

## Letter 5

### The term model as a model of set theory

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## §1 Introduction

We are currently working in the theory  $\mathbf{ZFC} + \mathbf{V=L} +$  “There is an inaccessible cardinal”. We let  $\theta$  be an inaccessible cardinal. Our goal is to prove the consistency of  $\mathbf{NFU}_*$

Let  $\alpha < \theta$  be a lbf. In the previous letter, we have associated to  $\alpha$  a term language  $\mathcal{L}$ .

A term model [say  $\mathcal{M}$ ] is given by putting an equivalence relation [say  $\equiv$ ] on the closed terms of  $\mathcal{L}$  which satisfies the analogues of the usual axioms of equality. [This is fully discussed in letter 4.] The goal of this letter is to describe a condition on  $\mathcal{M}$  [being *well-instantiated*] which allows us to make  $\mathcal{M}$  into a model of set-theory.

To be quite pedantic, we will construct a model of a certain version of set-theory,  $T$  [to be described precisely in a moment] whose underlying set is the collection of equivalence classes of the equivalence relation  $\equiv$ . But we shall, in the future, abuse notation, and refer to the model thus created as  $\mathcal{M}$  also.

### 1. 1 The theory $T$

The theory  $T$  will have the following axioms:

1. The extensionality, foundation, infinity, union, power set, choice and pairing axioms.
2. The  $\Sigma_1$  replacement axiom. This says that if  $\phi(x, y)$  is a  $\Sigma_1$  relation which is “functional” [that is, if  $\phi(x, y_1)$  and  $\phi(x, y_2)$  then  $y_1 = y_2$ ] and  $A$  is a set then  $\{y \mid (\exists x \in A)\phi(x, y)\}$  is a set. [Notice we do *not* require that  $\forall x \in A \exists y \phi(x, y)$ .]
3.  $\mathbf{V=L}$ .
4. There are arbitrarily large lbf’s.

This is a fairly strong set-theory. [Though it is much weaker than **ZFC**.] Certainly, almost all of mathematics as it is done by the “mathematician in the street” can be carried out in this theory.

### 1. 2 An addition to letter 4

In letter 4 we made intuitive remarks about the purposes of the various function symbols of the language  $\mathcal{L}$ . We wish to make the following additional remark:

**Remark 1.1** It is our intention that the functions  $h_i$  will have the following properties.

1. If  $x$  is not an ordinal, then  $h_i(x) = 0$ .
2. If  $\beta$  is an ordinal,  $h_i(\beta) \leq \beta$ .

The purpose of this condition will be to insure that the  $\xi_i$  are cofinal in the ordinals of the term model  $\mathcal{M}$

## §2 Well-instantiated term models

Let  $\mathcal{M}$  be a term model of  $\mathcal{L}$ . We are going to introduce the notion of  $\mathcal{M}$  being *well-instantiated*. This will serve to insure that  $\mathcal{M}$  does indeed conform to the intuitive remarks of the preceding letter 4.

### 2. 1

We begin with a series of preliminary definitions. Let  $\alpha_1 < \alpha$ . Let  $n \in \omega$ . We say that a closed term  $\tau$  of  $\mathcal{L}$  has *rank at most*  $(\alpha_1, n)$  iff:

1. If  $\bar{\gamma}$  occurs in  $\tau$ , then  $\gamma < \alpha_1$ .
2. If  $h_i$  occurs in  $\tau$ , then  $i < n$ .
3. If  $\xi_i$  occurs in  $\tau$ , then  $|i| \leq n$ .

### 2. 2

An  $(\alpha_1, n)$  instantiation model  $M$  consists of the following:

1. The underlying set of  $M$  is  $L_\lambda$  for some lbf  $\lambda$  such that  $\alpha_1 < \lambda$ .
2. There are lbfps  $\beta_{-n} \dots \beta_n$  such that:

$$\alpha_1 < \beta_{-n} < \dots < \beta_n < \lambda.$$

3. For  $i < n$  there is a map  $H_i : L_\lambda \mapsto L_\lambda$  which satisfies the requirements discussed in Remark 1.1.

**2. 3**

Suppose that  $\tau$  is a closed term of  $\mathcal{L}$  of rank at most  $(\alpha_1, n)$ . We show how to associate a value to  $\tau$ ,  $\tau^M$ , lying in  $L_\lambda$ .

For this it suffices to give the interpretation of all the function symbols and constants appearing in  $\tau$ . We do this as follows:

1. If  $\gamma < \alpha_1$ , then the interpretation of  $\bar{\gamma}$  is  $\gamma$ .
2. The interpretation of  $\xi_i$  is  $\beta_i$  (for  $-n \leq i \leq n$ ).
3. The interpretation of  $f_{n,i}$  is as discussed in letter 4.
4. Let  $i < n$ . Then the interpretation of  $h_i$  is given by  $H_i$ .

**2. 4**

Let  $\alpha_1 < \alpha$ . Let  $n \in \omega$ . Let  $M$  be an  $(\alpha_1, n)$  instantiation model. Then  $M$  *instantiates*  $\mathcal{M}$  if the following holds:

Let  $\tau_1$  and  $\tau_2$  be closed terms of  $\mathcal{L}$  of rank at most  $(\alpha_1, n)$ . Then  $\tau_1^M = \tau_2^M$  iff  $\tau_1 \equiv \tau_2$  in the term model  $\mathcal{M}$ .

**2. 5**

We now define what it means for  $\langle M, Y \rangle$  to be an  $(\alpha_1, n)$  *pre-instantiation model*.

1.  $M$  is a first-order structure for the language of set-theory. The underlying set of  $M$  is  $L_\lambda$  for some lbf  $\lambda$  such that  $\alpha_1 < \lambda$ .
2.  $Y$  is a subset of  $\lambda$  consisting of lbfs. The order-type of  $Y$  is a limit ordinal.
3. For  $i < n$  there is a map  $H_i : L_\lambda \mapsto L_\lambda$  which satisfies the requirements discussed in Remark 1.1.

**2. 6**

Finally, we can say what it means for  $\mathcal{M}$  to be well-instantiated:

For every limit ordinal  $\alpha_1 < \alpha$ , for every  $n \in \omega$ , there is a  $(\alpha_1, n)$  pre-instantiation model  $\langle M_1, Y \rangle$  such that if

$$\beta_{-n} < \dots < \beta_n$$

are an increasing  $2n + 1$ -tuple of elements of  $Y$ , then  $\langle M_1, \beta_{-n}, \dots, \beta_n \rangle$  instantiates  $\mathcal{M}$ .

### §3 Making $\mathcal{M}$ into a model of set-theory.

Some terminology: I shall refer to the  $f_{n,i}$ 's as *local functions*. As I have already remarked this conflicts with my terminology in my proof that NFUA yields  $n$ -Mahlos.

We are going to associate a model of the theory  $T$  to the term model  $\mathcal{M}$ . We have already said that the underlying set of this model will be the set of  $\equiv$  equivalence classes.

#### 3. 1

Next, we have to define the  $\epsilon$ -relation of the model. Let  $g_1$  be the binary function [one of the  $f_{2,i}$ 's] that does the following: If  $x \in y$ , then  $g_1(x, y) = 1$ . Otherwise,  $g_1(x, y) = 0$ .

Let  $\tau$  be a closed term of  $\mathcal{L}$ . We let  $[\tau]$  be the equivalence class of  $\tau$  with respect to the relation  $\equiv$  of  $\mathcal{M}$ . The various function symbols of  $\mathcal{L}$  determine, in an obvious way, functions which act on the set of equivalence classes. We use the same notation for these derived functions as for the corresponding symbols.

Let  $\tau_1$  and  $\tau_2$  be closed terms. We put  $[t_1] \epsilon_{\mathcal{M}} [\tau_2]$  iff  $g_1([\tau_1], [\tau_2]) = 1$ . [By abuse of language, for  $\gamma < \alpha$ , we write  $\gamma$  for  $[\bar{\gamma}]$ .]

#### 3. 2

It will take us a while to establish that  $T$  holds in the model just described. Before taking the first step along the way, we need to nail down precisely what we mean by  $\Sigma_0$ -formulas. These are the class of formulas inductively defined by the following requirements:

1. If  $v$  and  $w$  are variables then " $v \in w$ " and " $v = w$ " are  $\Sigma_0$  formulas.
2. If  $\psi$  is a  $\Sigma_0$  formula, then so is  $\neg\psi$ .
3. If  $\psi_1$  and  $\psi_2$  are  $\Sigma_0$  formulas so is  $\psi_1 \vee \psi_2$ .
4. If  $v$  and  $w$  are distinct variables and  $\psi$  is a  $\Sigma_0$  formula, then so is  $(\exists v \in w)\psi$ .

It is no real loss of generality to require in clause 4 that neither  $v$  or  $w$  appears bound in  $\psi$ . We shall assume this in what follows.

**3. 3**

**Lemma 3.1** Let  $\mathcal{M}$  be well-instantiated. Let  $\psi(v_0, \dots, v_m)$  be a  $\Sigma_0$  formula. Let  $\alpha_1$  be a limit ordinal less than  $\alpha$  and let  $n_1 \in \omega$ .

We suppose that  $\tau_0, \dots, \tau_r$  are closed terms of  $\mathcal{L}$  of rank at most  $(\alpha_1, n)$ . Let  $M$  be an  $(\alpha_1, n_1)$  instantiation model for  $\mathcal{M}$ .

Then the following are equivalent:

1.  $\mathcal{M} \models \psi([\tau_0], \dots, [\tau_r])$ .
2.  $M \models \psi(\tau_0^M, \dots, \tau_r^M)$ .

**Proof:** The proof proceeds by induction on the length of  $\psi$ . We shall only consider the case when  $\psi$  has the form  $(\exists v \in v_0)\chi(v, v_0, v_1, \dots, v_r)$ . The other cases [when  $\psi$  is atomic or a boolean combination of shorter formulas] will be left to the reader.

We fix  $\alpha_1 < \alpha$  and  $n_1 \in \omega$ . Let  $\tau_0, \dots, \tau_r$  be closed subterms of  $\mathcal{L}$  of rank at most  $(\alpha_1, n_1)$ . Let  $M$  be an  $(\alpha_1, n_1)$  instantiation model for  $\mathcal{M}$ .

This case reduces to two subclaims.

**Claim 3.1.A** Suppose that  $M \models \psi(\tau_0^M, \dots, \tau_r^M)$ . Then  $\mathcal{M} \models \psi([\tau_0], \dots, [\tau_r])$ .

**Proof:** Let  $g_2(x_0, \dots, x_r)$  be a local function with the following property:

1. If  $(\exists x \in a_0)\chi(x, a_0, \dots, a_r)$  then  $g_2(a_0, \dots, a_r)$  is the  $L$ -least such  $x$ .
2. Otherwise,  $g_2(a_0, \dots, a_r) = 0$ .

Let  $\tau^*$  be the term  $g_2(\tau_0, \dots, \tau_r)$ . Then clearly  $\tau^*$  has rank at most  $(\alpha_1, n_1)$ .

Moreover, by the hypothesis of this claim and the definition of  $g_2$  we have

$$M \models [\chi(\tau^*, \tau_0, \dots, \tau_r) \wedge \tau^* \in \tau_0].$$

Using our inductive hypothesis, we conclude that

$$\mathcal{M} \models [\chi([\tau^*], [\tau_0], \dots, [\tau_r]) \wedge [\tau^*] \in [\tau_0]].$$

Hence  $\mathcal{M} \models \psi([\tau_0], \dots, [\tau_r])$ .  $\square_{\text{Claim}}$

**Claim 3.1.B** Suppose that  $\mathcal{M} \models \psi([\tau_0], \dots, [\tau_r])$ . Then  $M \models \psi(\tau_0^M, \dots, \tau_r^M)$ .

**Proof:** Let  $\tau^*$  be a closed term such that

$$\mathcal{M} \models [\chi(\tau^*, \tau_0, \dots, \tau_r) \wedge \tau^* \in \tau_0].$$

We choose  $\alpha_2 < \alpha$  and  $n_2 \in \omega$  such that  $\tau^*, \tau_0, \dots, \tau_r$  are closed terms of rank at most  $(\alpha_2, n_2)$ . Let  $M_2$  be an  $(\alpha_2, n_2)$  instantiation model for  $\mathcal{M}$ .

Let  $g_3(x_0, \dots, x_r)$  be a local function with the following properties:

1.  $g_3(x_0, \dots, x_r) \in \{0, 1\}$  for any choice of  $x_0, \dots, x_r$ .
2.  $g_3(x_0, \dots, x_r) = 1$  iff  $\psi(x_0, \dots, x_r)$ .

Let  $\tau^\#$  be the closed term  $g_3(\tau_0, \dots, \tau_r)$ . Then  $\tau^\#$  has rank at most  $(\alpha_1, n_1)$ .

By our inductive hypothesis,  $M_2 \models [\chi(\tau^\#^{M_2}, \tau_0^{M_2}, \dots, \tau_r^{M_2}) \wedge \tau^\#^{M_2} \in \tau_0^{M_2}]$ .

Hence  $M_2 \models \psi(\tau_0^{M_2}, \dots, \tau_r^{M_2})$ . It follows that in  $M_2$ ,  $g_3((\tau_0^{M_2}, \dots, \tau_r^{M_2})) = 1$ .

Since  $M_2$  is an  $(\alpha_2, n_2)$  instantiation model for  $\mathcal{M}$ , we have  $g_3(\tau_0^{M_2}, \dots, \tau_r^{M_2}) \equiv \bar{1}$  in  $\mathcal{M}$ .

Since  $M$  is an  $(\alpha_1, n_1)$  instantiation model for  $\mathcal{M}$  we have, in  $M$ ,

$$g_3(\tau_0^M, \dots, \tau_r^M) = 1.$$

Hence,  $M \models \psi(\tau_0^M, \dots, \tau_r^M)$ .  $\square_{\text{Claim}}$

$\square_{\text{Lemma}}$

### 3. 4

The proofs of the following lemmas will be left to the reader. Their proofs employ both the statement of Lemma 3.1 and ideas used in its proof. [Using local functions as Skolem functions; checking assertions by using auxiliary models [like  $M_2$  in the proof of Claim 3.1.B.]]

**Lemma 3.2** *The axioms of foundation, extensionality, pairing, union, power set and infinity hold in  $\mathcal{M}$ .*

### 3. 5

It is well-known that there is a sentence  $\sigma$  such that:

1. If  $\lambda$  is an uncountable cardinal, then  $L_\lambda$  models  $\sigma$ .
2. If  $x$  is a transitive set such that  $\sigma$  holds in  $x$ , then  $x$  has the form  $L_\lambda$  for some limit ordinal  $\lambda$ .

**Remark.** I have stated a version of this result which suffices for our applications and which is easy to prove. In fact, one can arrange that the transitive models of  $\sigma$  are precisely the sets  $L_\lambda$  for some ordinal  $\lambda > 0$ . But this takes considerably more delicate arguments.

When I say  $\mathbf{V=L}$  in the following lemma, I mean: every set  $x$  is a member of a transitive set  $y$  such that  $y$  models  $\sigma$ .

**Lemma 3.3**  $\mathbf{V=L}$  holds in  $\mathcal{M}$ .

**Corollary 3.4** *The axiom of choice holds in  $\mathcal{M}$ .*

**3. 6**

In checking that an ordinal  $\beta$  is a lbf, it suffices to see that a certain sentence [which we will not construct in detail] holds in the power set of  $L_\beta$ . Using this and Remark 1.1, the following is easy to prove:

**Lemma 3.5** *The elements  $[\xi_i]$  (for  $i \in \omega$ ) are lbf's in  $\mathcal{M}$  and are cofinal in the ordinals of  $\mathcal{M}$ .*

**3. 7**

Using the previous lemmas and the fact that for  $\beta, \beta'$  bfp's with  $\beta < \beta'$ ,  $\Sigma_1$  formulas are absolute between  $L_\beta$  and  $L_{\beta'}$ , it is easy to prove the following lemma:

**Lemma 3.6**  *$\Sigma_1$ -replacement holds in  $\mathcal{M}$ .*

We have proved:

**Theorem 3.7**  *$\mathcal{M}$  is a model of  $T$ .*

This ends letter 5.