biased classical shadows

Can knowledge of the coefficients *locally* improve random guessing à la classical shadows

$$\beta = \left( \prod_{i \in [n]} \beta_i \right) : \{X, Y, Z\}^n \rightarrow \mathbb{R}^+$$

$$\text{Var}(\nu) \leq \mathbb{E}(\nu^2) = \sum_{P, Q} \alpha_P \alpha_Q g(P, Q, \beta) \text{Tr}(PQ\rho)$$

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$$\text{cost}(\beta | \rho_{HF}) = \dots$$

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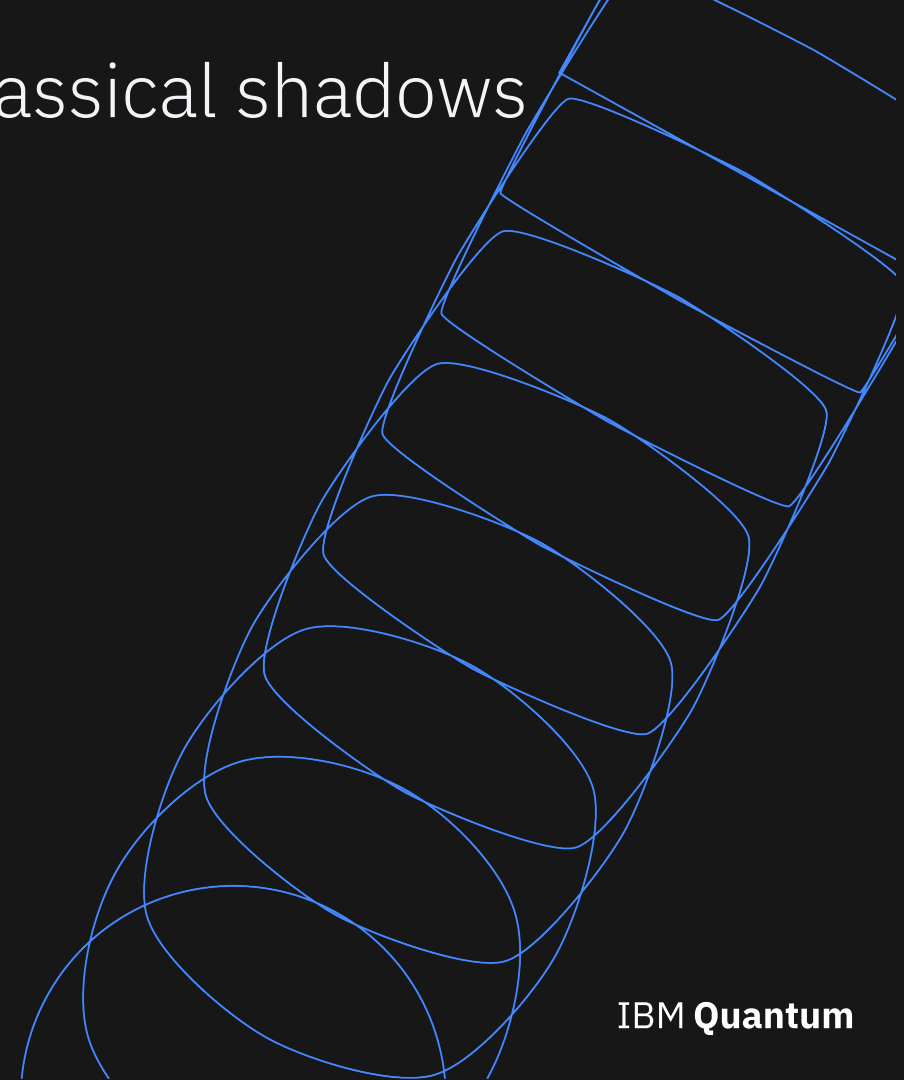
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$$\text{cost}_{\text{diag}}(\beta) = \sum_P \alpha_P^2 \frac{1}{\prod_{i|P_i \neq I} \beta_i(P_i)}$$

# Solution 3: Globally-biased classical shadows

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## Decision Diagrams for Quantum Measurements with Shallow Circuits

Stefan Hillmich,<sup>1,\*</sup> Charles Hadfield,<sup>2,†</sup> Rudy Raymond,<sup>3,4,‡</sup> Antonio Mezzacapo,<sup>2</sup> and Robert Wille<sup>1,5</sup>

<sup>1</sup>*Johannes Kepler University Linz, 4040 Linz, Austria*

<sup>2</sup>*IBM Quantum, IBM T.J. Watson Research Center, Yorktown Heights, NY 10598*

<sup>3</sup>*IBM Quantum, IBM Japan, 19-21 Nihonbashi Chuo-ku, Tokyo, 103-8510, Japan*

<sup>4</sup>*Quantum Computing Center, Keio University, 3-14-1 Hiyoshi,*

*Kohoku-ku, Yokohama, Kanagawa, 223-8522, Japan*

<sup>5</sup>*Software Competence Center Hagenberg (SCCH) GmbH, 4232 Hagenberg, Austria*

We consider the problem of estimating quantum observables on a collection of qubits, given as a linear combination of Pauli operators, with shallow quantum circuits consisting of single-qubit rotations. We introduce estimators based on randomised measurements, which use decision diagrams to sample from probability distributions on measurement bases. This approach generalises previously known uniform and locally-biased randomised estimators. The decision diagrams are constructed given target quantum operators and can be optimised considering different strategies. We show numerically that the estimators introduced here can produce more precise estimates on some quantum chemistry Hamiltonians, compared to previously known randomised protocols and Pauli grouping methods.



# Decision diagrams

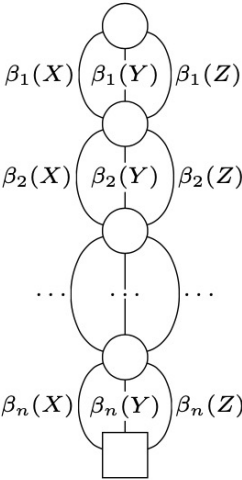


Figure 6: The Decision Diagram of LBCS

# Decision diagrams

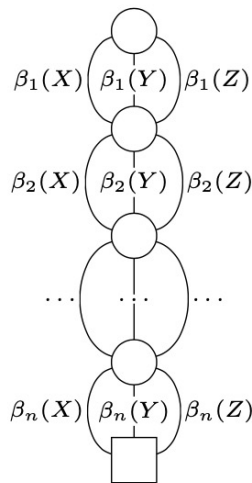


Figure 6: The Decision Diagram of LBCS

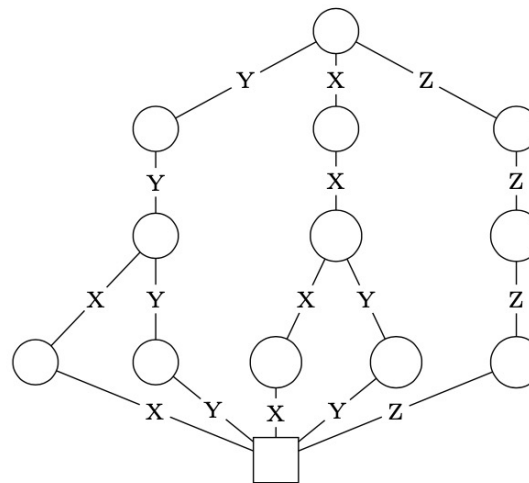


Figure 7: The unoptimised decision diagram of LDF-based Pauli Grouping of  $H_2$  (4 qubits) in Jordan-Wigner encoding

# Decision diagrams

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## Algorithm 2 Construction of a decision diagram (DD) from Hamiltonian $H$

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Take absolute values of coefficients in  $H$

Merge compatible terms to get reduced positive Pauli list  $\mathcal{R}(H)$

▷ Preprocessing

**for** Each term and coefficient in  $\mathcal{R}(H)$  **do**

▷ Initialisation of DD

    Take existing path covering the longest prefix of term

    Create new edges for remaining Pauli operators up to the last

    Create edge to terminal with the last Pauli op and coefficient as edge weight

**for** Vertex in decision diagram in breadth-first order from terminal **do**

▷ Normalisation of DD

    Calculate sum of weights on out-going edges

    Divide weights on out-going edges by sum and multiply sum to in-coming edge weights

**for** Vertex in decision diagram in breadth-first order from terminal **do**

▷ Merge equivalent vertices in DD

    Calculate hash of vertex and if equivalent vertex exists, merge both

Remove identities in DD

    Replace “lonely” identity edges with virtual edges

    Remove identity edges where other edge with same source and target exists

    Merge targets of identity edges with target vertices of other edge

**for** Vertex in decision diagram in breadth-first order from terminal **do**

▷ Merge equivalent vertices in DD

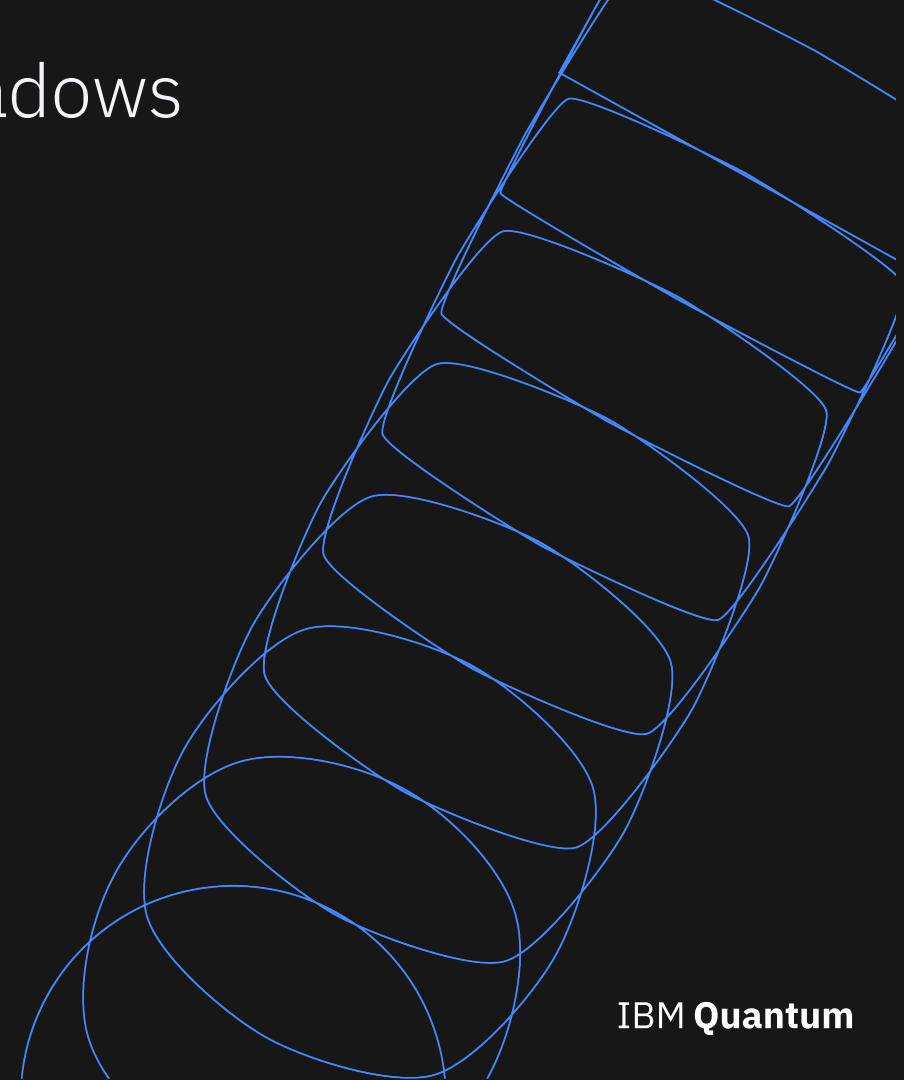
    Calculate hash of vertex and if equivalent vertex exists, merge both

**return** decision diagram

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# Solution 4: Adaptive Pauli shadows

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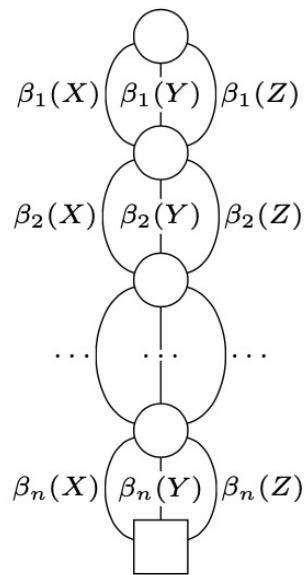


## ADAPTIVE PAULI SHADOWS FOR ENERGY ESTIMATION

CHARLES HADFIELD

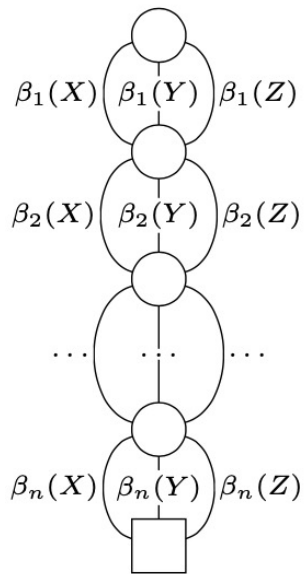
ABSTRACT. Locally-biased classical shadows allow rapid estimation of energies of quantum Hamiltonians. Recently, derandomised classical shadows have emerged claiming to be even more accurate. This accuracy comes at a cost of introducing classical computing resources into the energy estimation procedure. This present note shows, by adding a fraction of this classical computing resource to the locally-biased classical shadows setting, that the modified algorithm, termed *Adaptive Pauli Shadows* is state-of-the-art for energy estimation.

# Adaptive Pauli shadows



$$\text{cost}_{\text{diag}}(\beta) = \sum_P \alpha_P^2 \frac{1}{\prod_{i|P_i \neq I} \beta_i(P_i)}$$

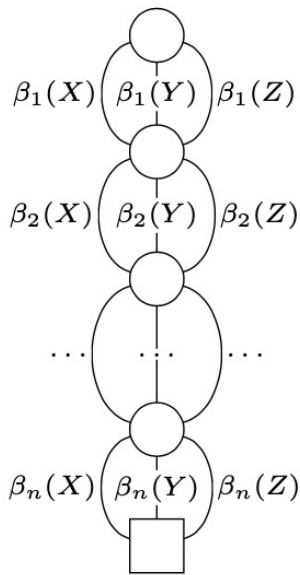
# Adaptive Pauli shadows



Let's choose each qubit's probability *on-the-fly*.

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# Adaptive Pauli shadows



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Let's start with the first qubit:

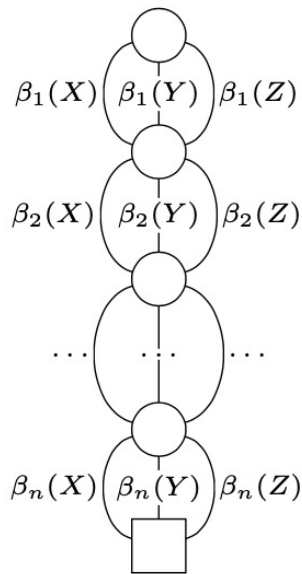
$$\text{cost}(\beta_1) = \frac{c_X}{\beta_1(X)} + \frac{c_Y}{\beta_1(Y)} + \frac{c_Z}{\beta_1(Z)}$$

Subject to beta being a probability distribution.

$$\text{cost}_{\text{diag}}(\beta) = \sum_P \alpha_P^2 \frac{1}{\prod_{i|P_i \neq I} \beta_i(P_i)}$$



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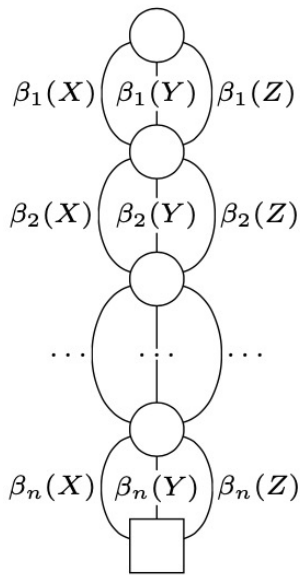
Subject to beta being a probability distribution.

*This cost function has an analytical solution!*

Pick the basis  $B_1$  from this distribution.

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# Adaptive Pauli shadows



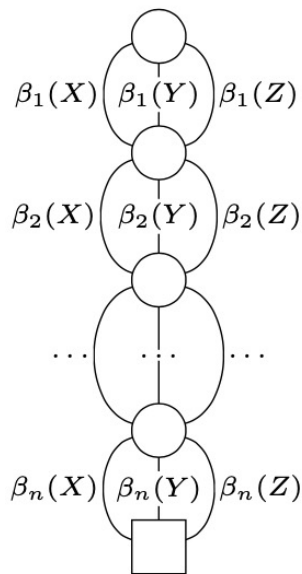
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For  $i^{\text{th}}$  qubit, only look at Pauli terms for which it is still possible to provide an estimate upon eventual measurement

$$\text{cost}(\beta_i) = \frac{c_X}{\beta_i(X)} + \frac{c_Y}{\beta_i(Y)} + \frac{c_Z}{\beta_i(Z)}$$

$$\text{cost}_{\text{diag}}(\beta) = \sum_P \alpha_P^2 \frac{1}{\prod_{i|P_i \neq I} \beta_i(P_i)}$$

# Adaptive Pauli shadows



Let's choose each qubit's probability *on-the-fly*.

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### Algorithm 2 Choice of measurement basis for Adaptive Pauli Shadows

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Randomly choose a bijection  $i : [n] \rightarrow [n]$

**for**  $j \in [n]$  **do**

Set  $\beta_{i(j)} : \mathcal{B} \rightarrow \mathbb{R}^+$  by solving the optimisation problem in Eq. (2)

Choose  $B_{i(j)}$  randomly according to distribution  $\beta_{i(j)}$

**return**  $B = \otimes_{i \in [n]} B_i$ .

---

$$\text{cost}_{\text{diag}}(\beta) = \sum_P \alpha_P^2 \frac{1}{\prod_{i|P_i \neq I} \beta_i(P_i)}$$

## Time to run VQE for some molecules

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