John R. Steel University of California, Berkeley

January 2017

**Preliminaries** 

Definition of least branch hod pair

Comparison of least branch hod pairs

Problem: Analyze HOD in models of determinacy.

Conjecture 1. Assume  $AD^+ + V = L(P(\mathbb{R}))$ ; then HOD  $\models$  GCH.

*Conjecture 2.* There is  $M \models AD^+ + V = L(P(\mathbb{R}))$  such that  $HOD^M \models$  "there is a subcompact cardinal".

#### **Preliminaries**

Definition of least branch hod pair

Comparison of east branch hod pairs

Problem: Analyze HOD in models of determinacy.

Conjecture 1. Assume  $AD^+ + V = L(P(\mathbb{R}))$ ; then HOD  $\models$  GCH.

*Conjecture 2.* There is  $M \models AD^+ + V = L(P(\mathbb{R}))$  such that  $HOD^M \models$  "there is a subcompact cardinal".

## Definition

"No long extenders" (NLE) is the assertion: there is no countable, iterable pure extender mouse with a long extender on its sequence.

#### **Preliminaries**

Definition of least branch hod pair

Comparison of east branch hod pairs

### Theorem

Suppose that  $\kappa$  is supercompact, and there are arbitrarily large Woodin cardinals. Suppose that V is uniquely iterable above  $\kappa$ ; then

- (1) for any  $\Gamma \subseteq \operatorname{Hom}_{\infty}$  such that  $L(\Gamma, \mathbb{R}) \models \mathsf{NLE}$ ,  $\operatorname{HOD}^{L(\Gamma, \mathbb{R})} \models \mathsf{GCH}$ , and
- (2) there is a  $\Gamma \subseteq \operatorname{Hom}_{\infty}$  such that  $\operatorname{HOD}^{L(\Gamma,\mathbb{R})} \models$  "there is a subcompact cardinal".

#### Preliminaries

Definition of least pranch hod pair

Comparison of east branch hod pairs

### Theorem

Suppose that  $\kappa$  is supercompact, and there are arbitrarily large Woodin cardinals. Suppose that V is uniquely iterable above  $\kappa$ ; then

- (1) for any  $\Gamma \subseteq \operatorname{Hom}_{\infty}$  such that  $L(\Gamma, \mathbb{R}) \models \mathsf{NLE}$ ,  $\operatorname{HOD}^{L(\Gamma, \mathbb{R})} \models \mathsf{GCH}$ , and
- (2) there is a  $\Gamma \subseteq \operatorname{Hom}_{\infty}$  such that  $\operatorname{HOD}^{\mathcal{L}(\Gamma,\mathbb{R})} \models$  "there is a subcompact cardinal".

*Moral:* Below long extenders, there is a simple general notion of *hod pair*, and a general comparison theorem for them. They have a fine structure. *Modulo the existence of iteration strategies*, they can be used to analyze HOD, and they can have subcompact cardinals.

#### Preliminaries

Definition of least branch hod pair

Comparison of east branch hod pairs

## **A Glossary**

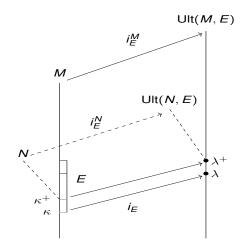
(a) An extender *E* over *M* is a system of measures on *M* coding an elementary  $i_E : M \rightarrow \text{Ult}(M, E)$ . *E* is short iff all its component measures concentrate on  $\text{crit}(i_E)$ .

$$Ult(M, E) = \{ [a, f]_E^M \mid f \in M \text{ and } a \in [\lambda]^{<\omega} \},$$
  
where  $\lambda = \lambda(E) = i_E(crit(E)).$ 

#### **Preliminaries**

Definition of least branch hod pair

Comparison of least branch hod pairs



Definition of least branch hod pair

Comparison of east branch hod pairs

Hod pair capturing and HOD.

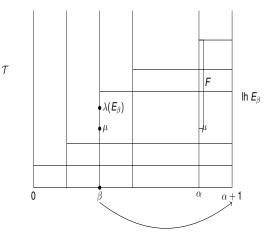
*M* agrees with Ult(M, E) and Ult(N, E) to  $(\lambda^+)^{Ult(M, E)}$ .

## (b) A normal iteration tree on M is an iteration tree T on M in which the extenders used have increasing strengths, and are applied to the longest possible initial segment of the earliest possible model. (So along branches of T, generators are not moved.)

#### **Preliminaries**

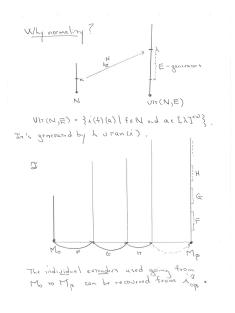
Definition of least branch hod pair

Comparison of least branch hod pairs



Definition of least branch hod pair

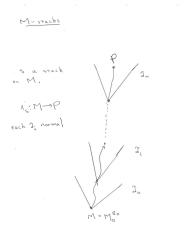
Comparison of least branch hod pairs



Definition of least branch hod pair

Comparison of least branch hod pairs

(c) An *M*-stack is a sequence s = ⟨T<sub>0</sub>,...,T<sub>n</sub>⟩ of normal trees such that T<sub>0</sub> is on *M*, and T<sub>i+1</sub> is on the last model of T<sub>i</sub>.



#### **Preliminaries**

Definition of least branch hod pair

Comparison of east branch hod pairs

(d) An *iteration strategy* Σ for *M* is a function that is defined on *M*-stacks *s* that are by Σ whose last tree has limit length, and picks a cofinal wellfounded branch of that tree.

#### **Preliminaries**

Definition of least branch hod pair

Comparison of east branch hod pairs

- (d) An *iteration strategy* Σ for *M* is a function that is defined on *M*-stacks *s* that are by Σ whose last tree has limit length, and picks a cofinal wellfounded branch of that tree.
- (e) If *s* is an *M*-stack, then  $\Sigma_s$  is the *tail strategy* given by  $\Sigma_s(t) = \Sigma(s^{-}t)$ .

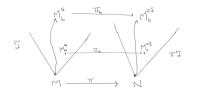
Definition of least branch hod pair

Comparison of east branch hod pairs

- (d) An *iteration strategy* Σ for *M* is a function that is defined on *M*-stacks *s* that are by Σ whose last tree has limit length, and picks a cofinal wellfounded branch of that tree.
- (e) If *s* is an *M*-stack, then  $\Sigma_s$  is the *tail strategy* given by  $\Sigma_s(t) = \Sigma(s^{-}t)$ .
- (f) It  $\pi: M \to N$  is elementary, and  $\Sigma$  is an iteration strategy for *N*, then  $\Sigma^{\pi}$  is the *pullback strategy* given by:  $\Sigma^{\pi}(s) = \Sigma(\pi s)$ .

Definition of least branch hod pair

Comparison of east branch hod pairs



if 
$$b = \Sigma(\pi I)$$
  
then  $\Sigma^{\pi}(I) = b$ 

Definition of least pranch hod pair

Comparison of east branch hod pairs

## Definition

A least branch premouse (lpm) is a structure  $\mathcal{M}$  constructed from a coherent sequence  $\dot{E}^{\mathcal{M}}$  of extenders, and a predicate  $\dot{\Sigma}^{\mathcal{M}}$  for an iteration strategy for  $\mathcal{M}$ .

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

## Definition

A *least branch premouse* (lpm) is a structure  $\mathcal{M}$  constructed from a coherent sequence  $\dot{\mathcal{E}}^{\mathcal{M}}$  of extenders, and a predicate  $\dot{\Sigma}^{\mathcal{M}}$  for an iteration strategy for  $\mathcal{M}$ .

## Remarks

(a)  $\mathcal{M}$  has a hierarchy, and a fine structure. By convention, there is a  $k = k(\mathcal{M})$  such that  $\mathcal{M}$  is *k*-sound. (I.e.,  $\mathcal{M} = \text{Hull}_k(\rho_k^{\mathcal{M}} \cup \rho_k^{\mathcal{M}})$ .)

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

## Definition

A *least branch premouse* (lpm) is a structure  $\mathcal{M}$  constructed from a coherent sequence  $\dot{\mathcal{E}}^{\mathcal{M}}$  of extenders, and a predicate  $\dot{\Sigma}^{\mathcal{M}}$  for an iteration strategy for  $\mathcal{M}$ .

## Remarks

- (a)  $\mathcal{M}$  has a hierarchy, and a fine structure. By convention, there is a  $k = k(\mathcal{M})$  such that  $\mathcal{M}$  is *k*-sound. (I.e.,  $\mathcal{M} = \text{Hull}_k(\rho_k^{\mathcal{M}} \cup p_k^{\mathcal{M}}).$ )
- (b) We use Jensen indexing for the extenders in  $E^{\mathcal{M}}$ .

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

## Definition

A *least branch premouse* (lpm) is a structure  $\mathcal{M}$  constructed from a coherent sequence  $\dot{E}^{\mathcal{M}}$  of extenders, and a predicate  $\dot{\Sigma}^{\mathcal{M}}$  for an iteration strategy for  $\mathcal{M}$ .

## Remarks

- (a)  $\mathcal{M}$  has a hierarchy, and a fine structure. By convention, there is a  $k = k(\mathcal{M})$  such that  $\mathcal{M}$  is *k*-sound. (I.e.,  $\mathcal{M} = \text{Hull}_k(\rho_k^{\mathcal{M}} \cup \rho_k^{\mathcal{M}})$ .)
- (b) We use Jensen indexing for the extenders in  $\dot{E}^{\mathcal{M}}$ .
- (c) At strategy-active stages  $\alpha$ , we consider the  $\mathcal{M}|\alpha$ -least  $\langle \nu, k, \mathcal{T} \rangle$  such that  $\mathcal{T}$  is a normal tree of limit length on  $\mathcal{M}|\langle \nu, k \rangle$  that is by  $\dot{\Sigma}^{\mathcal{M}|\alpha}$ , and  $\dot{\Sigma}^{\mathcal{M}|\alpha}(\mathcal{T})$  is undefined. Then  $\dot{\Sigma}^{\mathcal{M}|(\alpha+1)} = \dot{\Sigma}^{\mathcal{M}|\alpha} \cup \{\langle \nu, k, \mathcal{T}, b \rangle\}$ , where *b* is some cofinal branch of  $\mathcal{T}$ .

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

A *least branch hod pair* (lbr hod pair) with with scope Z is a pair  $(P, \Sigma)$  such that

- (1) P is an lpm,
- (2)  $\Sigma$  is an iteration strategy defined on all *P*-stacks  $s \in Z$ ,

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

A *least branch hod pair* (lbr hod pair) with with scope Z is a pair  $(P, \Sigma)$  such that

- (1) P is an lpm,
- (2)  $\Sigma$  is an iteration strategy defined on all *P*-stacks  $s \in Z$ ,
- (3) if *Q* is a  $\Sigma$ -iterate of *P* via *s*, then  $\dot{\Sigma}^{Q} \subseteq \Sigma_{s}$ , and

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

A *least branch hod pair* (lbr hod pair) with with scope Z is a pair  $(P, \Sigma)$  such that

- (1) P is an lpm,
- (2)  $\Sigma$  is an iteration strategy defined on all *P*-stacks  $s \in Z$ ,
- (3) if *Q* is a  $\Sigma$ -iterate of *P* via *s*, then  $\dot{\Sigma}^{Q} \subseteq \Sigma_{s}$ , and
- (4) Σ is self-consistent, normalizes well, and has strong hull condensation.

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

A *least branch hod pair* (lbr hod pair) with with scope Z is a pair  $(P, \Sigma)$  such that

- (1) P is an lpm,
- (2)  $\Sigma$  is an iteration strategy defined on all *P*-stacks  $s \in Z$ ,
- (3) if *Q* is a  $\Sigma$ -iterate of *P* via *s*, then  $\dot{\Sigma}^{Q} \subseteq \Sigma_{s}$ , and
- (4) Σ is self-consistent, normalizes well, and has strong hull condensation.

 $\Sigma$  is *self-consistent* iff the part of  $\Sigma$  that is a strategy for  $\mathcal{M}|\langle \nu, k \rangle$  is consistent with the part of  $\Sigma$  that is a strategy for  $\mathcal{M}|\langle \mu, j \rangle$ .

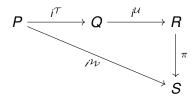
#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of east branch hod pairs

## Normalizing well

For  $\langle T, U \rangle$  a stack on *P*, and W = W(T, U) its embedding normalization, we have



#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

Hod pair capturing and HOD.

Then  $\Sigma$  2-normalizes well iff

 $\langle \mathcal{T}, \mathcal{U} \rangle$  is by  $\Sigma$  iff  $W(\mathcal{T}, \mathcal{U})$  is by  $\Sigma$ ,

and

$$\Sigma^{\pi}_{\langle \mathcal{W} \rangle} = \Sigma_{\langle \mathcal{T}, \mathcal{U} \rangle}.$$

for all such stacks  $\langle T, U \rangle$ .  $\Sigma$  normalizes well iff all its tails 2-normalize well.

# W(E,F)

Let  $\mathcal{T} = \langle E \rangle$  and  $\mathcal{U} = \langle F \rangle$ , with  $\operatorname{crit}(F) < \operatorname{crit}(E)$ .

$$N \xrightarrow{F} Q \xrightarrow{\tau} i_{F}^{M}(N) = \text{Ult}_{0}(P, i_{F}^{M}(E))$$

$$\stackrel{F}{\longrightarrow} P \xrightarrow{V_{F}^{M}(E)}$$

Preliminaries

#### Definition of least branch hod pair

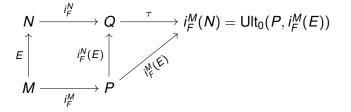
Comparison of least branch hod pairs

Hod pair capturing and HOD.

 $\tau$  is the natural embedding from Ult(*N*, *F*) to  $i_F^M(N)$ . That is,

$$\tau([a,g]_F^N)=[a,g]_F^M.$$

The extenders used in W(E, F) are *E*, then *F*, then  $i_F^M(E)$ .



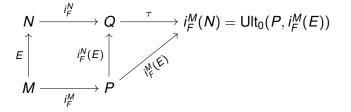
The full normalization X(E, F) uses E, then F, then  $i_F^N(E)$ .

#### Preliminaries

#### Definition of least branch hod pair

Comparison of east branch hod pairs

The extenders used in W(E, F) are *E*, then *F*, then  $i_F^M(E)$ .



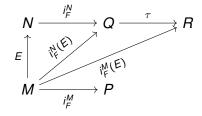
The full normalization X(E, F) uses E, then F, then  $i_F^N(E)$ .Ult(M, F) has  $i_F^N(E)$  on its sequence by Condensation.

#### Preliminaries

#### Definition of least branch hod pair

Comparison of east branch hod pairs

The situation when  $\operatorname{crit}(E) < \operatorname{crit}(F) < \lambda(E)$  is summarized by:



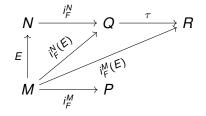
W(E, F) uses *E*, then *F*, then  $i_F^M(E)$ .

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

The situation when  $\operatorname{crit}(E) < \operatorname{crit}(F) < \lambda(E)$  is summarized by:



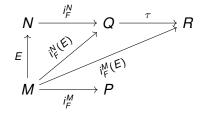
W(E, F) uses *E*, then *F*, then  $i_F^M(E).X(E, F)$ ) uses *E*, then *F*, then  $i_F^N(E)$ .

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

The situation when  $\operatorname{crit}(E) < \operatorname{crit}(F) < \lambda(E)$  is summarized by:



W(E, F) uses *E*, then *F*, then  $i_F^M(E).X(E, F)$ ) uses *E*, then *F*, then  $i_F^N(E)$ .

*Remark* So there are two ways *F* can appear in the branch to the final model: as itself, or buried inside  $i_F(E)$ .

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

# $W(\mathcal{T}, F)$

Let  $\mathcal{T}$  have last model  $M_{\theta}^{\mathcal{T}}$ , with F on its sequence. Let

 $\alpha = \text{least } \xi \text{ such that } F \text{ is on the } M_{\varepsilon}^{T} \text{-sequence},$ 

and

$$\beta = \text{least } \xi \text{ such that } \operatorname{crit}(F) < \lambda(E_{\xi}^{\mathcal{T}}).$$

Then

$$W(\mathcal{T}, F) = \mathcal{T} \upharpoonright (\alpha + 1)^{\frown} \langle F \rangle^{\frown} \Phi "\mathcal{T} \upharpoonright (\theta + 1 - \beta).$$

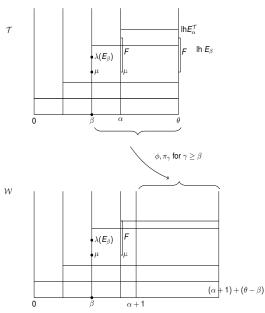
Here  $\Phi$  comes from a copying construction.

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of east branch hod pairs

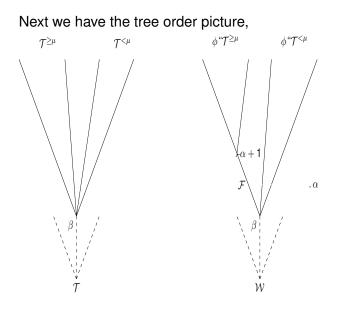
 $W(\mathcal{T},F)$ 



#### **Preliminaries**

#### Definition of least branch hod pair

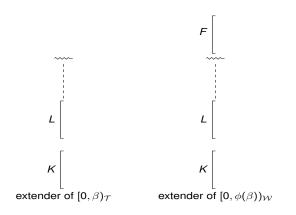
Comparison of least branch hod pairs



#### Definition of least branch hod pair

Comparison of least branch hod pairs

# We show how *F* gets inserted into the extender of the branch $\mathcal{T}$ ending at $M_{\xi}^{\mathcal{T}}$ . For $\xi = \beta$ :

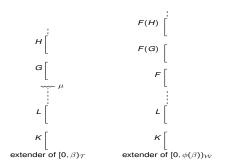


#### Preliminaries

#### Definition of least branch hod pair

Comparison of east branch hod pairs

For  $\xi > \beta$ , let *G* be the first extender used in  $[0, \xi)_T$  such that  $\nu(G) \ge \nu(E_{\beta}^T)$ . The picture depends on whether  $\mu \le \operatorname{crit}(G)$ . If  $\mu \le \operatorname{crit}(G)$ , it is



#### Preliminaries

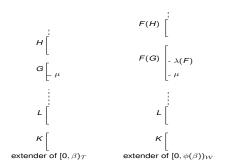
#### Definition of least branch hod pair

Comparison of least branch hod pairs

Hod pair capturing and HOD.

In this case, *F* is used on  $[0, \phi(\xi))_W$ , and the remaining extender used are the images of old ones under copy maps.

## If $\operatorname{crit}(G) < \mu < \nu(G)$ , the picture is



#### Preliminaries

#### Definition of least branch hod pair

Comparison of least branch hod pairs

Hod pair capturing and HOD.

In this case, the two branches use the same extenders until *G* is used on  $[0, \xi)_T$ . At that point and after,  $[0, \phi(\xi))_W$  uses the images of extenders under the copy maps.

# $W(\mathcal{T},\mathcal{U})$

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

Hod pair capturing and HOD.

We define  $W_{\gamma} = W(\mathcal{T}, \mathcal{U} \upharpoonright \gamma + 1)$  by induction on  $\gamma$ . It has last model  $R_{\gamma}$ , and we have  $\sigma_{\gamma}$  from  $M_{\gamma}^{\mathcal{U}}$  to  $R_{\gamma}$ .

# $W(\mathcal{T},\mathcal{U})$

#### **Preliminaries**

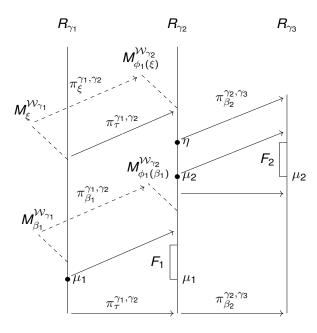
#### Definition of least branch hod pair

Comparison of least branch hod pairs

Hod pair capturing and HOD.

We define  $W_{\gamma} = W(\mathcal{T}, \mathcal{U} \upharpoonright \gamma + 1)$  by induction on  $\gamma$ . It has last model  $R_{\gamma}$ , and we have  $\sigma_{\gamma}$  from  $M_{\gamma}^{\mathcal{U}}$  to  $R_{\gamma}$ .

The  $W_{\gamma}$ 's constitute a tree of iteration trees, under the order  $<_U$  on the  $\gamma$ 's. If  $\gamma_1 <_U \gamma_2 <_U \gamma_3$ , the picture is:

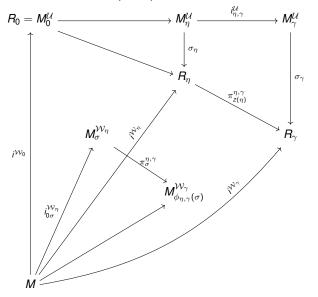


#### **Preliminaries**

## Definition of least branch hod pair

Comparison of least branch hod pairs

Another view of  $W(\mathcal{T}, \mathcal{U})$ :



#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

# Strong hull condensation

Roughly,  $\Sigma$  has *strong hull condensation* iff  $\mathcal{T}$  and  $\mathcal{U}$  are normal trees on P, and  $\mathcal{U}$  is by  $\Sigma$ , and  $\pi : \mathcal{T} \to \mathcal{U}$  is appropriately elementary, then  $\mathcal{T}$  is by  $\Sigma$ .

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

# Strong hull condensation

Roughly,  $\Sigma$  has *strong hull condensation* iff  $\mathcal{T}$  and  $\mathcal{U}$  are normal trees on P, and  $\mathcal{U}$  is by  $\Sigma$ , and  $\pi \colon \mathcal{T} \to \mathcal{U}$  is appropriately elementary, then  $\mathcal{T}$  is by  $\Sigma$ . One must be careful about the elementarity required of  $\pi$ , and in particular, the extent to which  $\pi$  is required to preserve exit extenders. There are several possible condensation properties here: hull condensation (Sargsyan), strong hull condensation, and still stronger ones.

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

Strong hull condensation means condensing under *psuedo-hull embeddings*. The natural embedding of  $\mathcal{T}$  into  $W(\mathcal{T}, \mathcal{U})$  is an example of a psuedo-hull embedding.

## Definition

A pseudo-hull embedding of  ${\mathcal T}$  into  ${\mathcal U}$  is a system

$$\langle \boldsymbol{u}, \langle \boldsymbol{t}^{\boldsymbol{0}}_{\boldsymbol{\beta}} \mid \boldsymbol{\beta} < \ln \mathcal{T} \rangle, \langle \boldsymbol{t}^{\boldsymbol{1}}_{\boldsymbol{\beta}} \mid \boldsymbol{\beta} + \mathbf{1} < \ln \mathcal{T} \rangle, \boldsymbol{p} \rangle$$

with various properties, including:

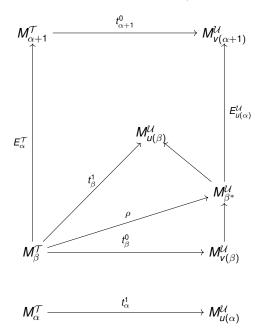
$$egin{aligned} \mathcal{P}(\mathcal{E}_{lpha}^{\mathcal{T}}) &= t_{lpha}^1(\mathcal{E}_{lpha}^{\mathcal{T}}) \ &= \mathcal{E}_{u(lpha)}^{\mathcal{U}}. \end{aligned}$$

## Preliminaries

## Definition of least branch hod pair

Comparison of least branch hod pairs

The diagram related to successor steps in T is:



#### Preliminaries

#### Definition of least branch hod pair

Comparison of least branch hod pairs

# Elementary properties of lbr hod pairs

## Lemma

Let  $(P, \Sigma)$  be an lbr hod pair with scope Z, and suppose  $\pi: Q \to P$  is elementary; then  $(Q, \Sigma^{\pi})$  is an lbr hod pair with scope Z.

#### Preliminaries

#### Definition of least branch hod pair

Comparison of east branch hod pairs

# Elementary properties of lbr hod pairs

## Lemma

Let  $(P, \Sigma)$  be an lbr hod pair with scope Z, and suppose  $\pi: Q \to P$  is elementary; then  $(Q, \Sigma^{\pi})$  is an lbr hod pair with scope Z.

## Lemma

(Pullback consistency) Let  $(P, \Sigma)$  be an lbr hod pair with scope Z, and let s be a P-stack by  $\Sigma$  giving rise to the iteration map  $\pi: P \to Q$ ; then  $(\Sigma_s)^{\pi} = \Sigma$ .

#### Preliminaries

#### Definition of least branch hod pair

Comparison of least branch hod pairs

# Elementary properties of lbr hod pairs

## Lemma

Let  $(P, \Sigma)$  be an lbr hod pair with scope Z, and suppose  $\pi: Q \to P$  is elementary; then  $(Q, \Sigma^{\pi})$  is an lbr hod pair with scope Z.

## Lemma

(Pullback consistency) Let  $(P, \Sigma)$  be an lbr hod pair with scope Z, and let s be a P-stack by  $\Sigma$  giving rise to the iteration map  $\pi: P \to Q$ ; then  $(\Sigma_s)^{\pi} = \Sigma$ .

## Lemma

(Dodd-Jensen) The  $\Sigma$ -iteration map from  $(P, \Sigma)$  to  $(Q, \Psi)$  is pointwise a pointwise minimal embedding of  $(P, \Sigma)$  into  $(Q, \Psi)$ .

#### **Preliminaries**

#### Definition of least branch hod pair

Comparison of least branch hod pairs

# Comparison

## **Theorem (Comparison)**

Assume AD<sup>+</sup>, and let  $(P, \Sigma)$  and  $(Q, \Psi)$  be lbr hod pairs with scope HC; then there are normal trees T on P by  $\Sigma$ and U on Q by  $\Psi$  with last models R and S respectively, such that either

(1) 
$$R \trianglelefteq S$$
, and  $\Sigma_T \subseteq \Psi_U$ , or

(2) 
$$S \subseteq R$$
, and  $\Psi_{\mathcal{U}} \subseteq \Sigma_{\mathcal{T}}$ .

## Corollary

Assume  $AD^+$ ; then the mouse order  $\leq^*$  on lbr hod pairs with scope HC is a prewellorder.

Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

# *Proof of theorem.* Let $N^*$ be a countable, Γ-correct model with a Woodin cardinal, where $(P, \Sigma)$ and $(Q, \Psi)$ are in Γ.

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

(a) no extenders on the N side are used,

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

- (a) no extenders on the N side are used, and
- (b) no strategy disagreements show up.

That  $\boldsymbol{\Sigma}$  normalizes well and has strong hull condensation are crucial here.

#### **Preliminaries**

Definition of least branch hod pair

Comparison of least branch hod pairs

- (a) no extenders on the N side are used, and
- (b) no strategy disagreements show up.

That  $\Sigma$  normalizes well and has strong hull condensation are crucial here.

Since  $N^*$  has a Woodin cardinal,  $(P, \Sigma)$  cannot iterate past all such  $(N, \Omega)$ , and hence, some such  $(N, \Omega)$  is an iterate of  $(P, \Sigma)$ . Similarly for  $(Q, \Psi)$ , and we are done.

Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

Why are there no strategy disagreements?

Suppose we have produced an iteration tree  $\mathcal{T}$  on P with last model R, and that  $R|\alpha = N|\alpha$ , and that  $\mathcal{U}$  is a tree on  $R|\alpha = N|\alpha$  played by both  $\Sigma_{\mathcal{T},R|\alpha}$  (the tail of  $\Sigma$ ) and  $\Omega$ , the  $N^*$ -induced strategy for N. Let  $\mathcal{U}$  have limit length, and let  $b = \Omega(\mathcal{U})$ . We must see  $b = \Sigma(\langle \mathcal{T}, \mathcal{U} \rangle)$ .

For this, we look at the embedding normalization W(T, U) of  $\langle T, U \rangle$ , which also has limit length. Then

Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

for  $b = \Omega(\mathcal{T})$ :

- b generates (modulo *T*) a unique cofinal branch *a* of *W*(*T*,*U*).
- (ii) Letting *i*<sup>\*</sup><sub>b</sub>: *N*<sup>\*</sup> → *N*<sup>\*</sup><sub>b</sub> come from lifting *i*<sup>U</sup><sub>b</sub> to *N*<sup>\*</sup> via the iteration-strategy construction, *W*(*T*,*U*)<sup>^</sup>⟨*a*⟩ is a pseudo-hull of *i*<sup>\*</sup><sub>b</sub>(*T*).
- (iii) But  $i_b^*(\Sigma) \subseteq \Sigma$  because  $\Sigma$  was Suslin-co-Suslin captured by  $N^*$ , so  $i_b^*(\mathcal{T})$  is by  $\Sigma$ .
- (iv) Thus  $W(\mathcal{T}, \mathcal{U})^{\frown} \langle a \rangle$  is by  $\Sigma$ , because  $\Sigma$  has strong hull condensation.
- (v) But *a* determines *b*, so since  $\Sigma$  normalizes well,  $\Sigma(\langle T, U \rangle) = b$ , as desired.

Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

Phalanx comparisons work too. From this we get

## Theorem

Assume  $AD^+$ , and let  $(P, \Sigma)$  be an lbr hod pair with scope HC; then the standard parameter of P is solid and universal, and hence  $(P, \Sigma)$  has a core.

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

Phalanx comparisons work too. From this we get

## Theorem

Assume  $AD^+$ , and let  $(P, \Sigma)$  be an lbr hod pair with scope HC; then the standard parameter of P is solid and universal, and hence  $(P, \Sigma)$  has a core.

## Theorem

Assume  $AD^+$ , and let N be a countable, iterable, coarse  $\Gamma$ -Woodin model; then the hod pair construction of N does not break down.

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

Phalanx comparisons work too. From this we get

## Theorem

Assume  $AD^+$ , and let  $(P, \Sigma)$  be an lbr hod pair with scope *HC*; then the standard parameter of *P* is solid and universal, and hence  $(P, \Sigma)$  has a core.

## Theorem

Assume  $AD^+$ , and let N be a countable, iterable, coarse  $\Gamma$ -Woodin model; then the hod pair construction of N does not break down.

## Theorem

Suppose that V is uniquely iterable, and there are arbitrarliy large Woodin cardinals; then the hod pair construction of V does not break down. Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

Phalanx comparisons also yield Condensation, and

## Theorem

(Trang, S.) Assume AD<sup>+</sup>, and let  $(P, \Sigma)$  be an lbr hod pair with scope HC;  $P \models \forall \kappa (\Box_{\kappa} \Leftrightarrow \kappa \text{ is not subcompact}).$ 

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

Phalanx comparisons also yield Condensation, and

## Theorem

(Trang, S.) Assume AD<sup>+</sup>, and let  $(P, \Sigma)$  be an lbr hod pair with scope HC;  $P \models \forall \kappa (\Box_{\kappa} \Leftrightarrow \kappa \text{ is not subcompact}).$ 

Phalanx comparisons also give

## Theorem

Assume  $AD^+$ , and let  $(P, \Sigma)$  be an lbr hod pair with scope HC; then

- (1)  $\Sigma$  is positional,
- (2)  $\Sigma$  has very strong hull condensation, and
- (3)  $\Sigma$  fully normalizes well.

Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

# Hod pair capturing

Hod pairs can be used to compute HOD, provided that there are enough of them.

## Definition

(AD<sup>+</sup>) HOD *pair capturing* (HPC) is the statement: for every Suslin, co-Suslin set of reals *A*, there is an lbr hod pair ( $P, \Sigma$ ) with scope HC such that *A* is Wadge reducible to Code( $\Sigma$ ).

*Remark.* Under AD<sup>+</sup>, if  $(P, \Sigma)$  is an lbr pair with scope HC, then Code $(\Sigma)$  is Suslin and co-Suslin.

## Preliminaries

Definition of least branch hod pair

Comparison of east branch hod pairs

# Hod pair capturing

Hod pairs can be used to compute HOD, provided that there are enough of them.

## Definition

(AD<sup>+</sup>) HOD *pair capturing* (HPC) is the statement: for every Suslin, co-Suslin set of reals *A*, there is an lbr hod pair ( $P, \Sigma$ ) with scope HC such that *A* is Wadge reducible to Code( $\Sigma$ ).

*Remark.* Under AD<sup>+</sup>, if  $(P, \Sigma)$  is an lbr pair with scope HC, then Code $(\Sigma)$  is Suslin and co-Suslin.

## Theorem

Assume there is a supercompact cardinal, and arbitrarily large Woodin cardinals. Suppose V is uniquely iterable. Let  $\Gamma \subseteq \operatorname{Hom}_{\infty}$  be such that  $L(\Gamma, \mathbb{R}) \models \mathsf{NLE}$ ; then  $L(\Gamma, \mathbb{R}) \models \mathsf{HPC}$ .

## Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

Assume  $AD^+ + V = L(P(\mathbb{R})) + HPC$ ; then  $HOD | \theta$  is an *lpm. Thus*  $HOD \models GCH$ .

*Remark.* Under  $AD_{\mathbb{R}} + HPC$ ,  $HOD \mid \theta$  is the direct limit of all "full" lbr hod pairs with scope HC.

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

Assume  $AD^+ + V = L(P(\mathbb{R})) + HPC$ ; then  $HOD | \theta$  is an *lpm. Thus*  $HOD \models GCH$ .

*Remark.* Under  $AD_{\mathbb{R}} + HPC$ ,  $HOD | \theta$  is the direct limit of all "full" lbr hod pairs with scope HC.

## Theorem

Assume  $AD^+ + V = L(P(\mathbb{R})) + HPC$ ; then equivalent are:

(a)  $\delta$  is a cutpoint Woodin cardinal of HOD,

(b)  $\delta = \theta_0$ , or  $\delta = \theta_{\alpha+1}$  for some  $\alpha$ .

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

Assume  $AD^+ + V = L(P(\mathbb{R})) + HPC$ ; then  $HOD | \theta$  is an *lpm. Thus*  $HOD \models GCH$ .

*Remark.* Under  $AD_{\mathbb{R}} + HPC$ ,  $HOD | \theta$  is the direct limit of all "full" lbr hod pairs with scope HC.

## Theorem

Assume  $AD^+ + V = L(P(\mathbb{R})) + HPC$ ; then equivalent are:

(a)  $\delta$  is a cutpoint Woodin cardinal of HOD,

(b) 
$$\delta = \theta_0$$
, or  $\delta = \theta_{\alpha+1}$  for some  $\alpha$ .

Thus  $\theta_0$  is the least Woodin cardinal of HOD.

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

Assume  $AD^+ + V = L(P(\mathbb{R})) + HPC$ ; then  $HOD | \theta$  is an *lpm. Thus*  $HOD \models GCH$ .

*Remark.* Under  $AD_{\mathbb{R}} + HPC$ ,  $HOD | \theta$  is the direct limit of all "full" lbr hod pairs with scope HC.

## Theorem

Assume  $AD^+ + V = L(P(\mathbb{R})) + HPC$ ; then equivalent are:

(a)  $\delta$  is a cutpoint Woodin cardinal of HOD,

(b)  $\delta = \theta_0$ , or  $\delta = \theta_{\alpha+1}$  for some  $\alpha$ .

Thus  $\theta_0$  is the least Woodin cardinal of HOD.

*Remark.* Woodin showed  $\theta_0$  and the  $\theta_{\alpha+1}$  are Woodin in HOD. He proved an approximation to their being cutpoints.

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

## $\label{eq:conjecture.} \mbox{(AD}^+ + \mbox{NLE}) \Rightarrow \mbox{HPC}.$

**Preliminaries** 

Definition of least branch hod pair

Comparison of east branch hod pairs

## **Conjecture.** $(AD^+ + NLE) \Rightarrow HPC.$

*Remark.* HPC is a cousin of Sargsyan's "Generation of full pointclasses". It holds in the minimal model of  $AD_{\mathbb{R}} + \theta$  is regular, and somewhat beyond, by Sargsyan's work.

#### Preliminaries

Definition of least branch hod pair

Comparison of east branch hod pairs

**Conjecture.**  $(AD^+ + NLE) \Rightarrow HPC.$ 

*Remark.* HPC is a cousin of Sargsyan's "Generation of full pointclasses". It holds in the minimal model of  $AD_{\mathbb{R}} + \theta$  is regular, and somewhat beyond, by Sargsyan's work.

HPC localizes:

## Theorem

```
Assume AD^+ + HPC, and let \Gamma \subseteq P(\mathbb{R}); then L(\Gamma, \mathbb{R}) \models HPC.
```

The key to localization of HPC is to compute optimal Suslin representations for the iteration strategies in lbr hod pairs. Preliminaries

Definition of least branch hod pair

Comparison of east branch hod pairs

# Hod pairs vs. Suslin cardinals

## **Definition**

(AD<sup>+</sup>) For  $(P, \Sigma)$  an lbr hod pair with scope HC,  $M_{\infty}(P, \Sigma)$  is the direct limit of all nondropping  $\Sigma$ -iterates of *P*, under the maps given by comparisons. Preliminaries

Definition of least branch hod pair

Comparison of east branch hod pairs

# Hod pairs vs. Suslin cardinals

## Definition

 $(AD^+)$  For  $(P, \Sigma)$  an lbr hod pair with scope HC,  $M_{\infty}(P, \Sigma)$  is the direct limit of all nondropping  $\Sigma$ -iterates of *P*, under the maps given by comparisons.

 $M_{\infty}(P, \Sigma)$  is well-defined by the Dodd-Jensen lemma. Moreover, it is OD from the rank of  $(P, \Sigma)$  in the mouse order. Thus  $M_{\infty}(P, \Sigma) \in \text{HOD}$ .

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

# Hod pairs vs. Suslin cardinals

## Definition

 $(AD^+)$  For  $(P, \Sigma)$  an lbr hod pair with scope HC,  $M_{\infty}(P, \Sigma)$  is the direct limit of all nondropping  $\Sigma$ -iterates of *P*, under the maps given by comparisons.

 $M_{\infty}(P, \Sigma)$  is well-defined by the Dodd-Jensen lemma. Moreover, it is OD from the rank of  $(P, \Sigma)$  in the mouse order. Thus  $M_{\infty}(P, \Sigma) \in \text{HOD.It}$  is an initial segment of the lpm hierarchy of HOD *if*  $(P, \Sigma)$  is "full".

#### Preliminaries

Definition of least branch hod pair

Comparison of east branch hod pairs

# Hod pairs vs. Suslin cardinals

### Definition

 $(AD^+)$  For  $(P, \Sigma)$  an lbr hod pair with scope HC,  $M_{\infty}(P, \Sigma)$  is the direct limit of all nondropping  $\Sigma$ -iterates of *P*, under the maps given by comparisons.

 $M_{\infty}(P, \Sigma)$  is well-defined by the Dodd-Jensen lemma. Moreover, it is OD from the rank of  $(P, \Sigma)$  in the mouse order. Thus  $M_{\infty}(P, \Sigma) \in \text{HOD.It}$  is an initial segment of the lpm hierarchy of HOD *if*  $(P, \Sigma)$  is "full". A tree  $\mathcal{T}$  by  $\Sigma$  is  $M_{\infty}$ -*relevant* iff there is a normal  $\mathcal{U}$  by  $\Sigma$ extending  $\mathcal{T}$  with last model Q such that the branch P-to-Q does not drop.  $\Sigma^{\text{rel}}$  is the restriction of  $\Sigma$  to  $M_{\infty}$ -relevant trees.

### **Preliminaries**

Definition of least branch hod pair

Comparison of least branch hod pairs

Recall that *A* is  $\kappa$ -Suslin iff A = p[T] for some tree *T* on  $\omega \times \kappa$ .

### Theorem

(AD<sup>+</sup>) Let  $(P, \Sigma)$  be an lbr hod pair with scope HC; then Code $(\Sigma^{rel})$  is  $\kappa$ -Suslin, for  $\kappa = |M_{\infty}(P, \Sigma)|$ .

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

Recall that *A* is  $\kappa$ -Suslin iff A = p[T] for some tree *T* on  $\omega \times \kappa$ .

### Theorem

(AD<sup>+</sup>) Let  $(P, \Sigma)$  be an lbr hod pair with scope HC; then  $Code(\Sigma^{rel})$  is  $\kappa$ -Suslin, for  $\kappa = |M_{\infty}(P, \Sigma)|$ . Remark. Code $(\Sigma^{rel})$  is not  $\alpha$ -Suslin, for any  $\alpha < |M_{\infty}(P, \Sigma)|$ , by Kunen-Martin. So  $|M_{\infty}(P, \Sigma)|$  is a Suslin cardinal.

#### **Preliminaries**

Definition of least branch hod pair

Comparison of least branch hod pairs

# *Proof sketch.* $M_{\infty}(P, \Sigma)$ is the direct limit along a generic stack *s* of trees by $\Sigma$ .

#### **Preliminaries**

Definition of least branch hod pair

Comparison of least branch hod pairs

*Proof sketch.*  $M_{\infty}(P, \Sigma)$  is the direct limit along a generic stack *s* of trees by  $\Sigma$ .But *s* can be fully normalized, so there is a single normal tree  $\mathcal{W}$  on *P* with last model  $M_{\infty}(P, \Sigma)$  such that every countable "weak hull" of  $\mathcal{W}$  is by  $\Sigma$ .

#### Preliminaries

Definition of least branch hod pair

Comparison of east branch hod pairs

*Proof sketch.*  $M_{\infty}(P, \Sigma)$  is the direct limit along a generic stack *s* of trees by  $\Sigma$ .But *s* can be fully normalized, so there is a single normal tree  $\mathcal{W}$  on *P* with last model  $M_{\infty}(P, \Sigma)$  such that every countable "weak hull" of  $\mathcal{W}$  is by  $\Sigma$ .But then for  $\mathcal{T}$  countable and  $M_{\infty}$ -relevant,

 $\mathcal{T}$  is by  $\Sigma \Leftrightarrow \mathcal{T}$  is a weak hull of  $\mathcal{W}$ .

The right-to-left direction follows from very strong hull condensation for  $\Sigma$ .

### **Preliminaries**

Definition of least branch hod pair

Comparison of least branch hod pairs

*Proof sketch.*  $M_{\infty}(P, \Sigma)$  is the direct limit along a generic stack *s* of trees by  $\Sigma$ .But *s* can be fully normalized, so there is a single normal tree  $\mathcal{W}$  on *P* with last model  $M_{\infty}(P, \Sigma)$  such that every countable "weak hull" of  $\mathcal{W}$  is by  $\Sigma$ .But then for  $\mathcal{T}$  countable and  $M_{\infty}$ -relevant,

 $\mathcal{T}$  is by  $\Sigma \Leftrightarrow \mathcal{T}$  is a weak hull of  $\mathcal{W}$ .

The right-to-left direction follows from very strong hull condensation for  $\Sigma$ .

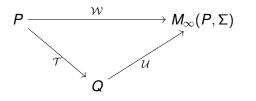
For left-to-right direction, we may assume  $\mathcal{T}$  has last model Q, and P-to-Q does not drop. We then have a normal  $\mathcal{U}$  on Q with last model  $M_{\infty}(P, \Sigma)$  such that all countable weak hulls of  $\mathcal{U}$  are by  $\Sigma$ .

### **Preliminaries**

Definition of least branch hod pair

Comparison of least branch hod pairs

We have



Then

$$\mathcal{W} = X(\mathcal{T}, \mathcal{U})$$

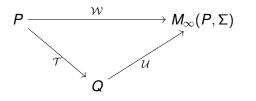
is the full normalization of  $\langle T, U \rangle$ . The construction of X(T, U) produces a weak hull embedding from T into X(T, U), which is what we want.

**Preliminaries** 

Definition of least branch hod pair

Comparison of east branch hod pairs

We have



Then

$$\mathcal{W} = X(\mathcal{T}, \mathcal{U})$$

is the full normalization of  $\langle T, U \rangle$ . The construction of X(T, U) produces a weak hull embedding from T into X(T, U), which is what we want.

Thus our Suslin representation verifies that  $\mathcal{T}$  is in the  $M_{\infty}$ -relevant part of  $\Sigma$  by producing a weak hull embedding of  $\mathcal{T}$  into  $\mathcal{W}$ .

Preliminaries

Definition of least branch hod pair

Comparison of east branch hod pairs

# Definition

Let P be an Ipm.

# (a) $\eta^{P}$ is the nonstrict sup of all lh(E), for *E* on the *P*-sequence.

### Preliminaries

Definition of least branch hod pair

Comparison of east branch hod pairs

# Definition

- Let P be an Ipm.
- (a)  $\eta^{P}$  is the nonstrict sup of all lh(E), for *E* on the *P*-sequence.
- (b) *P* has a top block iff there is a  $\kappa < \eta^{P}$  such that  $o(\kappa)^{P} = \eta^{P}$ .

### Preliminaries

Definition of least branch hod pair

Comparison of east branch hod pairs

# Definition

Let P be an Ipm.

- (a)  $\eta^{P}$  is the nonstrict sup of all lh(E), for *E* on the *P*-sequence.
- (b) *P* has a top block iff there is a κ < η<sup>P</sup> such that o(κ)<sup>P</sup> = η<sup>P</sup>. If so, then β<sup>P</sup> is the least such κ. We say β<sup>P</sup> begins the top block of P.

### Preliminaries

Definition of least branch hod pair

Comparison of east branch hod pairs

## Definition

Let P be an Ipm.

- (a)  $\eta^{P}$  is the nonstrict sup of all lh(E), for *E* on the *P*-sequence.
- (b) *P* has a top block iff there is a κ < η<sup>P</sup> such that o(κ)<sup>P</sup> = η<sup>P</sup>. If so, then β<sup>P</sup> is the least such κ. We say β<sup>P</sup> begins the top block of P.
- (c) Let  $\mathcal{T}$  be a normal tree on P with last model Q. We say  $\mathcal{T}$  is *short* iff P-to-Q drops, or for  $\pi \colon P \to Q$  the iteration map,  $\eta^Q < \pi(\eta^P)$ .

### Preliminaries

Definition of least branch hod pair

Comparison of east branch hod pairs

## Definition

Let P be an Ipm.

- (a)  $\eta^{P}$  is the nonstrict sup of all lh(E), for *E* on the *P*-sequence.
- (b) *P* has a top block iff there is a κ < η<sup>P</sup> such that o(κ)<sup>P</sup> = η<sup>P</sup>. If so, then β<sup>P</sup> is the least such κ. We say β<sup>P</sup> begins the top block of *P*.
- (c) Let  $\mathcal{T}$  be a normal tree on P with last model Q. We say  $\mathcal{T}$  is *short* iff P-to-Q drops, or for  $\pi \colon P \to Q$  the iteration map,  $\eta^Q < \pi(\eta^P)$ .

### Theorem

(AD<sup>+</sup>) Let  $(P, \Sigma)$  be an lbr hod pair with scope HC, and suppose P has a top block. Let  $\Psi$  be the restriction of  $\Sigma^{rel}$  to short trees, and  $\pi: P \to M_{\infty}(P, \Sigma)$  be the iteration map; then Code( $\Psi$ ) is  $\pi(\beta^{P})$ -Suslin, but not  $\alpha$ -Suslin for any  $\alpha < |\pi(\beta^{P})|$ .

### Preliminaries

Definition of least branch hod pair

Comparison of east branch hod pairs

We believe that under  $AD^+ + HPC$ , all Suslin cardinals  $\kappa$  arise in one of these two ways. That is, the set that is Suslin for the first time at  $\kappa$  is either a complete iteration strategy for an lpm, or a short tree strategy for an lpm.

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

We believe that under  $AD^+ + HPC$ , all Suslin cardinals  $\kappa$  arise in one of these two ways. That is, the set that is Suslin for the first time at  $\kappa$  is either a complete iteration strategy for an lpm, or a short tree strategy for an lpm.

This suggests proving HPC, assuming  $AD^+ + NLE$ , via an induction on Suslin cardinals, or equivalently, pointclasses with the Scale Property. Crossing gaps in scales is not actually a problem:

#### **Preliminaries**

Definition of least branch hod pair

Comparison of least branch hod pairs

### Theorem

Assume AD<sup>+</sup>, and let  $\Gamma$  be an inductive-like pointclass with the scale property. Suppose that the iteration strategies of lbr hod pairs are Wadge cofinal in  $\Gamma \cap \check{\Gamma}$ ; then

 (a) there is a short-tree-strategy pair (P, Ψ) such that Code(Ψ) is in Γ \ Γ, and

#### Preliminaries

Definition of least branch hod pair

Comparison of least branch hod pairs

### Theorem

Assume AD<sup>+</sup>, and let  $\Gamma$  be an inductive-like pointclass with the scale property. Suppose that the iteration strategies of lbr hod pairs are Wadge cofinal in  $\Gamma \cap \check{\Gamma}$ ; then

- (a) there is a short-tree-strategy pair  $(P, \Psi)$  such that  $Code(\Psi)$  is in  $\Gamma \setminus \check{\Gamma}$ , and
- (b) if all sets in Γ are Suslin, then there is an lbr hod pair (P,Σ) such that Code(Σ) is not in Γ.

#### **Preliminaries**

Definition of least branch hod pair

Comparison of least branch hod pairs

# **Determinacy models from hod pairs**

### Theorem

(Sargsyan,S.) Assume AD<sup>+</sup>, and that there is an lbr hod pair (P, Σ) such that P ⊨ ZFC + "δ is a Woodin limit of Woodin cardinals + "there are infinitely many Woodin cardinals above δ". Then there is a pointclass Γ such that
(1) L(Γ, ℝ) ⊨ "the largest Suslin cardinal exists, and belongs to the Solovay sequence" (LSA), and
(2) L(Γ, ℝ) ⊨ "if A is a set of reals that is OD(s) for some

 $s: \omega \rightarrow \theta$ , then A is Suslin and co-Suslin".

### **Preliminaries**

Definition of least branch hod pair

Comparison of least branch hod pairs

# **Determinacy models from hod pairs**

### Theorem

(Sargsyan,S.) Assume AD<sup>+</sup>, and that there is an lbr hod pair (P, Σ) such that P ⊨ ZFC + "δ is a Woodin limit of Woodin cardinals + "there are infinitely many Woodin cardinals above δ". Then there is a pointclass Γ such that
(1) L(Γ, ℝ) ⊨ "the largest Suslin cardinal exists, and belongs to the Solovay sequence" (LSA), and
(2) L(Γ, ℝ) ⊨ "if A is a set of reals that is OD(s) for some

 $s: \omega \rightarrow \theta$ , then A is Suslin and co-Suslin".

Part (1) is due to Sargsyan, and requires weaker hypotheses on *P*. The insight that Woodin limits of Woodins are what you need for (2) is due to Sargsyan.

### **Preliminaries**

Definition of least branch hod pair

Comparison of least branch hod pairs

# **Determinacy models from hod pairs**

### Theorem

(Sargsyan,S.) Assume AD<sup>+</sup>, and that there is an lbr hod pair (P, Σ) such that P ⊨ ZFC + "δ is a Woodin limit of Woodin cardinals + "there are infinitely many Woodin cardinals above δ". Then there is a pointclass Γ such that
(1) L(Γ, ℝ) ⊨ "the largest Suslin cardinal exists, and belongs to the Solovay sequence" (LSA), and
(2) L(Γ, ℝ) ⊨ "if A is a set of reals that is OD(s) for some

(2)  $L(\Gamma, \mathbb{R}) \models$  "if A is a set of reals that is OD(s) for some  $s: \omega \to \theta$ , then A is Suslin and co-Suslin".

Part (1) is due to Sargsyan, and requires weaker hypotheses on *P*. The insight that Woodin limits of Woodins are what you need for (2) is due to Sargsyan. Part (2) is a step toward a model of  $AD_{\mathbb{R}}$  that satisfies " $\omega_1$ is *X*-supercompact, for all sets *X*".

### **Preliminaries**

Definition of least branch hod pair

Comparison of east branch hod pairs