

Math 125 A – Fall 2013  
Midterm 2: November 4

Name:..... Solutions .....

/30

**Problem 1:** (10 points) Decide whether the following statements are True or False. Circle the right answer. You don't need to justify your answers.

T  F Once we prove compactness, the notions of elementary class and of weakly elementary class become equivalent.

T F If an embedding between two  $\mathcal{L}$ -structures is onto, it is an elementary embedding.

T  F Let  $\mathcal{L}$  be countable. Then every infinite  $\mathcal{L}$ -structure has a proper elementary substructure (proper means not equal).

T  F Let  $\mathcal{A}$  and  $\mathcal{B}$  be  $\mathcal{L}$ -structures. Then  $\mathcal{A} \preceq \mathcal{B}$  if and only if  $\mathcal{A} \subseteq \mathcal{B}$  and  $\mathcal{A} \equiv \mathcal{B}$ .

T F Let  $\mathcal{M}, \mathcal{N}$  be  $\mathcal{L}$ -structures.  $\mathcal{M}$  and  $\mathcal{N}$  are elementary equivalent if and only if  $Th(\mathcal{N}) \subseteq Th(\mathcal{M})$ .

**Problem 3:** (6 points)

Let  $\Sigma = \{\varphi_n : n \in \mathbb{N}\} \subseteq \text{Sent}_{\mathcal{L}}$  be such that  $\varphi_{n+1} \models \varphi_n$  for every  $n$ . Prove that if  $\text{Mod}(\Sigma)$  is elementary, then there exists some  $n$  such that  $\models \varphi_n \leftrightarrow \varphi_{n+1}$ .

Suppose  $\text{Mod}(\Sigma)$  is elementary and  $\text{Mod}(\Sigma) = \text{Mod}(\mathcal{T}) \Rightarrow \Sigma \models \mathcal{T}$   
Then, by Thm from class,  $\Sigma \neq \Sigma_0$  for some finite  $\Sigma_0 \in \Sigma$

Let  $n_0 = \text{largest } n, \varphi_n \in \Sigma_0$ .

Since  $\varphi_{n_0} \models \varphi_n \quad \forall n < n_0, \quad \varphi_{n_0} \models \Sigma_0 \models \Sigma$

Is  $\varphi_{n_0} \models \varphi_{n+1}$

Problem 3: (8 points) Consider  $\mathcal{Q} = (\mathbb{Q}, +)$ .

(a) Find four subsets of  $\mathbb{Q}$  which are definable in  $\mathcal{Q}$ , and give the formulas that define them.

(b) Prove that there are no other subsets of  $\mathbb{Q}$  which are definable in  $\mathcal{Q}$ .

- (a)
- $\{0\} \subseteq \mathbb{Q}$  definable by  $x+x=x$
  - $\mathbb{Q} - \{0\}$  definable by  $x+x \neq x$
  - $\mathbb{Q}$  definable by  $x=x$
  - $\emptyset$  definable by  $x \neq x$

(b) Suppose  $D \subseteq \mathbb{Q}$  is definable.

$\forall D \subseteq \{0\}$  then either  $D = \emptyset$  or  $D = \{0\}$

So, suppose  $D \not\subseteq \{0\}$  and  $d \in D - \{0\}$ .

Take any  $q \in \mathbb{Q} - \{0\}$ . The  $h(p) = \frac{q}{d}p$  is an automorphism of  $(\mathbb{Q}, +)$

So  $q = h(d) \in D$

Thus  $\mathbb{Q} - \{0\} \subseteq D$ . Hence either  $\mathbb{Q} - \{0\} = D$   
or  $\mathbb{Q} = D$

**Problem 4:** (6 points) Let  $\mathcal{L} = \{\leq\}$ , and let  $\Sigma \subset \text{Sent}_{\mathcal{L}}$  be such that all the models of  $\Sigma$  are linear orderings. Suppose also that for every  $n \in \mathbb{N}$ , there is an  $\mathcal{A} \models \Sigma$  such that  $|\mathcal{A}| \geq n$ .

Prove that for every linear ordering  $\mathcal{C} = (C, \leq^{\mathcal{C}})$  there is an  $\mathcal{A} \models \Sigma$  for which there exists an  $\mathcal{L}$ -embedding of  $\mathcal{C}$  into  $\mathcal{A}$ .

$$\text{Let } \mathcal{L}' = \{\leq\} \cup \{c_b : b \in C\}$$

$$\text{Let } \Gamma = \Sigma \cup \{c_b \leq c_{b'} : \begin{matrix} b, b' \in C \\ b \leq^{\mathcal{C}} b' \end{matrix}\}$$

$\Gamma$  is finitely satisfiable, because if  $\Gamma_0 \subseteq \Gamma$  is finite it mentions only finitely many constants  $c_{b_1}, \dots, c_{b_k}$ .

Then, let  $\mathcal{A} \models \Sigma$  with  $|\mathcal{A}| > b_k$

We can interpret  $c_{b_1}, \dots, c_{b_k}$  in  $\mathcal{A}$  so that they are in the correct ordering and get a model of  $\Gamma_0$ .

So, by compactness,  $\Gamma$  is satisfiable.

Let  $\mathcal{A}' \models \Gamma$  and let  $\mathcal{A}$  be the restriction of  $\mathcal{A}'$  to  $\mathcal{L}$ .

The  $\mathcal{A} \models \Sigma$  and the map  $h: \mathcal{C} \rightarrow \mathcal{A}$   
 $h(b) = c_b^{\mathcal{A}'}$  is an order preserving embedding.

**Problem 5:** (5 points)

Let  $\mathcal{L}$  be a countable language and  $\mathcal{M}$  an infinite  $\mathcal{L}$ -structure. Show that there exists a countable  $\mathcal{L}$ -structure  $\mathcal{N}$  such that  $\mathcal{M} \preceq \mathcal{N}$  and  $\mathcal{M} \neq \mathcal{N}$ .

Use  $\cup_p$  LST to get  $\mathcal{N}' \succeq \mathcal{M}$ ,  $\mathcal{N}'$  uncountable

Let  $X = \mathcal{M} \cup \{c\}$  for some  $c \in \mathcal{N}' \setminus \mathcal{M}$ . (note  $X$  is countable)

Use Down LST to get  $\mathcal{N}^p$ ,  $X \subseteq \mathcal{N}^p$ ,  $\mathcal{N}^p$  countable  
 $\mathcal{N}^p \preceq \mathcal{N}'$

Then  $\mathcal{M} \preceq \mathcal{N}^p$  and  $c \in \mathcal{N}^p \setminus \mathcal{M}$ .