

LECTURE 8

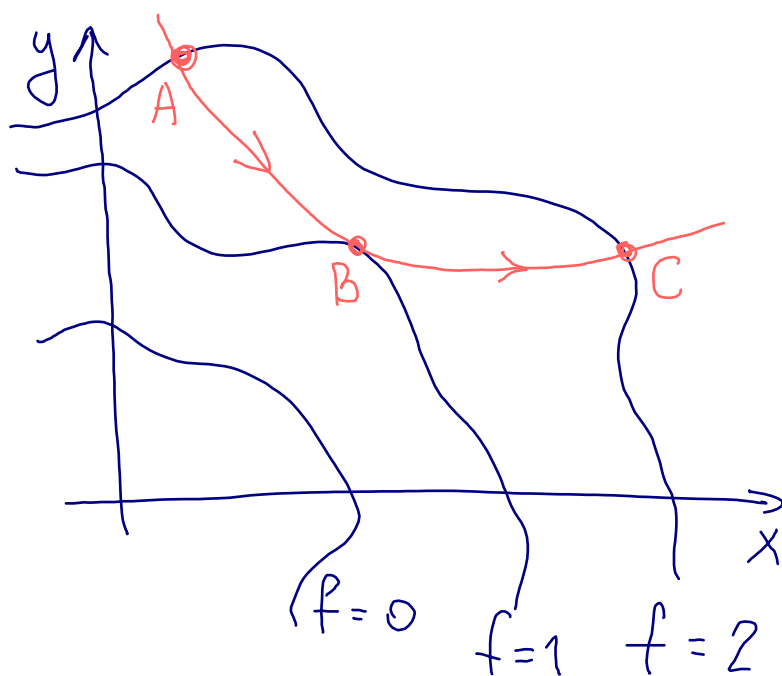
Here we do two examples and applications of what we learned so far

§8.1. Functions restricted to curves

Exercise: We are given

- level curves of a function $f(x, y)$

- a curve C with 3 points marked: A, B, C



- ① Draw the direction of ∇f at the points A, B, C
(you cannot get the magnitude from the plot)
- ② If C is parametrized as $(x, y) = (x(t), y(t))$ where t grows in the direction of the arrow, is the function $f(x(t), y(t))$ increasing/decreasing/has a local max/min at the values of t corresponding to A, B, C ?

Solution:

① ∇f has to be orthogonal to the level curves and point in the direction of increasing f

(Recall: the directional der'ive $f=0$ $f=1$ $f=2$)

$$D_{\vec{v}} f(x,y) = \nabla f(x,y) \cdot \vec{v}$$

② Recall the Chain Rule:

$$D_t(f(x(t), y(t))) = \nabla f(x(t), y(t)) \cdot \vec{v}(t)$$

where $\vec{v}(t) = (x'(t), y'(t))$ is the velocity vector.

• At A, ∇f & \vec{v} are at an obtuse angle

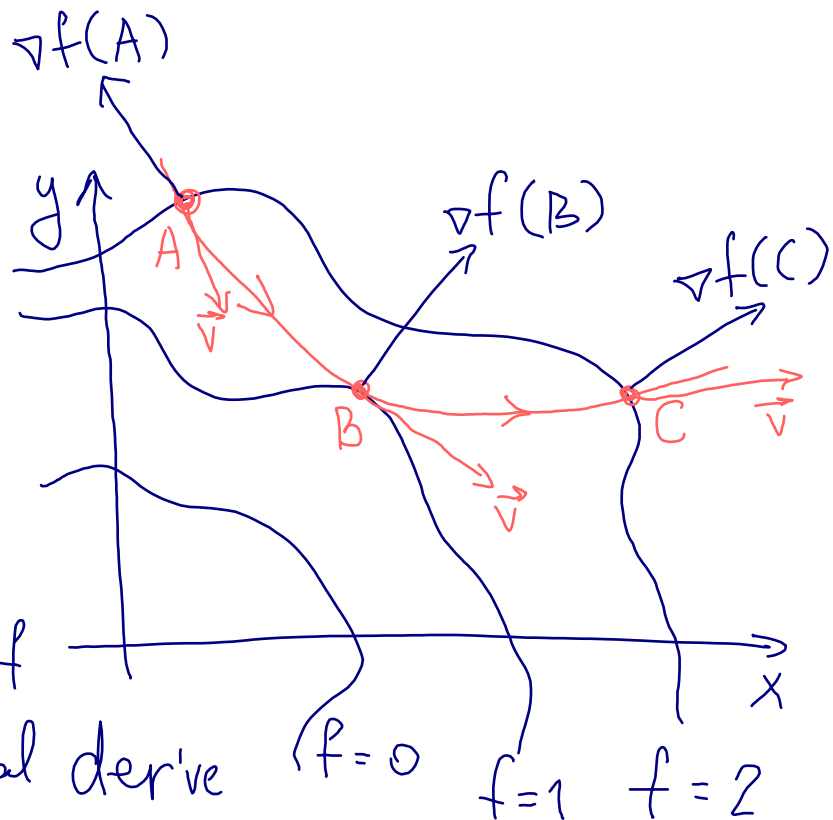
So $D_t(\dots) < 0 \Rightarrow f$ decreasing on \mathcal{C}

• At B, $\nabla f \perp \vec{v}$. So $D_t(\dots) = 0$

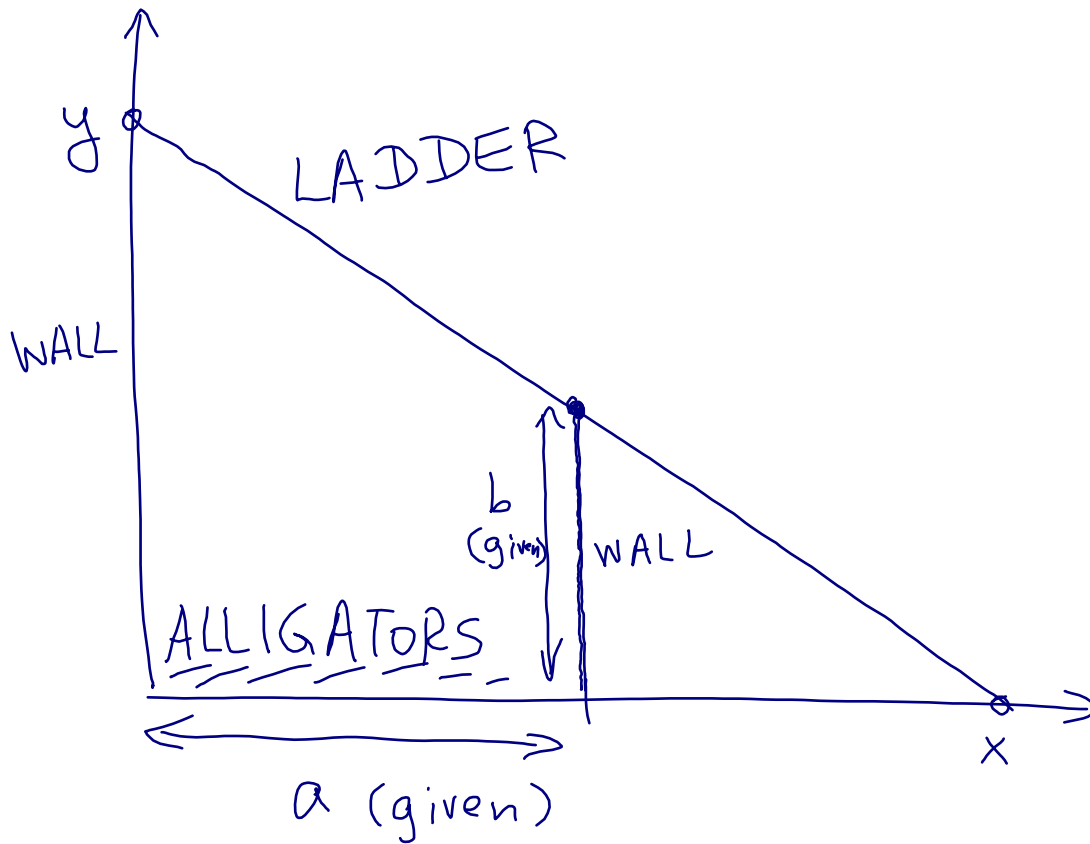
Note: $f(B) = 1$ but $f > 1$ elsewhere on \mathcal{C} (look at level lines). So get a local min on \mathcal{C}

• At C, ∇f & \vec{v} are at an acute angle

So $D_t(\dots) > 0 \Rightarrow f$ increasing on \mathcal{C}



§ 8.2. Alligator moat (exercise 13.9.63 in the book)



Exercise: find the minimal possible length of the ladder by setting up a constrained optimization problem on (x, y)



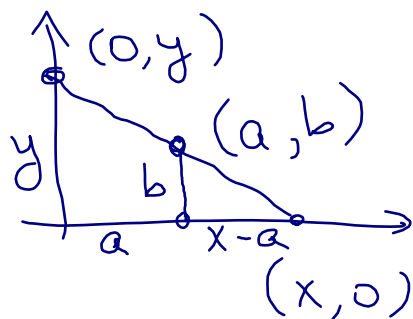
Solution part 1: set up the problem.

$f(x,y) = x^2 + y^2$ is the square of the length of the ladder.
We want to minimize f .

What is the constraint?

There should be a line (the ladder) passing through $(x,0)$, $(0,y)$ and (a,b)

This means:



$$\frac{y}{x} = \frac{b}{x-a}$$

Can simplify to $(x-a)y = bx$

That is, $\boxed{\frac{a}{x} + \frac{b}{y} = 1}$

Note: we restrict to $x > a$, $y > b$ out of geometric considerations. If $x \rightarrow a$ then $y \rightarrow \infty$ and $f(x,y) \rightarrow \infty$ and same for $y \rightarrow b$. So it's enough to find all LOCAL minima.

Solution part 2: Solving the optimization problem:

$$\text{minimize } \boxed{f(x, y) = x^2 + y^2}$$

under the constraint

$$g(x, y) = 0 \quad \text{where } \boxed{g(x, y) = \frac{a}{x} + \frac{b}{y} - 1}$$

Use Lagrange multipliers to find local extrema:

$$\begin{cases} f_x(x, y) = \lambda g_x(x, y) \\ f_y(x, y) = \lambda g_y(x, y) \end{cases}$$

$$\begin{cases} 2x = -\frac{\lambda a}{x^2} \\ 2y = -\frac{\lambda b}{y^2} \end{cases} \Leftrightarrow \begin{cases} x = \sqrt[3]{-\frac{\lambda a}{2}} \\ y = \sqrt[3]{-\frac{\lambda b}{2}} \end{cases}$$

To simplify, denote $\boxed{C = \sqrt[3]{-\frac{\lambda}{2}}}$, then

$$x = C \sqrt[3]{a}, \quad y = C \sqrt[3]{b}$$

Use the constraint: $1 = \frac{a}{x} + \frac{b}{y} = \frac{1}{C} (a^{2/3} + b^{2/3})$

$$\text{Thus } \boxed{C = a^{2/3} + b^{2/3}}$$

And we get

$$x = c\sqrt[3]{a} = a^{1/3} (a^{2/3} + b^{2/3})$$

$$y = c\sqrt[3]{b} = b^{1/3} (a^{2/3} + b^{2/3})$$

Only 1 extremum point, so (using Note at the end of part 1)
the minimum of f is

$$f(x,y) = x^2 + y^2 = (a^{2/3} + b^{2/3})^3$$

So the minimal length of the ladder is

$$\sqrt{f(x,y)} = (a^{2/3} + b^{2/3})^{3/2}.$$